Using REXX and z/OS UNIX System Services
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About this document

This document presents the information you need to write REXX™ programs that access z/OS UNIX System Services (z/OS UNIX). It describes the features and usage requirements for the z/OS UNIX REXX extensions, or syscall commands, which are interfaces between the z/OS operating system and the functions specified in the POSIX.1 standard (ISO/IEC 9945-1:1990[E] IEEE Std 1003.1-1990: First edition 1990-12-07; Information technology—Portable Operating System Interface [ POSIX] Part 1; System Application Program Interface [API] [C Language]). These functions are used by z/OS UNIX. This document also describes syscall commands that are not related to the standards.

Who should read Using REXX and z/OS UNIX System Services

This document is for programmers who are already familiar with the REXX language and experienced with the workings of TSO/E and z/OS UNIX. It describes how to include in a REXX program syscall commands that access z/OS UNIX services.

BPX messages from the REXX processor are documented in z/OS MVS System Messages, Vol 3 (ASB-BPX). You can also access the messages directly from the LookAt Web site at www.ibm.com/servers/eserver/zseries/zos/bkserv/lookat/.

Where to find more information

Where necessary, this document references information in other documents about the elements and features of the z/OS® system. For complete titles and order numbers for all z/OS documents, see z/OS Information Roadmap.

Direct your request for copies of any IBM® publication to your IBM representative or to the IBM branch office serving your locality.

There is also a toll-free customer support number (1-800-879-2755) available Monday through Friday from 6:30 a.m. through 5:00 p.m. Mountain Time. You can use this number to:
• Order or inquire about IBM publications
• Resolve any software manufacturing or delivery concerns
• Activate the program reorder form to provide faster and more convenient ordering of software updates

Softcopy publications

The z/OS UNIX® library is available on the z/OS Collection Kit, SK2T-6700. This softcopy collection contains a set of z/OS and related unlicensed product documents. The CD-ROM collection includes the IBM Library Reader™, a program that enables customers to read and print the softcopy documents.

PDF versions of the publications are available at http://www.ibm.com/systems/z/os/zos/bkserv/ Use Adobe® Reader to view them.

IBM Systems Center publications

IBM Systems Centers produce IBM Redbooks® publications that can be helpful in setting up and using z/OS UNIX. See the IBM Redbooks site at http://www.ibm.com/redbooks
These documents have not been subjected to any formal review nor have they been checked for technical accuracy, but they represent current product understanding at the time of their publication and provide information on a wide range of topics. You must order them separately. A selected list of these documents is on the z/OS UNIX Web site at http://www.ibm.com/servers/eserver/zseries/zos/unix/bpxa1pub.html/.

Porting information for z/OS UNIX

A Porting Guide is available at http://www.ibm.com/servers/eserver/zseries/zos/unix/bpxa1por.html. It covers a range of useful topics, including sizing a port, setting up a porting environment, ASCII-EBCDIC issues, performance, and much more.

The porting page also features a variety of porting tips and lists porting resources that will help you in your port.

z/OS UNIX courses

For a current list of courses that you can take, go to http://www.ibm.com/services/learning/.

You can also ask your IBM representative or call 1-800-IBM-TEACH (1-800-426-8322).

z/OS UNIX home page


Some of the tools available from the Web site are ported tools, and some are unsupported tools designed for z/OS UNIX. The code works in our environment at the time we make it available, but is not officially supported. Each tool has a README file that describes the tool and lists any restrictions.

The simplest way to reach these tools is through the z/OS UNIX home page. From the home page, click on Tools and Toys.

The code is also available from ftp://ftp.software.ibm.com/s390/zos/unix/ through anonymous FTP.

**Restrictions**

Because the tools are not officially supported, APARs cannot be accepted.

Discussion list

Customers and IBM participants also discuss z/OS UNIX on the mvs-oe discussion list. This list is not operated or sponsored by IBM.

To subscribe to the mvs-oe discussion, send a note to:
listserv@vm.marist.edu

Include the following line in the body of the note, substituting your given name and family name as indicated:
subscribe mvs-oe given_name family_name
After you are subscribed, you will receive further instructions on how to use the mailing list.

The z/OS Basic Skills Information Center

The z/OS Basic Skills Information Center is a Web-based information resource intended to help users learn the basic concepts of z/OS, the operating system that runs most of the IBM mainframe computers in use today. The Information Center is designed to introduce a new generation of Information Technology professionals to basic concepts and help them prepare for a career as a z/OS professional, such as a z/OS system programmer.

Specifically, the z/OS Basic Skills Information Center is intended to achieve the following objectives:

- Provide basic education and information about z/OS without charge
- Shorten the time it takes for people to become productive on the mainframe
- Make it easier for new people to learn z/OS.

To access the z/OS Basic Skills Information Center, open your Web browser to the following Web site, which is available to all users (no login required):

http://publib.boulder.ibm.com/infocenter/zos/basics/index.jsp

Finding more information about REXX


How to read syntax diagrams

This section describes how to read syntax diagrams. It defines syntax diagram symbols, items that may be contained within the diagrams (keywords, variables, delimiters, operators, fragment references, operands) and provides syntax examples that contain these items.

Syntax diagrams pictorially display the order and parts (options and arguments) that comprise a command statement. They are read from left to right and from top to bottom, following the main path of the horizontal line.

Symbols

The following symbols may be displayed in syntax diagrams:

<table>
<thead>
<tr>
<th>Symbol</th>
<th>Definition</th>
</tr>
</thead>
<tbody>
<tr>
<td>➔</td>
<td>Indicates the beginning of the syntax diagram.</td>
</tr>
<tr>
<td>➔</td>
<td>Indicates that the syntax diagram is continued to the next line.</td>
</tr>
<tr>
<td>➔</td>
<td>Indicates that the syntax is continued from the previous line.</td>
</tr>
<tr>
<td>➔</td>
<td>Indicates the end of the syntax diagram.</td>
</tr>
</tbody>
</table>

Syntax items

Syntax diagrams contain many different items. Syntax items include:

- Keywords - a command name or any other literal information.
- Variables - variables are italicized, appear in lowercase, and represent the name of values you can supply.
• Delimiters - delimiters indicate the start or end of keywords, variables, or operators. For example, a left parenthesis is a delimiter.
• Operators - operators include add (+), subtract (-), multiply (*), divide (/), equal (=), and other mathematical operations that may need to be performed.
• Fragment references - a part of a syntax diagram, separated from the diagram to show greater detail.
• Separators - a separator separates keywords, variables or operators. For example, a comma (,) is a separator.

Note: If a syntax diagram shows a character that is not alphanumeric (for example, parentheses, periods, commas, equal signs, a blank space), enter the character as part of the syntax.

Keywords, variables, and operators may be displayed as required, optional, or default. Fragments, separators, and delimiters may be displayed as required or optional.

<table>
<thead>
<tr>
<th>Item type</th>
<th>Definition</th>
</tr>
</thead>
<tbody>
<tr>
<td>Required</td>
<td>Required items are displayed on the main path of the horizontal line.</td>
</tr>
<tr>
<td>Optional</td>
<td>Optional items are displayed below the main path of the horizontal line.</td>
</tr>
<tr>
<td>Default</td>
<td>Default items are displayed above the main path of the horizontal line.</td>
</tr>
</tbody>
</table>

## Syntax examples

The following table provides syntax examples.

<table>
<thead>
<tr>
<th>Item</th>
<th>Syntax example</th>
</tr>
</thead>
<tbody>
<tr>
<td>Required item.</td>
<td>Required items appear on the main path of the horizontal line. You must specify these items.</td>
</tr>
<tr>
<td>Required choice.</td>
<td>A required choice (two or more items) appears in a vertical stack on the main path of the horizontal line. You must choose one of the items in the stack.</td>
</tr>
<tr>
<td>Optional item.</td>
<td>Optional items appear below the main path of the horizontal line.</td>
</tr>
<tr>
<td>Optional choice.</td>
<td>An optional choice (two or more items) appears in a vertical stack below the main path of the horizontal line. You may choose one of the items in the stack.</td>
</tr>
</tbody>
</table>
Table 1. Syntax examples (continued)

<table>
<thead>
<tr>
<th>Item</th>
<th>Syntax example</th>
</tr>
</thead>
<tbody>
<tr>
<td>Default.</td>
<td>![Default Example Diagram]</td>
</tr>
</tbody>
</table>

Default items appear above the main path of the horizontal line. The remaining items (required or optional) appear on (required) or below (optional) the main path of the horizontal line. The following example displays a default with optional items.

| Variable. | ![Variable Example Diagram] |

Variables appear in lowercase italics. They represent names or values.

| Repeatable item. | ![Repeatable Item Example Diagram] |

An arrow returning to the left above the main path of the horizontal line indicates an item that can be repeated.

A character within the arrow means you must separate repeated items with that character.

An arrow returning to the left above a group of repeatable items indicates that one of the items can be selected, or a single item can be repeated.

| Fragment. | ![Fragment Example Diagram] |

The fragment symbol indicates that a labeled group is described below the main syntax diagram. Syntax is occasionally broken into fragments if the inclusion of the fragment would overly complicate the main syntax diagram.
Summary of changes

Summary of changes for SA22-7806-12
Version 1 Release 11

This document contains information previously in z/OS Using REXX and z/OS UNIX System Services SA22-7806-11, which supports z/OS Version 1 Release 10.

New information
- A new variable, MTM_SAMEMODE, has been added to "unmount" on page 191.
- A new common key, SUBSYS, has been added to "Requesting dynamic allocation" on page 249.
- Two syscall commands, setreuid and setregid, have been added to Chapter 3, "The syscall commands," on page 19.

Changed information
- The variable MNTE_MODE has been changed in "getmntent" on page 99.

Deleted information
- The REXX-related BPX messages are now in z/OS MVS System Messages, Vol 3 (ASB-BPX).

You may notice changes in the style and structure of some content in this document—for example, headings that use uppercase for the first letter of initial words only, and procedures that have a different look and format. The changes are ongoing improvements to the consistency and retrievability of information in our documents.

This document contains terminology, maintenance, and editorial changes. Technical changes or additions to the text and illustrations are indicated by a vertical line to the left of the change.

Summary of changes for SA22-7806-11
Version 1 Release 10

This document contains information previously in z/OS Using REXX and z/OS UNIX System Services SA22-7806-10, which supports refresh of z/OS Version 1 Release 8.

New information
- New error message BPXW9092I has been added to the appendix.

Changed information
- For getmntent, new variables have been added to the MNTE parameter.

This document contains terminology, maintenance, and editorial changes, including changes to improve consistency and retrievability.
Chapter 1. Using TSO/E REXX for z/OS UNIX processing

A REXX program is recognized by the word REXX (not case-sensitive) as the first word on the first line and within a REXX comment. For example, the following is a simple REXX program:

```rexx
/* rexx */
say 'hello world'
```

Restriction: Blank spaces cannot precede the opening REXX comment symbol (/*).

The set of z/OS UNIX extensions to the TSO/E Restructured Extended Executor (REXX) language enable REXX programs to access z/OS UNIX callable services. The z/OS UNIX extensions, called syscall commands, have names that correspond to the names of the callable services they invoke—for example, access, chmod, and chown.

You can run an interpreted or compiled REXX program with syscall commands from TSO/E, from MVS™ batch, from the z/OS shells, or from a program. You can run a REXX program with syscall commands only on a system with z/OS UNIX System Services installed. For a complete description of each of the syscall commands, see Chapter 3, “The syscall commands,” on page 19.

The set of z/OS UNIX REXX functions also extend the REXX language on z/OS in the z/OS UNIX environment. There are functions that provide:
- Standard REXX I/O
- Access to some common file services and environment variables
All of the z/OS UNIX functions, except bpxwunix() and syscalls(), must be run in a z/OS UNIX environment. For a complete description of each of the functions, see Chapter 5, “z/OS UNIX REXX functions,” on page 215.

For dynamic allocation and dynamic output, BPXWDYN is a text interface to a subset of the SVC 99 and SVC 109 services that is designed to be called from REXX. For a complete description of BPXWDYN, see Chapter 6, “BPXWDYN: a text interface to dynamic allocation and dynamic output,” on page 245.

Host command environments for z/OS UNIX processing

Host command environments and external function packages that are available in the MVS REXX environment can be used by a REXX program that has z/OS UNIX extensions. These additional host command environments are also available:

SYSCALL For a REXX program with syscall commands that will be run from TSO/E or MVS batch, you need to initialize the environment by beginning a REXX program with a syscalls('ON') call.

SH For a REXX program with syscall commands that will be run from a z/OS shell or from a program, SH is the initial host environment. The SYSCALL environment is automatically initialized as well, so you do not need to begin the REXX program with a syscalls('ON') call. Syscall commands within the REXX program (for example, chmod) are interpreted as z/OS shell commands, not as syscall commands.

TSO A REXX program can run TSO/E commands, but you cannot use TSO commands to affect your REXX environment, or have REXX...
statements or other host command environments affect your TSO process. Commands that are addressed to TSO will be run in a TMP running in a separate address space and process from your REXX program. The TSO process is started when the first TSO command is run, and persists until your REXX program terminates or you run the TSO LOGOFF command.

When a REXX program is run from a z/OS shell or from a program, both the SH and SYSCALL host command environments are available to it. When a REXX program is run from TSO/E or MVS batch, only the SYSCALL environment is available.

For background information on the concept of a host command environment, see [z/OS TSO/E User's Guide](#).

### The SYSCALL environment

The SYSCALL environment can be used by any REXX program with syscall commands, whether it runs from TSO/E or the z/OS shells (where the environment is automatically initialized).

### Running a REXX program from TSO/E or MVS batch

To run a REXX program with syscall commands from TSO/E or MVS batch, use the `syscalls('ON')` function at the beginning of the REXX program. This function:

- Ensures that the SYSCALL command environment (ADDRESS syscall) is established.
- Ensures that the address space is a process; this is known as *dubbing*.
- Initializes the predefined variables in the current REXX variable pool.
- Sets the signal process mask to block all signals that can be blocked. See "Using the REXX signal services" on page 9 for more information on signals.
- Clears the __argv. and __environment. stems. For this reason, it is not recommended that you use `syscalls('ON')` in a z/OS shell environment.

For REXX programs run from TSO/E or MVS batch, you use the `syscalls()` function to control the SYSCALL host command environment. You control the SYSCALL environment by using:

- `syscalls('ON')` to establish the SYSCALL environment
- `syscalls('OFF')` to end the SYSCALL environment
- `syscalls('SIGON')` to establish the signal interface routine
- `syscalls('SIGOFF')` to delete the signal interface routine

**Rule:** The words ON, OFF, SIGON, and SIGOFF must be in uppercase letters.

### Establishing the SYSCALL environment

The `syscalls('ON')` function establishes the SYSCALL environment. It sets up the REXX predefined variables and blocks all signals. The function sets this return value:

- 0  Successful completion.
- 4  The signal process mask was not set.
- 7  The process was dubbed, but the SYSCALL environment was not established.

---

2  z/OS V1R11.0 Using REXXX and z/OS UNIX System Services
The process could not be dubbed.

The following example shows how you can use the sycalls('ON') function at the beginning of a REXX program:

```rexx
if sycalls('ON')>3 then
   do
      say 'Unable to establish the SYSCALL environment'
      return
   end
```

**Ending the SYSCALL environment**

The sycalls('OFF') function ends the connection between the current task and z/OS UNIX. The task is undubbed and the REXX program continues.

In general, it is not necessary to make a sycalls('OFF') call.

The sycalls('OFF') function has one return value:

0  Successful completion.

**Establishing and deleting the signal interface routine**

The sycalls('SIGON') function establishes the signal interface routine (SIR). After you establish the SIR, use the sigaction syscall command to catch the signals you want to process and the sigprocmask syscall command to unblock those signals.

**Note:** For a REXX program run from a z/OS shell or from a program, the SIR is established by default.

The sycalls('SIGON') function has these return values:

0  Successful completion.
4  The SIR could not be established. The usual cause for this is that another SIR has already been established for the process.

If you are writing a REXX program that runs a program that requires a signal interface routine (for example, a program that uses the C runtime library), you must delete the SIR. The sycalls('SIGOFF') function deletes the SIR and uses sigprocmask() to reset the signal process mask so that it blocks all signals that can be blocked.

The sycalls('SIGOFF') function has two return values:

0  Successful completion.
4  The SIR could not be deleted. The usual cause for this is that a SIR did not exist for the process.

**The SH environment**

The SH environment is the default host command environment when a REXX program is run from a z/OS shell or from a program using exec(); it is available to a REXX program only in those two situations. In the SH environment, a syscall command runs as a z/OS shell command that has been issued this way:

```
/bin/sh -c shell_command
```

If you are running the REXX program from a z/OS shell or from a program, the SYSCALL environment is automatically initialized.
See Chapter 4 for some sample REXX programs that show how REXX and z/OS shell commands can work together—for example, a REXX program that can read output from a z/OS shell command. The `mount` and `unmount` sample programs in that chapter are shipped in the `/samples` directory as files `mountx` and `unmountx`.

Running a REXX program from the z/OS shells or from a program
You can run a REXX program from the z/OS shells, or you can call it from any program just as you would call an executable program. The REXX program runs as a separate process; it does not run in a TSO/E address space. You cannot use TSO/E commands in the REXX program.

A REXX program that is invoked from a z/OS shell or from a program must be a text file or a compiled REXX program that resides in the z/OS UNIX file system. It must have read and execute access permissions. Each line in the text file must be terminated by a newline character and must not exceed 2048 characters. Lines are passed to the REXX interpreter as they are. Sequence numbers are not supported; if you are using the ISPF editor to create the REXX program, be sure to set NUMBER OFF.

If you are working in a z/OS shell environment and use only a filename to invoke the REXX program, the PATH environment variable is used to locate it. For example, `myrexx` uses PATH to locate the program, but `.myrexx` searches only the working directory.

For a REXX program that is run from a z/OS shell or from a program, the SIR is established by default. If the REXX program calls a C program that is running POSIX(ON) or a program that requires an SIR, use the `syscalls('SIGOFF')` function to delete the SIR before calling that program.

CEXEC output from the REXX compiler is supported in the z/OS shell environments. To compile and put CEXEC output into the z/OS UNIX file system, you can use the REXXOEC cataloged procedure; it compiles under TSO/E and then uses the TSO/E OCOPY command to copy the compiled program from a data set to a file in the file hierarchy.

Using external functions and subroutines
You can call external functions and subroutines from a REXX program that resides in the z/OS UNIX file system. The search path for an external routine is similar to that used for a REXX program that is invoked from a z/OS shell or from a program. If only the filename is used on the call to the function or subroutine, the PATH environment variable is used to locate it; otherwise, the function name determines the search. For an executable module, the link pack area (LPA), link list, and STEPLIB may also be searched. The default z/OS environment searches for executable modules first. See "Customizing the z/OS UNIX REXX environment" on page 12.

The search order for modules and execs that are invoked as functions or subroutines is controlled by the `FUNCOSFL` flag in the REXX parameter module. For a description of that flag, see `/OS TSO/E REXX Reference`.

The following rule must be observed in naming and calling an external function or subroutine:
- If the name contains special characters or lowercase characters, you must enclose it in quotes—for example:
  ```
  ans='myfunc'(p1,p2)
  ```
If the name is not quoted, REXX folds the name to uppercase. The function call then fails, because the file is not found.

Executable external functions or subroutines that are written in a language other than interpreted REXX and located in the z/OS UNIX file system are not supported.

The TSO command environment

The TSO command environment (ADDRESS TSO) can be used from a z/OS UNIX REXX environment, and is initialized with:

```
address tso [command]
```

where `command` may be any TSO/E command, CLIST, or REXX exec that can run in a batch TSO TMP.

Commands addressed to TSO are run in a TSO TMP that is running in a separate address space and process from your REXX program. This provides you with the capability to run TSO commands. It does not provide you with the capability to use TSO commands to affect your REXX environment, or to have REXX statements or other host command environments affect your TSO process.

The TSO process is started when the first TSO/E command is run, and persists until your REXX program terminates or you run the TSO LOGOFF command. You can use the `ps` shell command to observe this process as the program `bpxwrtso`. Unexpected termination of the TSO process causes the next TSO command to fail with return code 16. A subsequent command starts a new TSO process.

Command input

Most native TSO commands, including commands that prompt for missing arguments, use TGET for input. This results in a command error, and the command usually terminates.

For commands that are able to read input, the source of the input is first any data that is currently on your stack, and then any data in your REXX program's standard input stream. Regardless of whether the command processes input, all data on the stack is queued to the TSO command. The stack is empty after any TSO command has been run.

The standard input stream may also be queued as input to the TSO command. For example, if you have a file redirected as input and you run a TSO command before processing that file, some or all of the file may be queued to the TSO command. If input is the terminal, queued input may be queued to the TSO command. This characteristic can be used to interact with some TSO commands.

You can disable command input by using the `rexxopt()` function with NOTSOIN specified.

Command output

By default, all command output is directed to your REXX process's standard output stream. You can use the `outtrap()` function to trap command output in variables.
Return codes

The special REXX variable RC usually contains the return code from the TSO command. If the command abends or is not found, or if another error is detected, special return codes are set, and a descriptive message may be written to the standard error stream:

- `-3` usually means that the TSO command was not found.
- Other negative numbers are usually abend codes. These should be accompanied by a message containing an abend reason code.
- `16` usually means that a processing error was encountered.

Examples

To run the TSO/E TIME command:

```
address tso 'time'
```

To trap command output and print it:

```
call outtrap out.
  address tso 'listc'
  do i=1 to out.0
    say out.i
  end
```

To run a REXX exec in TSO/E:

```
address tso
  "alloc fi(sysexec) da('schoen.rexx') shr"
  "myexec"

This is a functional replacement for the tsocmd utility:

```
/* rexx */
  address tso arg(1)
  return rc
```

The tsocmd utility is available for download from the Tools and Toys z/OS UNIX Web page at [http://www.ibm.com/servers/eserver/zseries/zos/unix/bpxa1toy.html](http://www.ibm.com/servers/eserver/zseries/zos/unix/bpxa1toy.html)

Variable scope

When the REXX program is initialized and the SYSCALL environment is established, the predefined variables are set up. If you call an internal subroutine that uses the PROCEDURE instruction to protect existing variables by making them unknown to that subroutine (or function), the predefined variables also become unknown. If some of the predefined variables are needed, you can either list them on the PROCEDURE EXPOSE instruction or issue another syscalls('ON') to reestablish the predefined variables. The predefined variables are automatically set up for external functions and subroutines. For example:

```
subroutine: procedure
  junk = syscalls('ON')
  parse arg dir
  'readdir (dir) dir. stem.'
```

Writing setuid and setgid REXX programs

Setting the set-group-ID-on-execution (setgid) permission means that when a file is run, the calling process’s effective GID is set to the file’s owner GID; the process seems to be running under the GID of the file’s owner, rather than that of the actual invoker.
Setting the set-user-ID-on-execution (setuid) permission means that when a file is run, the calling process's effective UID is set to the file's owner UID; the process seems to be running under the UID of the file's owner, rather than that of the actual invoker.

Like any other setuid or setgid program, a REXX program should not allow the user of the program to get control in your environment. Some examples of instructions that can let a user obtain control are:

- Interactive trace.
- Calling external functions or subroutines: Using a relative pathname can let the user get control if the user sets the PATH variable. External functions and subroutines run under the UID and GID of the main program, regardless of their setuid and setgid mode bits.

---

**Input and output for z/OS UNIX processing**

When a REXX program runs, open file descriptors are inherited from the process that issued the `exec()`.

**Using standard input, output, and error (file descriptors 0, 1, and 2)**

For a REXX program that is run from a z/OS shell or from a program, file descriptors 0, 1, and 2 (conventionally, standard input, standard output, and standard error files) are typically inherited.

**Attention:** A read or write error on file descriptors 0, 1, or 2 results in a halt interruption if the read or write was from a PARSE EXTERNAL instruction, SAY instruction, or EXECIO.

If the REXX program issues a PARSE EXTERNAL instruction, either explicitly or implicitly (such as from a PARSE PULL instruction with an empty stack), it reads standard input for a single text record. The newline character is stripped from the record before it is returned to the REXX program. Standard input is assumed to be a text file, such as your terminal input.

If the REXX program issues a SAY instruction, the text is directed to standard output, and a newline character is appended to the end of the text. Messages issued by REXX, including error and trace messages, are similarly directed to standard output.

If PARSE EXTERNAL is used after standard input has reached the end of the file, null lines are returned. The end-of-file condition can be detected by EXECIO. For more information, see “Using EXECIO” on page 8.

**Using SYSCALL commands**

The SYSCALL host command environment gives you more direct control over input and output. You can use:

- `readfile` to read an entire text file. See “readfile” on page 145 for more information.
- `writefile` to write an entire text file. See “writefile” on page 201 for more information.
- `read` to read bytes from any kind of file. See “read” on page 142 for more information.
write to write bytes to any kind of file. See "write" on page 199 for more information.

Using EXECIO

EXECIO differs from readfile and writefile in that it operates on open files. To read or write a file in segments (for example, a line at a time), use this TSO/E REXX command:

```
address MVS "EXECIO"
```

The data can come from and go to the stack or a stem. You can also use it to read or write an entire file. As shown in the following diagram, z/OS UNIX supports all the TSO/E REXX operands except OPEN, DISKRU, and linenum.

**Note:** When using EXECIO to read a file, the maximum allowable length of a line in the file is 1024 characters, including the newline character. For information on reading blocks of data, see "read" on page 142.

For the ddname operand, you can use the following pseudo-ddnames for processing, when the REXX program is run from a shell or from a program:

- File descriptors 0 to 7
- STDIN, STDOUT, STDERR

For information on EXECIO, see [z/OS TSO/E REXX Reference](#).

Exit status from a REXX program

When a REXX program is run from a z/OS shell or from a program it can return a return code. If the program returns a value in the range 0–255, that value is returned. Otherwise, a value of 255 is returned. If a program is terminated, REXX returns a value of 255.

Tokens returned from the PARSE SOURCE instruction

The tokens that are returned from the PARSE SOURCE instruction depend on where the REXX program is run: from a z/OS shell, from a program, from TSO/E, or from batch.

Running from a z/OS shell or from a program

When a REXX program runs in a z/OS shell environment or is called from a program, the PARSE SOURCE instruction returns nine tokens, in this order:

1. The string TSO
2. The string COMMAND, FUNCTION, or SUBROUTINE, depending on whether the program was invoked as a host command, from a function call in an expression, or with the CALL instruction
3. The first 8 characters of the name of the REXX program
4. The string PATH
5. The first 44 characters of the pathname of the REXX program
6. ? (question mark)
7. The name of the initial host command environment in uppercase: SH
8. The name of the address space in uppercase: OMVS
9. An 8-character user token: OpenMVS

To determine whether the REXX program was run from a z/OS shell, use token 8 or 9.

**Example**

If `myexec` is invoked from the z/OS shells and resides in the working directory, and if PATH is set to `.:/bin`, a PARSE SOURCE instruction returns the following tokens:

```
TSO COMMAND myexec PATH ./myexec ? SH OMVS OpenMVS
```

**Running from TSO/E or batch**

If the REXX program runs from TSO/E or MVS batch, the PARSE SOURCE instruction returns the tokens that are described in [z/OS TSO/E REXX Reference](#).

**Using the REXX signal services**

The REXX signal services consist of the following syscall commands:

- alarm
- kill
- pause
- sigaction
- sigpending
- sigprocmask
- sigsuspend
- sleep

REXX does not include a service that allows you to attach your own signal catcher. Instead, you have the following options:

- To use the REXX signal catcher as the action for a signal, you can specify the SIG_CAT variable as the signal handler on `sigaction`. SIG_CAT can terminate various wait conditions without causing the process to end. If a signal arrives when the process is not currently waiting and the signal is not blocked, it may be lost.
  
  There are two primary uses for SIG_CAT: when you are using the `alarm` command, and when you want to avoid unexpected process termination for other unblocked signals.
  
  SIG_CAT causes a signal to interrupt conditions such as waits and blocks, but the application cannot determine which signal was delivered. It is not a traditional signal catcher, as implemented in the C language.

- To set the action to the default action, you can specify SIG_DFL as the signal handler on `sigaction`.

- To set the action to ignore the signal, you can specify SIG_IGN as the signal handler on `sigaction`. 
POSIX.1 defines several C functions to manipulate signal sets. REXX does not define these functions; however, you can define each function using a single REXX statement, as shown in Table 2.

### Table 2. REXX statements for defining signal sets

<table>
<thead>
<tr>
<th>C Function</th>
<th>Equivalent REXX Statement</th>
</tr>
</thead>
<tbody>
<tr>
<td>sigsetempty()</td>
<td>sigsetempty: return copies(0,64)</td>
</tr>
<tr>
<td></td>
<td>• Parameters: none</td>
</tr>
<tr>
<td></td>
<td>• Returns: signal set</td>
</tr>
<tr>
<td>sigfillset()</td>
<td>sigfillset: return copies(1,64)</td>
</tr>
<tr>
<td></td>
<td>• Parameters: none</td>
</tr>
<tr>
<td></td>
<td>• Returns: signal set</td>
</tr>
<tr>
<td>sigaddset()</td>
<td>sigaddset: return overlay(1, arg(1), arg(2))</td>
</tr>
<tr>
<td></td>
<td>• Parameters: signal set, signal number</td>
</tr>
<tr>
<td></td>
<td>• Returns: signal set</td>
</tr>
<tr>
<td>sigdelset()</td>
<td>sigdelset: return overlay(0, arg(1), arg(2))</td>
</tr>
<tr>
<td></td>
<td>• Parameters: signal set, signal number</td>
</tr>
<tr>
<td></td>
<td>• Returns: signal set</td>
</tr>
<tr>
<td>sigismember()</td>
<td>sigismember: return substr(arg(1), arg(2), 1)</td>
</tr>
<tr>
<td></td>
<td>• Parameters: signal set, signal number</td>
</tr>
<tr>
<td></td>
<td>• Returns: 0 (not member) or 1 (is member)</td>
</tr>
</tbody>
</table>

### Using immediate commands

Immediate commands are TSO/E REXX commands, provided with the TSO/E implementation of the language. Immediate commands change characteristics that control the execution of an exec or program.

In response to an interrupt signal, usually <Ctrl-C>, the REXX interrupt handler suspends execution of the REXX program and prompts for an immediate command. The command is specified by number. In the z/OS UNIX REXX environment, the following commands are supported:

<table>
<thead>
<tr>
<th>Command</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>Continue execution</td>
</tr>
<tr>
<td>2</td>
<td>Issue a Halt Interruption</td>
</tr>
<tr>
<td>3</td>
<td>Start trace</td>
</tr>
<tr>
<td>4</td>
<td>End trace</td>
</tr>
<tr>
<td>5</td>
<td>Halt type</td>
</tr>
<tr>
<td>6</td>
<td>Resume type</td>
</tr>
</tbody>
</table>

You can use the `rexxopt()` function to disable this capability or attach this signal handler to other signals. As with any signal handler, the kernel may defer delivery of the signal depending on what the program is executing at the time.

**Note:** REXX programs that are run as setuid or setgid programs cannot be interrupted to issue an immediate command.

### Moving a REXX program from TSO/E to a z/OS shell

If you write a REXX program to run in TSO/E, it is likely that you will have to alter the REXX program to run it in a z/OS shell environment. Some of the differences between the two environments that you need to consider are:
You can use the `spawn` syscall command to run z/OS shell commands from the TSO/E environment.

Using the `syscalls('ON')` function at the beginning of the REXX program is required in TSO/E, but not in a z/OS shell environment. If you use `syscalls('ON')` in a z/OS shell environment, it clears the `__argv` and `__environment` stems. For this reason, it is not recommended that you use `syscalls('ON')` in a z/OS shell environment. Using `syscalls('ON')` in a z/OS shell environment also sets up the REXX predefined variables and blocks all signals. On entry to a REXX program in a z/OS shell environment, the REXX predefined variables are already set.

In TSO/E, the `syscalls('OFF')` function ends the process, but the REXX program continues to run. In the z/OS shells, the `syscalls('OFF')` function undubs the task and the REXX program continues.

`PARSE SOURCE` returns different tokens in TSO/E and in a z/OS shell environment. A REXX program uses the tokens to determine how it was run.

In TSO/E, the variables `__argv.0` and `__environment.0` are set to zero (0).

See [The TSO command environment](#) on page 5 for information about running TSO/E commands from a REXX program.

### Using `argv` and `environment` variables

Environment variables are text strings in the form `VNAME=value`, where `VNAME` is the name of the variable and `value` is its value. The stem variables `__argv` and `__environment` are always set to the original values passed to the first-level REXX program, and they are visible to external REXX functions. You may want to use `PARSE ARG` instead of the `__argv` stem in external REXX programs. As the following two sample programs show, using the `__argv` stem from an external exec returns the same data as it did from the initial exec. In order for an external REXX program to get the arguments a caller is sending it, it must use `arg()` or `PARSE ARG`:

**PGM1:**
```rexx
/* rexx */
say 'this is the main pgm'
say 'it was passed' __argv.0 'arguments:'
do i = 1 to __argv.0
   say ' Argument' i': '"__argv.i'"'
end

call 'pgm2' 'arguments', 'to pgm2'
```

**PGM2:**
```rexx
/* rexx */
say 'This is pgm2'
say 'Using __argv stem, there are' __argv.0 'arguments. They are:'
do i = 1 to __argv.0
   say ' Argument' i': '"__argv.i'"'
end

say 'Using arg(), there are' arg() 'arguments:'
do i = 1 to arg()
   say ' Argument' i': '"arg(i)'"'
end
```

**Sample execution**

$ pgm1 'arguments to' 'pgm1'
this is the main pgm
it was passed 3 arguments:
   Argument 1: "pgm1"
Argument 2: "arguments to"
Argument 3: "pgm1"
This is pgm2
Using __argv stem, there are 3 arguments. They are:
Argument 1: "pgm1"
Argument 2: "arguments to"
Argument 3: "pgm1"
Using arg(), there are 2 arguments:
Argument 1: "arguments"
Argument 2: "to pgm2"

Customizing the z/OS UNIX REXX environment

When a REXX program is run from the z/OS shells or called from a program using exec(), the z/OS UNIX REXX environment that is established is created from the module BPXWRXEV. The source for this module is member BPXWRX01 in SYS1.SAMPLIB.

This environment is inherited from the default MVS REXX environment. However, the default handling of error messages from the REXX processor is overridden so that the messages are written to STDOUT. This is the same place to which output from the SAY instruction and trace information is sent.

You can further customize the sample member to alter the REXX environment for REXX programs running under z/OS UNIX without affecting REXX programs running in the z/OS environment. For detailed information on how to change the default values for initializing an environment, see z/OS TSO/E REXX Reference.

Performance in the SYSCALL environment

syscalls('ON') ensures that the SYSCALL host command environment is available in your REXX environment. If the call detects that SYSCALL is not available in your environment, it dynamically adds it.

Performance characteristics for dynamically added host commands are not as good as for host commands that are included in the initial environment: Every time a command is directed to the SYSCALL host command environment, the TSO/E REXX support loads the module for the SYSCALL host command.

To avoid this, include the SYSCALL host command in the three default TSO/E environments:

<table>
<thead>
<tr>
<th>Module name</th>
<th>SYS1.SAMPLIB member name</th>
<th>REXX environment</th>
</tr>
</thead>
<tbody>
<tr>
<td>IRXPARMS</td>
<td>IRXREXX1</td>
<td>MVS</td>
</tr>
<tr>
<td>IRXTSPRM</td>
<td>IRXREXX2</td>
<td>TSO</td>
</tr>
<tr>
<td>IRXISPRM</td>
<td>IRXREXX3</td>
<td>ISPF</td>
</tr>
</tbody>
</table>

Customizing IRXISPRM provides dramatic performance improvement for REXX programs that use syscall commands from TSO/E or MVS batch.

Make the following changes to the SYS1.SAMPLIB members to add the SYSCALL host command to that default environment:
1. Find the label SUBCOMTB_TOTAL and add 1 to its value. For example:
   Change SUBCOMTB_TOTAL DC F'14' to SUBCOMTB_TOTAL DC F'15'.
2. Find the label SUBCOMTB_USED and add 1 to its value. For example:
Change SUBCOMTB_USED DC F'14' to SUBCOMTB_USED DC F'15'.

3. Find the end of the subcommand table, just before the label PACKTB or PACKTB_HEADER, and add the following lines:

```
SUBCOMTB_NAME_REXXIX DC CL8'SYSCALL'
SUBCOMTB_ROUTINE_REXXIX DC CL8'BPXWREXX'
SUBCOMTB_TOKEN_REXXIX DC CL16'
```

4. Assemble and link-edit the module and replace the default TSO/E module. These are normally installed in SYS1.LPALIB.

See [z/OS TSO/E REXX Reference](#) for additional information on customizing the default environments.

---

**Authorization**

Users authorized to perform special functions are defined as having appropriate privileges, and they are called superusers. Appropriate privileges also belong to users with:

- A user ID of zero
- RACF-supported user privileges trusted and privileged, regardless of their user ID

A user can switch to superuser authority (with an effective UID of 0) if the user is permitted to the BPX.SUPERUSER FACILITY class profile within RACF®. Either the ISPF Shell or the `su` shell command can be used to switch to superuser authority.

---

**Security**

This document assumes that your operating system contains Resource Access Control Facility (RACF). You could use an equivalent security product updated to handle z/OS UNIX security.
Chapter 2. z/OS UNIX REXX programming services

An application that supports scripting or macro languages, such as an editor, can use REXX as the macro language. An application that is written in a programming language such as C can create a z/OS UNIX REXX environment and run a REXX program directly. For information on using the TSO/E REXX programming services, such as IRXJCL and IRXEXEC, see z/OS TSO/E REXX Reference.

Creating a z/OS UNIX REXX environment from an application

To create a z/OS UNIX REXX environment, fetch and call the module BPXWRBLD from a key 8 problem state program. z/OS linkage for C (that is, standard OS linkage) is required.

The BPXWRBLD module requires the following parameters:

**16K area**
A 16,000-byte environment area. This must persist for the life of the REXX environment.

**arg count**
The count of the number of REXX initialization arguments.

**arg pointer array**
An array of pointers to null-terminated strings, one for each REXX initialization argument. The array and the null-terminated strings must persist for the life of the REXX environment.

**env count**
The count of the number of environment variables to be exported to the REXX program.

**env length pointer array**
An array of pointers to fullwords, one for each environment variable. The fullword contains the length of the string that defines the environment variable, including the terminating null. The last element of the array must point to a fullword of 0.

The array and the fullwords must persist for the life of the REXX environment.

**env pointer array**
An array of pointers to null-terminated strings, one for each environment variable. Each string defines one environment variable. The array and the null-terminated strings must persist for the life of the REXX environment.

The format of the string is

```
NAME=value
```

where NAME is the environment variable name, and value is the value for the environment variable followed by a null character.

**REXX env addr**
The address of a fullword where the address of the newly created REXX environment is returned.

If BPXWRBLD fails to create the environment, it returns the return code it received from the IRXINIT service. BPXWRBLD does not return any other codes.
The parameter list is a standard MVS variable-length parameter list. On entry, the following registers must be set:

Register 1  Address of the parameter list
Register 13  Address of a register save area
Register 14  Return address
Register 15  Entry point address

Register 1 contains the address of the parameter list:

```
R1 --> -----------
    |  ---------+ 16K area
    |  ---------+ arg count
    |  ---------+ arg pointer array
    |  ---------+ env count
    |  ---------+ env length pointer array
    |  ---------+ env pointer array
    |  1  ---------+ Rexx env addr
    -----------
```

When constructing arguments to the REXX program that are also passed to BPXWRBLD, keep in mind that:

- The only use of the argument count and argument array is to populate the `__argv` REXX variables. You can set the argument count to 0 if the REXX programs will always get their arguments using PARSE ARG or the `arg(1)` REXX function call. In this case, `__argv:0` is set to 0 when the REXX program is run.
- After the call to BPXWRBLD, do not alter the data that is pointed to by the environment pointer arrays or the arg pointer array.

Signals are not supported in this environment.

**Running the REXX program**

Before calling a TSO/E REXX service to run the program, ensure that file descriptors 0, 1, and 2 are open. The REXX program will fail if it attempts a PARSE EXTERNAL, EXECIO, or SAY and that function fails.

After the REXX environment is established, the program can call either the IRXJCL or the IRXEXEC TSO/E REXX service to run the REXX program.

- If the IRXJCL service is used, the name of the REXX program is the first word of the IRXJCL parameter string. It is limited to 8 characters.
- If you request the IRXEXEC service to load the program, you must provide the name of the REXX program in the member field of the EXECBLK. Set the DDNAME field to spaces. This also limits the name of the REXX program to 8 characters. Names longer than 8 characters can be supported with additional programming effort. You would need to preload the program and build an INSTBLK instead of an EXECBLK for the IRXEXEC call. If the REXX program is compiled in CEXEC format, load it as a single-record program.
- If the name of the REXX program does not contain a slash (/), the PATH environment variable is used to locate the program.
The current REXX environment must be the z/OS UNIX REXX environment. You cannot pass the environment to be used in Register 0.

When the REXX program is being loaded, the IRXEXEC or the IRXJCL service uses one file descriptor to open the file, read it, and close it. If no file descriptor is available because the maximum number of file descriptors are already open, the program cannot be loaded.

**Example: C/370 program**

This C/370™ program creates a REXX environment and runs a REXX program:

```c
#include <stdlib.h>
#include <string.h>
#include <stdio.h>

typedef int EXT();
#pragma linkage(EXT,OS)

int main(int argc, char *argv) {
    extern char **environ; /* access environ variables */
    EXT *irxjcl; /* pointer to IRXJCL routine */
    EXT *bpwrlbd; /* pointer to BPXWRBLD routine */
    char *penvb; /* addr of REXX environment */
    int i, j; /* temps */
    long rci; /* return code */
    int **environlp; /* ptr to env length pointers */
    int *environl; /* ptr to env lengths */
    char rxwork[16000]; /* OE MVS env work area */
    char *execname="execname"; /* name of exec up to 8 chars */
    char *execparm="exec parameter string"; /* parm to exec */
    struct s_rxparm {
        short len; /* halfword length of parm */
        char name[8]; /* area to hold exec name */
        char space; /* one space */
        char text[253]; /* big area for exec parm */
    } *rxparm;

    /* if stdin or stdout are not open you might want to open file */
    /* descriptors 0 and 1 here */

    /* if no environ, probably tso or batch - make one */
    if (environ==NULL) {
        environ=(char **)malloc(8); /* create one */
        environ[0]="PATH=."; /* set PATH to cwd */
        environ[1]=NULL; /* env terminator */
    }

    /* need to build the environment in the same format as expected by */
    /* the exec() callable service. See */
    /* Assembler Callable Services for UNIX System Services. */

    /* the environ array must always end with a NULL element */
    for (i=0; environ[i]!=NULL; i++); /* count vars */
    environlp=(int **)malloc(i*4+4); /* get array for len ptrs */
    environl=(int *)malloc(i*4+4); /* get words for len vals */
    for (j=0; j<i; j++) {
        environlp[j]=&environl[j]; /* point to len */
        environl[j]=strlen(environ[j])+1; /* set len word */
    }
    environl[i]=0; /* null entry at end */

    /* load routines */
    irxjcl=(EXT*)fetch("IRXJCL ");
}
```

/* build the REXX environment */
rcinit=bpxwrbld(rxwork,
    argc,argv,
    &penvb);
if (rcinit!=0) {
    printf("environment create failed rc=%d\n",rcinit);
    return 255;
};

/* if you need to add subcommands or functions to the environment, */
/* or create a new environment inheriting the current one, this is */
/* the place to do it. The user field in the environment is used */
/* by the z/OS UNIX REXX support and must be preserved. */

/* run exec */
rxparm=(struct s_rxparm *)malloc(strlen(execname)+
    strlen(execparm)+
    sizeof(struct s_rxparm));
memset(rxparm->name, '\', sizeof(rxparm->name));
memcpy(rxparm->name, execname, strlen(execname));
rxparm->space=' ';
memcpy(rxparm->text, execparm, i=strlen(execparm));
rxparm->len=sizeof(rxparm->name)+sizeof(rxparm->space)+i;
return irxjcl(rxparm);
Chapter 3. The syscall commands

In most cases, syscall commands invoke the z/OS UNIX callable service that corresponds to the command verb (the first word of the command). The parameters that follow the command verb are specified in the same order as in POSIX.1 and the z/OS UNIX callable services, where applicable.

For complete information about the processing of a particular syscall command, read about the callable service it invokes, as described in z/OS UNIX System Services Programming: Assembler Callable Services Reference.

Specifying a syscall command

You must specify the syscall command parameters in the order indicated in the syscall command description.

syscall command name

The syscall command name is case-insensitive: you can specify it as uppercase, lowercase, or mixed case.

Parameters

You can specify several types of parameters, but most fall into the following categories:

pathname

The pathname is case-sensitive, and it is specified as a string. The syscall commands can take a relative or absolute pathname as a parameter. The search for a relative pathname begins in your working directory:

- If you are running a REXX program from a z/OS shell, your working directory is inherited from your z/OS shell session.
- If you are running a REXX program in TSO/E, your working directory is typically your home directory.

Portable pathnames can use only the characters in the POSIX portable filename character set:

- Uppercase or lowercase A to Z
- Numbers 0 to 9
- Period (.)
- Underscore (_)
- Hyphen (-)

Do not include any nulls in a pathname.

mode

The mode is a three- or four-digit number that corresponds to the access permission bits. Each digit must be in the range 0–7, and at least three digits must be specified. See Appendix B for more information on permissions.

stem

The name of a stem variable. A stem can be used for input, output, or both. A stem is indicated by a . (period) at the end of the variable name.

- The variable name for the first value consists of the name of the stem variable with a 1 appended to it. The number is incremented for each value—for example, vara.1, vara.2, and vara.3.
• The variable name that contains the number of variables returned (excluding 0) consists of the name of the stem variable with a 0 appended to it—for example, `vara.0`

If you omit the period from the end of the variable name, a numeric suffix is appended to the name (for example, `foo` would become `foo0`, `foo1`, and so on).

The name of a stem variable is case-insensitive.

`variable`

The name of a REXX variable. The name is case-insensitive.

**Specifying numerics**

All numbers are numeric REXX strings. Negative numbers can be preceded by a minus sign (−); others must be unsigned.

The SYSCALL environment supports a 10-digit field. If you are performing arithmetic on a field longer than 9 digits, you must set precision to 10. A range of up to $2^{31}-1$ is supported.

**Specifying strings**

You can specify a string in any of these ways:

<table>
<thead>
<tr>
<th>String</th>
<th>Example</th>
</tr>
</thead>
<tbody>
<tr>
<td>Any series of characters not containing a space. This example shows a path name with no space in it.</td>
<td>&quot;creat /u/wjs/file 700&quot;</td>
</tr>
<tr>
<td>Any series of characters delimited by ' and not containing '. This example shows a path name with a space in it.</td>
<td>&quot;creat 'u/wjs/my file' 700&quot;</td>
</tr>
<tr>
<td>Any series of characters delimited by &quot; and not containing &quot;. This example shows a path name with a space in it.</td>
<td>'creat &quot;u/wjs/my file&quot; 700'</td>
</tr>
<tr>
<td>A variable name enclosed in parentheses. Strings that contain both the single and double quote characters must be stored in a variable, and you must use the variable name.</td>
<td>file='&quot;u/wjs/my file'&quot; &quot;creat (file) 700&quot;</td>
</tr>
</tbody>
</table>

The following example uses a variable enclosed in parentheses to avoid problems with a blank in the filename:

```
file='"u/wjs/my file'" 
"creat (file) 700"
```

If you incorrectly coded the second line as:

```
"creat /u/wjs/my file 700"
```

it would contain four tokens instead of three.
Using predefined variables

The predefined variables that are available to a REXX program in a z/OS shell environment make symbolic references easier and more consistent—for example, when you are specifying a flag or using a stem variable. Instead of coding a numeric value, you can specify the predefined variable that is used to derive that numeric value.

Appendix A lists all the predefined variables alphabetically and shows their numeric value and data type. If a variable is a stem variable, this shows the data type for the stem variable.

You can also use the index of this document as a reference: under the name of each syscall command are grouped the names of the predefined variables associated with it.

Return values

A command can be issued to the SYSCALL environment or the SH environment, and the return values are different in the two environments.

Returned from the SYSCALL environment

When a syscall command completes, the environment can set four reserved variables:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>RC</td>
<td>A numeric return code from the command execution.</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Value Range</th>
<th>Meaning</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>The command finished successfully. If there is an error code for the requested function, it is returned in RETVAL and ERRNO.</td>
</tr>
<tr>
<td>&gt;0</td>
<td>The command finished successfully, but a function-specific warning is indicated.</td>
</tr>
<tr>
<td>−3</td>
<td>The command environment has not been called. Probably the syscall('ON') function did not end successfully, or the current address environment is not SYSCALL.</td>
</tr>
<tr>
<td>−20</td>
<td>The command was not recognized, or there was an improper number of parameters specified on the command.</td>
</tr>
<tr>
<td>−21,−22, ...</td>
<td>The first, second, ... parameter is in error. (The parameter is indicated by the second digit.)</td>
</tr>
<tr>
<td>&lt;0</td>
<td>Other negative values may be returned by the REXX language processor. A negative value means that the command did not finish successfully.</td>
</tr>
</tbody>
</table>

| RETVAL     | A numeric return value from the callable service. This indicates the success or failure of the service. For most successful calls to services, RETVAL is set to zero; for unsuccessful calls, this value is −1. |

However, there are some services (such as getgrgid and getgrnam) that return zero instead of −1 when the service fails. In addition, some services return a positive RETVAL to indicate success. For
details about a specific service, see [z/OS UNIX System Services Programming: Assembler Callable Services Reference].

**ERRNO**
A hexadecimal error number from the callable service. This variable is valid only if the return code (RC) is not negative and RETVAL is −1.

**ERRNOJR**
A hexadecimal reason code from the callable service. This variable is valid only if the return code (RC) is not negative and RETVAL is −1.

**Returned from the SH environment**
When a command completes in the SH environment, the return code is set in the variable RC. Unusual situations cause the return code to be set to a negative value:

-1xxx Terminated by signal xxx
-2xxx Stopped by signal xxx
-3xxx Fork failed with error number xxx
-4xxx Exec failed with error number xxx
-5xxx Wait failed with error number xxx

**Syscall command descriptions**
Function
access invokes the access callable service to determine if the caller can access a file.

Parameters

pathname
The pathname of the file to be checked for accessibility.

flags
One or more numeric values that indicate the accessibility to be tested. You can specify a numeric value (see Appendix A) or the predefined variable used to derive the appropriate numeric value. The predefined variables you can specify are:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>F_OK</td>
<td>Test for file existence</td>
</tr>
<tr>
<td>R_OK</td>
<td>Test for permission to read</td>
</tr>
<tr>
<td>W_OK</td>
<td>Test for permission to write</td>
</tr>
<tr>
<td>X_OK</td>
<td>Test for permission to execute</td>
</tr>
</tbody>
</table>

For example, R_OK+W_OK tests for read and write permission.

Usage notes
1. Testing for file permissions is based on the real user ID (UID) and real group ID (GID), not the effective UID or effective GID of the calling process.
2. The caller can test for the existence of a file or for access to the file, but not both.
3. In testing for permission, the caller can test for any combination of read, write, and execute permission. If the caller is testing a combination of permissions, a −1 is returned if any one of the accesses is not permitted.
4. If the caller has appropriate privileges, the access test is successful even if the permission bits are off, except when testing for execute permission. When the caller tests for execute permission, at least one of the execute permission bits must be on for the test to be successful.

Example
To test for permission to execute grep:
"access '/bin/grep'" x_ok
Function

`acldelete` deletes an access control list (ACL) associated with `pathname`.

Parameters

**pathname**

The pathname of the file or directory the ACL is associated with.

**acltype**

Indicates the type of ACL. This parameter can have the following values:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>ACL_TYPE_ACCESS (1)</td>
<td>An access ACL</td>
</tr>
<tr>
<td>ACL_TYPE_FILEDEFAULT (2)</td>
<td>A file default ACL</td>
</tr>
<tr>
<td>ACL_TYPE_DIRDEFAULT (3)</td>
<td>A directory default ACL</td>
</tr>
</tbody>
</table>

Usage notes

1. For regular files, the `acltype` must indicate it is an access ACL that is to be deleted.
2. For a directory, the `acltype` must indicate one of the three types of ACLs (access, file default, or directory default).

Example

To delete the access ACL from the `/tmp` directory, this example assumes the user has set the appropriate stem variable before the call:

```
"acldelete /tmp acl." acl_type_access
```

For a complete example that uses several of the ACL services to list ACLs, see "List the ACL entries for a file" on page 214.

For more information about access control lists, see Using access control lists (ACLs) in *z/OS UNIX System Services Planning*. 
Function

`acldeleteentry` deletes a specific entry in the access control list (ACL) represented by `variable`.

Parameters

`variable`

The name of a REXX variable that contains a token to access an ACL.

`stem`

The name of a stem variable that contains an ACL entry. STEM.0 contains a count of the number of variables set in the stem. The following variables may be used to access the stem variables. The number in parentheses is the actual value of the variable:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>ACL_ENTRY_TYPE (1)</td>
<td>Indicates the type of ACL entry:</td>
</tr>
<tr>
<td></td>
<td>ACL_ENTRY_USER (1) (User ACL)</td>
</tr>
<tr>
<td></td>
<td>ACL_ENTRY_GROUP (2) (Group ACL)</td>
</tr>
<tr>
<td>ACL_ID (2)</td>
<td>The numeric id, uid or gid of the entry</td>
</tr>
<tr>
<td>ACL_READ (3)</td>
<td>Indicates read access (1 = yes, 0 = no)</td>
</tr>
<tr>
<td>ACL_WRITE (4)</td>
<td>Indicates write access (1 = yes, 0 = no)</td>
</tr>
<tr>
<td>ACL_EXECUTE (5)</td>
<td>Indicates execute or search access (1 = yes, 0 = no)</td>
</tr>
<tr>
<td>ACL_DELETE (6)</td>
<td>Indicates that the ACL entry is deleted (1 = yes, 0 = no)</td>
</tr>
</tbody>
</table>

Usage notes

1. The entry to delete is identified by the entry type and ID, and is contained in `stem`.
2. If the entry does not exist, the service will return `retval=−1` and `errno=enoent`.

Example

To delete the ID in an ACL, this example assumes the user has set the appropriate stem variable before the call:

```
"acldeleteentry tokenvar acl." acl_id
```

For a complete example that uses several of the ACL services to list ACLs, see "List the ACL entries for a file" on page 214.

For more information about access control lists, see Using access control lists [ACLs] in z/OS UNIX System Services Planning.
aclfree

Function

`aclfree` releases resources associated with the access control list (ACL) represented by `variable` and obtained using the `aclinit` syscall command.

Parameters

- `variable`
  The name of a REXX variable that contains a token to access an ACL.

Example

```
aclfree tokenvar
```

For a complete example that uses several of the ACL services to list ACLs, see "List the ACL entries for a file" on page 214.

For more information about access control lists, see Using access control lists (ACLs) in z/OS UNIX System Services Planning.
Function

`aclget` reads an access control list (ACL) of the specified type associated with the file identified by `pathname`. The ACL is associated with the specified `variable`, and is accessed and altered using the `acldeleteentry`, `aclgetentry`, and `aclupdateentry` services.

Before using `aclget`, the variable must be initialized by using `aclinit`.

Parameters

`variable`

The name of a REXX variable that contains a token to access an ACL.

`pathname`

The pathname of the file or directory the ACL is associated with.

`acltype`

Indicates the type of ACL. This parameter can have the following values:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>ACL_TYPE_ACCESS (1)</td>
<td>An access ACL</td>
</tr>
<tr>
<td>ACL_TYPE_FILEDEFAULT (2)</td>
<td>A file default ACL</td>
</tr>
<tr>
<td>ACL_TYPE_DIRDEFAULT (3)</td>
<td>A directory default ACL</td>
</tr>
</tbody>
</table>

Example

```
"aclget tokenvar /lpp/payroll" acl_type_access
```

For a complete example that uses several of the ACL services to list ACLs, see "List the ACL entries for a file" on page 214.

For more information about access control lists, see Using access control lists (ACLs) in z/OS UNIX System Services Planning.
aclgetentry

Function

dclgetentry reads an access control list (ACL) entry from the ACL represented by variable.

Parameters

variable

The name of a REXX variable that contains a token to access an ACL.

stem

The name of a stem variable that contains an ACL entry. STEM.0 contains a count of the number of variables set in the stem. The following variables may be used to access the stem variables. The number in parentheses is the actual value of the variable:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>ACL_ENTRY_TYPE (1)</td>
<td>Indicates the type of ACL entry:</td>
</tr>
<tr>
<td></td>
<td>ACL_ENTRY_USER (1) (User ACL)</td>
</tr>
<tr>
<td></td>
<td>ACL_ENTRY_GROUP (2) (Group ACL)</td>
</tr>
<tr>
<td>ACL_ID (2)</td>
<td>The numeric id, uid or gid of the entry</td>
</tr>
<tr>
<td>ACL_READ (3)</td>
<td>Indicates read access (1 = yes, 0 = no)</td>
</tr>
<tr>
<td>ACL_WRITE (4)</td>
<td>Indicates write access (1 = yes, 0 = no)</td>
</tr>
<tr>
<td>ACL_EXECUTE (5)</td>
<td>Indicates execute or search access (1 = yes, 0 = no)</td>
</tr>
<tr>
<td>ACL_DELETE (6)</td>
<td>Indicates that the ACL entry is deleted (1 = yes, 0 = no)</td>
</tr>
</tbody>
</table>

index

Specifies the relative ACL entry to access. The first entry is 1.

Usage notes

1. An entry is identified by index, if index is specified. Otherwise, the entry is identified by the type and ID specified in stem.
2. If the entry does not exist, the service will return retval=−1 and errno=enoent.

Example

To read an ACL entry, this example assumes the user has set the appropriate stem variable before the call:

"aclgetentry tokenvar acl."

For a complete example that uses several of the ACL services to list ACLs, see "List the ACL entries for a file" on page 214.

For more information about access control lists, see [Using access control lists (ACLs)](z/OS UNIX System Services Planning)
aclinit

Function
aclinit obtains resources necessary to process access control lists (ACLs) and associates those resources with variable.

Parameters
variable
   The name of a REXX variable that contains a token to access an ACL.

Usage notes
1. The variable associated with the obtained resources must be passed to other services that operate on an ACL.
2. Any one variable can only represent one ACL at a time.
3. aclfree must be used to release the resources obtained by aclinit.

Example
"aclinit tokenvar"

For a complete example that uses several of the ACL services to list ACLs, see "List the ACL entries for a file" on page 214.

For more information about access control lists, see Using access control lists (ACLs) in z/OS UNIX System Services Planning.
aclset

Function

`aclset` replaces the access control list (ACL) associated with `pathname` with the ACL represented by `variable`.

Parameters

`variable`

The name of a REXX variable that contains a token to access an ACL.

`pathname`

The pathname of the file or directory the ACL is associated with.

`acltype`

Indicates the type of ACL. This parameter can have the following values:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>ACL_TYPE_ACCESS (1)</td>
<td>An access ACL</td>
</tr>
<tr>
<td>ACL_TYPE_FILEDEFAULT (2)</td>
<td>A file default ACL</td>
</tr>
<tr>
<td>ACL_TYPE_DIRDEFAULT (3)</td>
<td>A directory default ACL</td>
</tr>
</tbody>
</table>

Example

This example replaces the directory default ACL:

```
"aclset tokenvar /u/dept58" acl_type_dirdefault
```

For a complete example that uses several of the ACL services to list ACLs, see "List the ACL entries for a file" on page 214.

For more information about access control lists, see Using access control lists (ACLs) in z/OS UNIX System Services Planning.
Function

`aclupdateentry` updates an existing access control list (ACL) entry or creates a new entry if the entry does not already exist in the ACL contained in `variable`.

Parameters

**variable**

The name of a REXX variable that contains a token to access an ACL.

**stem**

The name of a stem variable that contains an ACL entry. STEM.0 contains a count of the number of variables set in the stem. The following variables may be used to access the stem variables. The number in parentheses is the actual value of the variable:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>ACL_ENTRY_TYPE</td>
<td>Indicates the type of ACL entry:</td>
</tr>
<tr>
<td>ACL_ENTRY_USER</td>
<td>ACL_ENTRY_USER (1) (User ACL)</td>
</tr>
<tr>
<td>ACL_ENTRY_GROUP</td>
<td>ACL_ENTRY_GROUP (2) (Group ACL)</td>
</tr>
<tr>
<td>ACL_ID</td>
<td>The numeric id, uid or gid of the entry</td>
</tr>
<tr>
<td>ACL_READ</td>
<td>Indicates read access (1 = yes, 0 = no)</td>
</tr>
<tr>
<td>ACL_WRITE</td>
<td>Indicates write access (1 = yes, 0 = no)</td>
</tr>
<tr>
<td>ACL_EXECUTE</td>
<td>Indicates execute or search access (1 = yes, 0 = no)</td>
</tr>
<tr>
<td>ACL_DELETE</td>
<td>Indicates that the ACL entry is deleted (1 = yes, 0 = no)</td>
</tr>
</tbody>
</table>

**index**

Specifies the relative ACL entry to access. The first entry is 1.

Usage notes

1. An entry is identified by either the entry type and ID contained in `stem` or by the relative entry number if `index` is specified. If `index` is 0 or greater than the current number of ACL entries, a new entry is created.

2. `aclupdateentry` can update a deleted ACL entry, mark an entry as deleted, mark the entry as not deleted, or add a new entry that is also marked as deleted by appropriate setting of `stem.acl_delete`. Also, duplicate entries can be added when using `index`. This may result in an unexpected ACL and should be avoided.

3. The read, write, execute, and deleted attributes are set based on the corresponding stem variables. A value of 1 indicates the attribute is to be set to 1. Any other value, including not setting the variable, results in the attribute being set to 0.

4. If the requested entry cannot be located, a new entry is created and the relative index for that entry is returned in `retval`. If an existing entry is updated, `retval` will contain 0. If `retval=−1` and `errno=enoent`, then a new entry could not be created because it would exceed the maximum number of entries (1024).
aclupdateentry

Example
To update an ACL entry, this example assumes the user has set the appropriate
stem variable before the call.
"aclupdateentry tokenvar acl."

For a complete example that uses several of the ACL services to list ACLs, see
"List the ACL entries for a file" on page 214.

For more information about access control lists, see Using access control lists
[ACLs] in z/OS UNIX System Services Planning
Function
alarm invokes the alarm callable service to generate a SIGALRM signal after the number of seconds specified have elapsed.

Parameters
seconds
The number of seconds to pass between receipt of this request and generation of the SIGALRM signal.

Usage notes
1. The default action for an alarm signal is to end a process.
2. The alarm callable service is always successful, and no return value is reserved to indicate an error.
3. An abend is generated when failures are encountered that prevent the alarm callable service from completing successfully.
4. Alarm requests are not stacked; only one SIGALRM can be scheduled to be generated at a time. If the previous alarm time did not expire and a new alarm is scheduled, the most recent alarm reschedules the time that SIGALRM is generated.
5. See "Using the REXX signal services" on page 9 for additional information on using signals.

Example
To generate a SIGALRM after 10 seconds:
"alarm 10"
**catclose**

**Function**
`catclose` closes a message catalog that was opened by `catopen`.

**Parameters**
- `catalog_descriptor`
  - The catalog descriptor (a number) returned by `catopen` when the message catalog was opened.

**Usage notes**
If it is unsuccessful, `catclose` returns −1 and sets ERRNO to indicate the error.

**Example**
See the example for "catgets" on page 35.
Function
The catalog descriptor (a number) returned by `catopen` when a message catalog was opened earlier.

set_number
A number that identifies a message set in the message catalog.

message_number
A number that identifies a message in a message set in the message catalog.

variable
The name of the buffer in which the message string is returned.

Usage notes
1. Set `variable` to a default message text prior to invoking the `catgets` command. If the message identified by `message_number` is not found, `variable` is not altered and can be used after the command has been invoked.
2. If the command is unsuccessful, `variable` is returned and ERRNO may be set to indicate the error.

Example
```c
"catopen mymsgs.cat"
cd=retval
:
msg='error processing request'
"catgets (cd) 1 3 msg"
say msg
:
"catclose" cd
```
Function

`catopen` opens a message catalog that has been built by the `gencat` utility. (For more information about `gencat`, see the `gencat` command description in `z/OS UNIX System Services Command Reference`.) The catalog descriptor is returned in `RETVAL`.

Parameters

catalog name
The pathname for the message catalog. If the pathname contains a slash (/), the environment variables `NLSPATH` and `LANG` do not affect the resolution of the pathname.

Usage notes
1. The catalog descriptor returned in `RETVAL` can be used with the `catgets` and `catclose` commands. Do not use the catalog descriptor with any other commands.
2. If it is unsuccessful, `catopen` returns a −1 and sets ERRNO to indicate the error.

Example

See the example for "catgets" on page 35.
Function

`cert` calls BPX1SEC to register or deregister a certificate with the calling user.

Parameters

`cert_variable`
- The name of the variable that contains the certificate.

`cert_length`
- The length of the certificate

`flag`
- Valid values for the flag are:
  - 0 — register
  - 1 — deregister

Usage notes

1. The intent of the `cert` service is to provide a way for the caller to associate/disassociate a certificate with the calling user. No new security environment is created and no authentication of the user is conducted.

2. The caller needs access to the RACDCERT facility class (as defined in the initACEE documentation) to register/deregister a certificate. No other authority above that is required to use `cert`.

3. The certificate is a data area that includes a 4-byte length field, header information for some certificate types, the actual certificate, and trailer information for some certificate types. The length passed in on the `cert` syscall is the whole length of that data area. It is up to the caller to build the appropriate structure.

Example

```
'cert newcert' length(newcert) 1
```
chattr

Function
chattr invokes the chattr callable service to set the attributes associated with a file. You can change the file mode, owner, access time, modification time, change time, reference time, audit flags, general attribute flags, and file size.

Parameters
pathname
The pathname of the file.

attribute_list
A list of attributes to be set and their values. The attributes are expressed either as numeric values (see [Appendix A]), or as the predefined variables beginning with ST_, followed by arguments for that attribute. The attributes that may be changed and their parameters are:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>ST_CCSID</td>
<td>Coded character set ID; first 4 characters are the file tag.</td>
</tr>
<tr>
<td>ST_MODE</td>
<td>1 argument: permission bits as 3 octal digits.</td>
</tr>
<tr>
<td>ST_UID</td>
<td>2 arguments: UID and GID numbers.</td>
</tr>
<tr>
<td>ST_SIZE</td>
<td>1 argument: new file size.</td>
</tr>
<tr>
<td>ST_ATIME</td>
<td>1 argument for access time: new time or −1 for TOD.</td>
</tr>
<tr>
<td>ST_MTIME</td>
<td>1 argument for modification time: new time or −1 for TOD.</td>
</tr>
<tr>
<td>ST_CTIME</td>
<td>1 argument for change time: new time or −1 for TOD.</td>
</tr>
<tr>
<td>ST_SETUID</td>
<td>No arguments.</td>
</tr>
<tr>
<td>ST_SETGID</td>
<td>No arguments.</td>
</tr>
<tr>
<td>ST_AAUDIT</td>
<td>1 argument: new auditor audit value.</td>
</tr>
<tr>
<td>ST_UAUDIT</td>
<td>1 argument: new user audit value.</td>
</tr>
<tr>
<td>ST_STICKY</td>
<td>No arguments.</td>
</tr>
<tr>
<td>ST_GENVALUE</td>
<td>2 arguments: names of two variables. The first variable contains the general attribute mask and the second contains the general attribute value.</td>
</tr>
<tr>
<td>ST_RTIME</td>
<td>1 argument for reference time: new time or −1 for TOD.</td>
</tr>
<tr>
<td>ST_FILEFMT</td>
<td>Format of the file. To specify the format, you can specify a numeric value (see [Appendix A]), or one of the following predefined variables used to derive the appropriate numeric value:</td>
</tr>
<tr>
<td>S_FFBINARY</td>
<td>Binary data</td>
</tr>
<tr>
<td>S_FFCR</td>
<td>Text data delimited by a carriage return character</td>
</tr>
<tr>
<td>S_FFCRLF</td>
<td>Text data delimited by carriage return and line feed characters</td>
</tr>
<tr>
<td>S_FFLF</td>
<td>Text data delimited by a line feed character</td>
</tr>
<tr>
<td>S_FFLFCR</td>
<td>Text data delimited by a line feed and carriage return characters</td>
</tr>
<tr>
<td>S_FFNA</td>
<td>Text data with the file format not specified</td>
</tr>
<tr>
<td>S_FFNL</td>
<td>Text data delimited by a newline character</td>
</tr>
</tbody>
</table>
Usage notes

1. Some of the attributes changed by the chattr service can also be changed by other services.

2. When changing the mode:
   - The effective UID of the calling process must match the file's owner UID, or the caller must have appropriate privileges.
   - Setting the set-group-ID-on-execution permission (in mode) means that when this file is run (through the exec service), the effective GID of the caller is set to the file's owner GID, so that the caller seems to be running under the GID of the file, rather than that of the actual invoker.
   - The set-group-ID-on-execution permission is set to zero if both of the following are true:
     - The caller does not have appropriate privileges.
     - The GID of the file's owner does not match the effective GID, or one of the supplementary GIDs, of the caller.
   - Setting the set-user-ID-on-execution permission (in mode) means that when this file is run, the process's effective UID is set to the file's owner UID, so that the process seems to be running under the UID of the file's owner, rather than that of the actual invoker.

3. When changing the owner:
   - For changing the owner UID of a file, the caller must have appropriate privileges.
   - For changing the owner GID of a file, the caller must have appropriate privileges, or meet all of these conditions:
     - The effective UID of the caller matches the file's owner UID.
     - The owner UID value specified in the change request matches the file's owner UID.
     - The GID value specified in the change request is the effective GID, or one of the supplementary GIDs, of the caller.
   - When changing the owner, the set-user-ID-on-execution and set-group-ID-on-execution permissions of the file mode are automatically turned off.
   - When the owner is changed, both UID and GID must be specified as they are to be set. If you want to change only one of these values, you need to set the other to its present value for it to remain unchanged.

4. For general attribute bits to be changed, the calling process must have write permission for the file.

5. When changing the file size:
   - The change is made beginning from the first byte of the file. If the file was previously larger than the new size, the data from \texttt{file\_size} to the original end of the file is removed. If the file was previously shorter than \texttt{file\_size}, bytes between the old and new lengths are read as zeros. The file offset is not changed.
   - If \texttt{file\_size} is greater than the current file size limit for the process, the request fails with EFBIG, and the SIGXFSZ signal is generated for the process.
   - Successful change clears the set-user-ID, the set-group-ID, and the save-text (sticky bit) attributes of the file, unless the caller is a superuser.

6. When changing times:
chattr

- For the access time or the modification time to be set explicitly (using either
  `st_atime` or `st_mtime` with the new time), the effective ID must match that of
  the file's owner, or the process must have appropriate privileges.

- For the access time or modification time to be set to the current time (using
  either `st_atime` or `st_mtime` with −1), the effective ID must match that of the
  file's owner, the calling process must have write permission for the file, or the
  process must have appropriate privileges.

- For the change time or the reference time to be set explicitly (using either
  `st_ctime` or `st_rtime` with the new time), the effective ID must match that of
  the file's owner, or the process must have appropriate privileges.

- For the change time or reference time to be set to the current time (using
  either `st_ctime` or `st_rtime` with −1), the calling process must have write
  permission for the file.

- When any attribute field is changed successfully, the file's change time is also
  updated.

7. For auditor audit flags to be changed, the user must have auditor authority. The
   user with auditor authority can set the auditor options for any file, even those to
   which they do not have path access or authority to use for other purposes.
   
   Auditor authority is established by issuing the TSO/E command ALTUSER
   AUDITOR.

8. For the user audit flags to be changed, the user must have appropriate
   privileges or be the owner of the file.

Example
To set permissions for `/u/project` to 775:

"chattr /u/project" st_mode 775
Function
chaudit invokes the chaudit callable service to change audit flags for a file.

Parameters
pathname
The pathname of the file.

audit_flags
One or more numeric values that indicate the type of access to be tested. You can specify a numeric value (see Appendix A) or the predefined variable used to derive the appropriate numeric value. The predefined variables you can specify are:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>AUD_FREAD</td>
<td>Audit failed read requests</td>
</tr>
<tr>
<td>AUD_SREAD</td>
<td>Audit successful read requests</td>
</tr>
<tr>
<td>AUD_FWRITE</td>
<td>Audit failed write requests</td>
</tr>
<tr>
<td>AUD_SWRITE</td>
<td>Audit successful write requests</td>
</tr>
<tr>
<td>AUD_FEXEC</td>
<td>Audit failed execute or search requests</td>
</tr>
<tr>
<td>AUD_SEXEC</td>
<td>Audit successful execute or search requests</td>
</tr>
</tbody>
</table>

option
A number indicating whether user-requested or auditor-requested auditing is being changed:
- 0 if user-requested auditing is being changed.
- 1 if auditor-requested auditing is being changed.

Usage notes
1. If option indicates that the auditor audit flags are to be changed, you must have auditor authority for the request to be successful. If you have auditor authority, you can set the auditor options for any file, even those to which you do not have path access or authority to use for other purposes.
   You can get auditor authority by entering the TSO/E command ALTUSER AUDITOR.
2. If option indicates that the user audit flags are to be changed, you must have appropriate privileges or be the owner of the file.

Example
In the following example, assume that pathname was assigned a value earlier in the exec. To change user-requested auditing so that failed read, write, and execute attempts for pathname are audited:

```
"chaudit (pathname)" aud_fread+aud_fwrite+aud_fexec 0
```
Function
chdir invokes the chdir callable service to change the working directory.

Parameters
pathname
  The pathname of the directory.

Usage notes
If you use chdir to change a directory in a REXX program that is running in a
TSO/E session, the directory is typically reset to your home directory when the
REXX program ends. When a REXX program changes directories and then exits,
the thread is undubbed. If this was the only thread dubbed in your TSO/E session,
the working directory is reset to the home directory the next time a syscall
command is issued. However, if there is more than one dubbed thread in the
address space, the remaining threads keep the working directory even when the
REXX program exits.

Example
To change the working directory to /u/lou/dirb:
"chdir /u/lou/dirb"
Function

`chmod` invokes the chmod callable service to change the mode of a file or directory.

Parameters

- **pathname**
  The pathname of the file or directory.

- **mode**
  A three- or four-digit number, corresponding to the access permission bits. Each digit must be in the range 0–7, and at least three digits must be specified. For more information on permissions, see Appendix B.

- **setuid**
  Sets the set-user-ID-on-execution permission. Specify 1 to set this permission on, or 0 to set it off. The default is 0.

- **setgid**
  Sets the set-group-ID-on-execution permission. Specify 1 to set this permission on, or 0 to set it off. The default is 0.

- **sticky**
  The sticky bit for a file indicates where the file should be fetched from. If the file resides in the link pack area (LPA), link list, or STEPLIB, specify 1. The default is 0.

Setting the sticky bit for a directory to 1 indicates that to delete or rename a file, the effective user ID of the process must be the same as that of the directory owner or file owner, or that of a superuser. Setting the sticky bit for a directory to 0 indicates that anyone who has write permission to the directory can delete or rename a file.

Usage notes

1. One bit sets permission for set-user-ID on access, set-group-ID on access, or the sticky bit. You can set this bit in either of two ways:
   - Specifying four digits on the `mode` parameter; the first digit sets the bit.
   - Specifying the `setuid`, `setgid`, or sticky parameters.
2. When a `chmod` or `fchmod` has occurred for an open file, `fstat` reflects the change in mode. However, no change in access authorization is apparent when the file is accessed through a previously opened file descriptor.
3. For mode bits to be changed, the effective UID of the caller must match the file’s owner UID, or the caller must be a superuser.
4. When the mode is changed successfully, the file’s change time is also updated.
5. Setting the set-group-ID-on-execution permission means that when this file is run (through the exec service), the effective GID of the caller is set to the file’s owner GID, so that the caller seems to be running under the GID of the file, rather than that of the actual invoker.
The set-group-ID-on-execution permission is set to zero if both of the following are true:

- The caller does not have appropriate privileges.
- The GID of the file’s owner does not match the effective GID or one of the supplementary GIDs of the caller.

6. Setting the set-user-ID-on-execution permission means that when this file is run, the process’s effective UID is set to the file’s owner UID, so that the process seems to be running under the UID of the file’s owner, rather than that of the actual invoker.

**Example**

In the following example, assume that `pathname` was assigned a value earlier in the exec. This example changes the mode of the file to read-write-execute for the owner, and read-execute for all others:

```bash
"chmod (pathname) 755"
```
**Function**

`chown` invokes the chown callable service to change the owner or group for a file or directory.

**Parameters**

- `pathname`
  - The pathname of a file or directory.
- `uid`
  - The numeric UID for the new owner of the file or the present UID, or −1 if there is no change.
- `gid`
  - The numeric GID for the group for the file or the present GID, or −1 if there is no change.

**Usage notes**

1. The `chown` service changes the owner UID and owner GID of a file. Only a superuser can change the owner UID of a file.
2. The owner GID of a file can be changed by a superuser, or if a caller meets all of these conditions:
   - The effective UID of the caller matches the file’s owner UID.
   - The `uid` value specified in the change request matches the file’s owner UID.
   - The `gid` value specified in the change request is the effective GID, or one of the supplementary GIDs, of the caller.
3. The set-user-ID-on-execution and set-group-ID-on-execution permissions of the file mode are automatically turned off.
4. If the change request is successful, the change time for the file is updated.
5. Values for both `uid` and `gid` must be specified as they are to be set. If you want to change only one of these values, the other must be set to its present value to remain unchanged.

**Example**

In the following example, assume that `pathname`, `uid`, and `gid` were assigned a value earlier in the exec:

```
"chown (pathname) (uid) (gid)"
```
Function

`close` invokes the close callable service to close a file.

Parameters

- **fd** The file descriptor (a number) for the file to be closed.

Usage notes

1. Closing a file closes, or frees, the file descriptor by which the file was known to the process. The system can then reassign the file descriptor to the same file or to another file when it is opened.
2. Closing a file descriptor also unlocks all outstanding byte range locks that a process has on the associated file.
3. If a file has been opened by more than one process, each process has a file descriptor. When the last open file descriptor is closed, the file itself is closed. If the file's link count is zero at that time, the file's space is freed and the file becomes inaccessible. When the last open file descriptor for a pipe or FIFO special file is closed, any data remaining in the file is discarded.

Example

In the following example, assume that `fd` was assigned a value earlier in the exec:

```
close fd
```
Function
closedir invokes the closedir callable service to close a directory.

Parameters
fd The file descriptor (a number) for the directory to be closed.

Usage notes
closedir closes a directory file descriptor opened by the opendir syscall command. The rddir command reads a directory in the readdir callable service format. You can use opendir, rewinddir, and closedir together with the rddir syscall command, but not with the readdir syscall command. Alternatively, you can simply use the readdir syscall command to read an entire directory and format it in a stem.

Example
In the following example, assume that fd was assigned a value earlier in the exec:
"closedir" fd
Function
create invokes the open callable service to open a new file. The file descriptor is returned in RETVAL.

Parameters

pathname
The pathname of a file.

mode
A three- or four-digit number, corresponding to the access permission bits. Each digit must be in the range 0–7, and at least three digits must be specified. For more information on permissions, see Appendix B.

Usage notes
Using create is the equivalent of using the open callable service with the create, truncate, and write-only options:

- When a file is created with the create option, the file permission bits as specified in mode are modified by the process's file creation mask (see "umask" on page 188) and then used to set the file permission bits of the file being created.
- The truncate option opens the file as though it had been created earlier, but never written into. The mode and owner of the file do not change (although the change time and modification time do), but the file's contents are discarded. The file offset, which indicates where the next write is to occur, points to the first byte of the file.

Example
To open a new file, /u/lou/test.exec, with read-write-execute permission for the owner only:
"create /u/lou/test.exec 700"
dup

Function
dup invokes the fcntl callable service to duplicate an open file descriptor. The file
descriptor is returned in RETVAL.

Parameters
fd  An opened file descriptor (a number) to be duplicated.

Usage notes
dup fd is equivalent to F_DUPFD fd 0.

Example
In the following example, assume that fd was assigned a value earlier in the exec:
"dup (fd)"
Function
dup2 invokes the fcntl callable service to duplicate an open file descriptor to the file descriptor of choice. The file descriptor returned is equal to \( fd2 \). If \( fd2 \) is already in use, it is closed and \( fd \) is duplicated. If \( fd \) is equal to \( fd2 \), \( fd2 \) is returned without closing it. The file descriptor is returned in RETVAL.

Parameters
\( fd \) An opened file descriptor (a number) to be duplicated.
\( fd2 \) The file descriptor (a number) to be changed.

Usage notes
dup \( fd \) \( fd2 \) is equivalent to F_DUPFD \( fd \) \( fd2 \).

Example
In the following example, assume that \( fd1 \) and \( fd2 \) were assigned values earlier in the exec:
"dup2" \( fd1 \) \( fd2 \)
There is no exec syscall command. Instead of using exec, see "spawn" on page 172.
Function
`extlink` invokes the `extlink` callable service to create a symbolic link to an external name. This creates a symbolic link file.

Parameters
extname
   The external name of the file for which you are creating a symbolic link.

linkname
   The pathname for the symbolic link.

Usage notes
1. The object identified by `extname` need not exist when the symbolic link is created, and refers to an object outside a hierarchical file system.
2. The external name contained in an external symbolic link is not resolved. The `linkname` cannot be used as a directory component of a pathname.

Example
To create a symbolic link named `mydsn` for the file `WJS.MY.DSN`:
"extlink WJS.MY.DSN /u/wjs/mydsn"
Function

fchattr invokes the fchattr callable service to modify the attributes that are associated with a file represented by a file descriptor. You can change the mode, owner, access time, modification time, change time, reference time, audit flags, general attribute flags, and file size.

Parameters

fd  The file descriptor for the file.

attribute_list

A list of attributes to be set and their values. The attributes are expressed either as numeric values (see Appendix A), or as the predefined variables beginning with ST_ followed by arguments for that attribute. The attributes that may be changed and their parameters are:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
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<td>Coded character set ID; first 4 characters are the file tag.</td>
</tr>
<tr>
<td>ST_MODE</td>
<td>1 argument: permission bits as 3 octal digits.</td>
</tr>
<tr>
<td>ST_UID</td>
<td>2 arguments: UID and GID numbers.</td>
</tr>
<tr>
<td>ST_SIZE</td>
<td>1 argument: new file size.</td>
</tr>
<tr>
<td>ST_ATIME</td>
<td>1 argument for access time: new time or -1 for TOD.</td>
</tr>
<tr>
<td>ST_MTIME</td>
<td>1 argument for modification time: new time or -1 for TOD.</td>
</tr>
<tr>
<td>ST_CTIME</td>
<td>1 argument for change time: new time or -1 for TOD.</td>
</tr>
<tr>
<td>ST_SETUID</td>
<td>No arguments.</td>
</tr>
<tr>
<td>ST_SETGID</td>
<td>No arguments.</td>
</tr>
<tr>
<td>ST_AAUDIT</td>
<td>1 argument: new auditor audit value.</td>
</tr>
<tr>
<td>ST_UAUDIT</td>
<td>1 argument: new user audit value.</td>
</tr>
<tr>
<td>ST_STICKY</td>
<td>No arguments.</td>
</tr>
<tr>
<td>ST_GENVALUE</td>
<td>2 arguments: names of two variables. The first variable contains the general attribute mask and the second contains the general attribute value.</td>
</tr>
<tr>
<td>ST_RTIME</td>
<td>1 argument for reference time: new time or -1 for TOD.</td>
</tr>
</tbody>
</table>
| ST.FILEFMT   | Format of the file. To specify the format, you can specify a numeric value (see Appendix A) or one of the following predefined variables used to derive the appropriate numeric value: S_FFBINARY Binary data S_FFCR Text data delimited by a carriage return character S_FFCRLF Text data delimited by carriage return and line feed characters S_FFLF Text data delimited by a line feed character S_FFLFCR Text data delimited by a line feed and carriage return characters S_FFNA Text data with the file format not specified S_FFNL Text data delimited by a newline character
Usage notes

1. Some of the attributes changed by the fchattr service can also be changed by other services.

2. When changing the mode:
   - The effective UID of the calling process must match the file’s owner UID, or the caller must have appropriate privileges.
   - Setting the set-group-ID-on-execution permission (in mode) means that when this file is run (through the exec service), the effective GID of the caller is set to the file’s owner GID, so that the caller seems to be running under the GID of the file, rather than that of the actual invoker.
     The set-group-ID-on-execution permission is set to zero if both of the following are true:
     - The caller does not have appropriate privileges.
     - The GID of the file’s owner does not match the effective GID, or one of the supplementary GIDs, of the caller.
   - Setting the set-user-ID-on-execution permission (in mode) means that when this file is run, the process’s effective UID is set to the file’s owner UID, so that the process seems to be running under the UID of the file’s owner, rather than that of the actual invoker.

3. When changing the owner:
   - For changing the owner UID of a file, the caller must have appropriate privileges.
   - For changing the owner GID of a file, the caller must have appropriate privileges, or meet all of these conditions:
     - The effective UID of the caller matches the file’s owner UID.
     - The owner UID value specified in the change request matches the file’s owner UID.
     - The GID value specified in the change request is the effective GID, or one of the supplementary GIDs, of the caller.
   - When the owner is changed, the set-user-ID-on-execution and set-group-ID-on-execution permissions of the file mode are automatically turned off.
   - When the owner is changed, both UID and GID must be specified as they are to be set. If you want to change only one of these values, you need to set the other to its present value for it to remain unchanged.

4. For general attribute bits to be changed, the calling process must have write permission for the file.

5. When changing the file size:
   - The change is made beginning from the first byte of the file. If the file was previously larger than the new size, the data from file_size to the original end of the file is removed. If the file was previously shorter than file_size, bytes between the old and new lengths are read as zeros. The file offset is not changed.
   - If file_size is greater than the current file size limit for the process, the request fails with EFBIG and the SIGXFSZ signal is generated for the process.
   - Successful change clears the set-user-ID, set-group-ID, and save-text (sticky bit) attributes of the file unless the caller is a superuser.

6. When changing times:
fchattr

- For the access time or the modification time to be set explicitly (using either `st_atime` or `st_mtime` with the new time), the effective ID must match that of the file's owner, or the process must have appropriate privileges.
- For the access time or modification time to be set to the current time (using either `st_atime` or `st_mtime` with −1), the effective ID must match that of the file's owner, the calling process must have write permission for the file, or the process must have appropriate privileges.
- For the change time or the reference time to be set explicitly (using either `st_ctime` or `st_rtime` with the new time) the effective ID must match that of the file's owner, or the process must have appropriate privileges.
- For the change time or reference time to be set to the current time (using either `st_ctime` or `st_rtime` with −1), the calling process must have write permission for the file.
- When any attribute field is changed successfully, the file's change time is also updated.

7. For auditor audit flags to be changed, the user must have auditor authority. The user with auditor authority can set the auditor options for any file, even those to which they don't have path access or authority to use for other purposes.
   Auditor authority is established by issuing the TSO/E command ALTUSER AUDITOR.

8. For the user audit flags to be changed, the user must have appropriate privileges or be the owner of the file.

Example
In the following example, assume that `fd` was assigned a value earlier in the exec. This truncates a file to 0 bytes and sets the file permissions to 600:

```
"fchattr" fd st_size 0 st_mode 600
```
Function
Function
fchaudit invokes the fchaudit callable service to change audit flags for a file
identified by a file descriptor. The file descriptor is specified by a number.

Parameters
Parameters
fd  The file descriptor for the file.

audit_flags
One or more numeric values that indicate the type of access to be tested. You
can specify a numeric value (see Appendix A) or the predefined variable used
to derive the appropriate numeric value. The predefined variables you can
specify are:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>AUD_FREAD</td>
<td>Audit failed read requests</td>
</tr>
<tr>
<td>AUD_SREAD</td>
<td>Audit successful read requests</td>
</tr>
<tr>
<td>AUD_FWRITE</td>
<td>Audit failed write requests</td>
</tr>
<tr>
<td>AUD_SWRITE</td>
<td>Audit successful write requests</td>
</tr>
<tr>
<td>AUD_FEXEC</td>
<td>Audit failed execute or search requests</td>
</tr>
<tr>
<td>AUD_SEXEC</td>
<td>Audit successful execute or search requests</td>
</tr>
</tbody>
</table>

option
A number indicating whether user-requested or auditor-requested auditing is
being changed:
• 0 if user-requested auditing is being changed.
• 1 if auditor-requested auditing is being changed.

Usage notes
1. If option indicates that the auditor audit flags are to be changed, you must have
auditor authority for the request to be successful. If you have auditor authority,
you can set the auditor options for any file, even those to which you do not
have path access or authority to use for other purposes.

You can get auditor authority by entering the TSO/E command ALTUSER
AUDITOR.

2. If option indicates that the user audit flags are to be changed, you must have
appropriate privileges or be the owner of the file.

Example
To change user-requested auditing so that failed read requests for the file identified
by file descriptor 0 are audited:
"fchaudit 0 (aud_fread) 0"
fchmod

Function

fchmod invokes the fchmod callable service to change the mode of a file or directory indicated by a file descriptor. The file descriptor is specified by a number.

Parameters

fd    The file descriptor for the file or directory.

mode  A three- or four-digit number, corresponding to the access permission bits. Each digit must be in the range 0–7, and at least three digits must be specified. For more information on permissions, see [Appendix B].

setuid  Sets the set-user-ID-on-execution permission. Specify 1 to set this permission on, or 0 to set it off. The default is 0.

setgid  Sets the set-group-ID-on-execution permission. Specify 1 to set this permission on, or 0 to set it off. The default is 0.

sticky  Sets the sticky bit to indicate where the file should be fetched from. If the file resides in the link pack area (LPA), link list, or STEPLIB, specify 1. The default is 0.

Usage notes

1. One bit sets permission for set-user-ID on access, set-group-ID on access, or the sticky bit. You can set this bit in either of two ways:
   - Specifying four digits on the mode parameter; the first digit sets the bit.
   - Specifying the setuid, setgid, or sticky parameters.

2. When a chmod or fchmod has occurred for an open file, fstat reflects the change in mode. However, no change in access authorization is apparent when the file is accessed through a previously opened file descriptor.

3. For mode bits to be changed, the effective UID of the caller must match the file’s owner UID, or the caller must be a superuser.

4. When the mode is changed successfully, the file’s change time is also updated.

5. Setting the set-group-ID-on-execution permission means that when this file is run, through the exec service, the effective GID of the caller is set to the file’s owner GID, so that the caller seems to be running under the GID of the file, rather than that of the actual invoker.

The set-group-ID-on-execution permission is set to zero if both of the following are true:
   - The caller does not have appropriate privileges.
   - The GID of the file’s owner does not match the effective GID or one of the supplementary GIDs of the caller.
6. Setting the set-user-ID-on-execution permission means that when this file is run, the process’s effective UID is set to the file’s owner UID, so that the process seems to be running under the UID of the file’s owner, rather than that of the actual invoker.

**Example**
In the following example, assume that `fd` was assigned a value earlier in the exec. This changes the mode for the file identified by the file descriptor so that only a superuser can access the file:

"fchmod (fd) 000"
Function

fchown invokes the fchown callable service to change the owner and group of a file or directory indicated by a file descriptor. The file descriptor is specified by a number.

Parameters

fd    The file descriptor for a file or directory.

uid   The numeric UID for the new owner of the file or the present UID, or −1 if there is no change.

gid   The numeric GID for the new group for the file or the present GID, or −1 if there is no change.

Usage notes

1. The fchown service changes the owner UID and owner GID of a file. Only a superuser can change the owner UID of a file.
2. The owner GID of a file can be changed by a superuser, or if a caller meets all of these conditions:
   • The effective UID of the caller matches the file’s owner UID.
   • The uid value specified in the change request matches the file’s owner UID.
   • The gid value specified in the change request is the effective GID, or one of the supplementary GIDs, of the caller.
3. The set-user-ID-on-execution and set-group-ID-on-execution permissions of the file mode are automatically turned off.
4. If the change request is successful, the change time for the file is updated.
5. Values for both uid and gid must be specified as they are to be set. If you want to change only one of these values, the other must be set to its present value to remain unchanged.

Example

In the following example, assume that fd, uid, and gid were assigned a value earlier in the exec:

"fchown" fd uid gid
**f_closfd**

**Function**

*f_closfd* invokes the fcntl callable service to close a range of file descriptors.

**Parameters**

*fd*  
The file descriptor (a number) for a file. This is the first file descriptor to be closed.

*fd2*  
The file descriptor (a number) for a file, which must be greater than or equal to *fd*. If a −1 is specified for *fd2*, all file descriptors greater than or equal to *fd* are closed.

**Usage notes**

1. A process can use *f_closfd* to close a range of file descriptors. *fd2* must be greater than or equal to *fd*, or it can also be −1, which indicates that all file descriptors greater than or equal to *fd* are to be closed.

2. Use of *f_closfd* is meant to be consistent with the close callable service. You cannot close file descriptors that could not also be closed using the close service.

3. When a file descriptor cannot be closed, it is considered an error, but the request continues with the next file descriptor in the range. File descriptors that are not in use are ignored.

**Example**

In the following example, assume that *fd* and *fd2* were assigned values earlier in the exec:

"f_closfd" fd fd2"
Function

*f_control_cvt* controls automatic file conversion and specifies the program and file CCSIDs (character code set identifiers) for an opened I/O stream.

Parameters

*fd* The file descriptor (a number) for the file. It must be a regular file, FIFO, or character special file.

*filetag* The name of a 4-byte hex variable describing the file tag. The CCSID (coded character set identifier) of the file tag occupies the first 2 bytes. For more information, see BPXYSTAT — Map the Response Structure for stat in z/OS UNIX System Services Programming: Assembler Callable Services Reference or the C header file *stat.h* in z/OS XL C/C++ Run-Time Library Reference.

*command* The *command* can be one of these variables:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>CVT_SETCVTOFF</td>
<td>Turns off any conversion that may be in effect. A hex value of 0 for both</td>
</tr>
<tr>
<td></td>
<td>program-ccsid and file-ccsid is recommended.</td>
</tr>
<tr>
<td>CVT_SETCVTON</td>
<td>Turns automatic conversion on for the stream and, optionally, sets the</td>
</tr>
<tr>
<td></td>
<td>program CCSID or file CCSID, or both. A hex value of 0 for program-ccsid</td>
</tr>
<tr>
<td></td>
<td>indicates using ThliCcsid (the current program CCSID) at the time of each</td>
</tr>
<tr>
<td></td>
<td>read or write. ThliCcsid is initially 1047, but can be reset directly by</td>
</tr>
<tr>
<td></td>
<td>the program, or indirectly by setting the appropriate run-time option or</td>
</tr>
<tr>
<td></td>
<td>environment variable. A hex value of 0 for file-ccsid indicates not changing</td>
</tr>
<tr>
<td></td>
<td>ThliCcsid.</td>
</tr>
<tr>
<td>CVT_SETAUTOCVTON</td>
<td>Conditionally turns automatic conversion on for the stream and, optionally,</td>
</tr>
<tr>
<td></td>
<td>sets the program CCSID or file CCSID, or both. Conversion will be in effect</td>
</tr>
<tr>
<td></td>
<td>for this stream only if the system or the local run-time environment has</td>
</tr>
<tr>
<td></td>
<td>been enabled for conversion. A hex value of 0 for program-ccsid indicates</td>
</tr>
<tr>
<td></td>
<td>using ThliCcsid at the time of each read or write. ThliCcsid is initially</td>
</tr>
<tr>
<td></td>
<td>1047, but can be changed by this variable, or by setting the appropriate</td>
</tr>
<tr>
<td></td>
<td>run-time option or environment variable. A hex value of 0 for file-ccsid</td>
</tr>
<tr>
<td></td>
<td>indicates not changing ThliCcsid.</td>
</tr>
<tr>
<td>CVT_QUERYCVT</td>
<td>Returns an indicator that lets you know whether automatic conversion is</td>
</tr>
<tr>
<td></td>
<td>in effect or not, and also returns the program and file CCSIDs that are</td>
</tr>
<tr>
<td></td>
<td>being used (if conversion is in effect). The variables command, program-</td>
</tr>
<tr>
<td></td>
<td>ccsid, and file-ccsid are set when you have a successful return. command is</td>
</tr>
<tr>
<td></td>
<td>set to either cvt_setcvtoff or cvt_setcvton. program-ccsid and file-ccsid</td>
</tr>
</tbody>
</table>
|                | are set to the CCSID values that would be in use if conversion were to occur.
**program-ccsid**

The name of a 2-byte hex variable that describes the CCSID for the running program.

**Note:** For EBCDIC (1047) the CCSID is '0417'X; for ASCII (819) it is '0333'X.

**file-ccsid**

The name of a 2-byte hex variable that describes the CCSID for the file opened with the file descriptor *fd*.

**Example**

This example turns automatic file conversion on for a stream opened with *fd* as the file descriptor:

```
pccsid = '0000'x
fccsid = '0000'x
cmd = cvt_setautocvt
"f_control_cvt (fd) cmd pccsid fccsid"
```

This example queries the current conversion state:

```
pccsid = '0000'x
fccsid = '0000'x
cmd = cvt_querycvt
"f_control_cvt (fd) cmd pccsid fccsid"
```
**Function**

fcntl is supported as a set of syscall commands whose names begin with f_:

- "f_closfd" on page 60
- "f_dupfd" on page 64
- "f_dupfd2" on page 65
- "f_getfd" on page 66
- "f_getfl" on page 67
- "f_getlk" on page 68
- "f_getlkw" on page 77
- "f_setfd" on page 75
- "f_setfl" on page 76
- "f_setlk" on page 77
- "f_setlkw" on page 80
Function

_f_dupfd_ invokes the fcntl callable service to duplicate the lowest file descriptor that is equal to or greater than _fd2_ and not already associated with an open file. The file descriptor is returned in RETVAL.

Parameters

_fd_ The file descriptor (a number) that you want to duplicate.

_fd2_ The file descriptor (a number) at which to start looking for an available file descriptor.

Example

In the following example, assume that _fd_ and _fd2_ were assigned values earlier in the exec:

"_f_dupfd_" fd fd2
Function

`f_dupfd2` invokes the fcntl callable service to duplicate a file descriptor that is equal to `fd2`.

Parameters

- **fd**: An opened file descriptor (a number) to be duplicated.
- **fd2**: The file descriptor of choice.

Usage notes

If `fd2` is already in use, `f_dupfd2` closes it. If `fd` is equal to `fd2`, `fd2` is returned but not closed.

Example

In the following example, assume that `fd` and `fd2` were assigned values earlier in the exec:

```
f_dupfd2 fd fd2
```
Function
f_getfd invokes the fcntl callable service to get the file descriptor flags for a file.

Parameters
fd The file descriptor (a number) for the file.

Usage notes
The file descriptor flags are returned in RETVAL. The only POSIX-defined flag is FCTLCLOEXEC. To determine if this flag is set, use this expression: (retval//2)=1

Example
To get the file descriptor flags for file descriptor 0:
"f_getfd 0"
Function

*f_getfl* invokes the fcntl callable service to get the file status flags for a file.

Parameters

*fd*  The file descriptor (a number) for the file.

Usage notes

RETVAL returns the file status flags as a numeric value (see Appendix A):

<table>
<thead>
<tr>
<th>Flag</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>O_CREAT</td>
<td>Create the file if it does not exist.</td>
</tr>
<tr>
<td>O_EXCL</td>
<td>Fail if the file does exist and O_CREAT is set.</td>
</tr>
<tr>
<td>O_NOCTTY</td>
<td>Do not make this file a controlling terminal for the calling process.</td>
</tr>
<tr>
<td>O_TRUNC</td>
<td>If the file exists, truncate it to zero length.</td>
</tr>
<tr>
<td>O_APPEND</td>
<td>Set the offset to EOF before each write.</td>
</tr>
<tr>
<td>O_NONBLOCK</td>
<td>An open, read, or write on the file will not block (wait for terminal input).</td>
</tr>
<tr>
<td>O_RDWR</td>
<td>Open for read and write.</td>
</tr>
<tr>
<td>O_RDONLY</td>
<td>Open for read-only.</td>
</tr>
<tr>
<td>O_WRONLY</td>
<td>Open for write-only.</td>
</tr>
<tr>
<td>O_SYNC</td>
<td>Force synchronous updates.</td>
</tr>
</tbody>
</table>

You can use the open flags to test specific values. The easiest way to test a value is to convert both the RETVAL and the flags to binary data, and then logically AND them. For example, to test O_WRITE and O_TRUNC:

```
wrtr=D2C(O_WRITE+O_TRUNC,4))
If BITAND(D2C(retval,4),wrtr)=wrtr Then
  Do /* o_write and o_trunc are set */
End
```

Example

In the following example, assume that *fd* was assigned a value earlier in the exec:

"*f_getfl"* *fd"
Function

f_getlk invokes the fcntl callable service to return information on a file segment for which locks are set, cleared, or queried.

Parameters

fd  The file descriptor (a number) for the file.

stem
The name of a stem variable that is the flock structure used to query, set, or clear a lock; or to return information. To specify the information, you can specify a numeric value (see Appendix A) or the predefined variables beginning with L_ used to derive the appropriate numeric value. For example, stem.1 and stem.L_type are both the lock-type request:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>L_LEN</td>
<td>The length of the byte range that is to be set, cleared, or queried.</td>
</tr>
<tr>
<td>L_PID</td>
<td>The process ID of the process holding the blocking lock, if one was found.</td>
</tr>
<tr>
<td>L_START</td>
<td>The starting offset byte of the lock that is to be set, cleared, or queried.</td>
</tr>
</tbody>
</table>
| L_TYPE     | The type of lock being set, cleared, or queried. To specify the information,
            | you can specify a numeric value (see Appendix A) or one of the following predefined variables used to derive the appropriate numeric value: |

F_RDLCK
Shared or read lock. This type of lock specifies that the process can read the locked part of the file, and other processes cannot write on that part of the file in the meantime. A process can change a held write lock, or any part of it, to a read lock, thereby making it available for other processes to read. Multiple processes can have read locks on the same part of a file simultaneously. To establish a read lock, a process must have the file accessed for reading.
<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>F_WRLCK</td>
<td>Exclusive or write lock. This type of lock indicates that the process can write on the locked part of the file, without interference from other processes. If one process puts a write lock on part of a file, no other process can establish a read lock or write lock on that same part of the file. A process cannot put a write lock on part of a file if there is already a read lock on an overlapping part of the file, unless that process is the only owner of that overlapping read lock. In such a case, the read lock on the overlapping section is replaced by the write lock being requested. To establish a write lock, a process must have the file accessed for writing.</td>
</tr>
<tr>
<td>F_UNLCK</td>
<td>Unlock. This is used to unlock all locks held on the given range by the requesting process.</td>
</tr>
<tr>
<td>L_WHENCE</td>
<td>The flag for the starting offset. To specify the information, you can specify a numeric value (see Appendix A) or one of the predefined variables beginning with SEEK_ used to derive the appropriate numeric value. Valid values are:</td>
</tr>
<tr>
<td>SEEK_CUR</td>
<td>The current offset</td>
</tr>
<tr>
<td>SEEK_END</td>
<td>The end of the file</td>
</tr>
<tr>
<td>SEEK_SET</td>
<td>The specified offset</td>
</tr>
</tbody>
</table>

Example
In the following example, assume that rec was assigned a value earlier in the exec:

```c
lock.l_len=40
lock.l_start=rec*40
lock.l_type=f_wrlck
lock.l_whence=seek_set
"f_getlk" fd "lock."
if lock.l_type=f_unlck then
 /* lock is available for the requested 40 byte record */
```
There is no `fork` syscall command. Instead of using `fork`, see "spawn" on page 172.
forkexecm

Function
forkexecm invokes the fork and execmvs callable services to fork and exec a program to be executed from the LINKLIB, LPALIB, or STEPLIB library.

Parameters
program
   The name of the program.
parm
   A parameter to be passed to the program.

Usage notes
1. If the exec fails, the child process ends with a SIGABRT signal.
2. If the fork succeeds, RETVAL contains the PID for the child process.
3. The child process has a unique process ID (PID) that does not match any active process group ID.
4. If a hierarchical file system (HFS) file has its FCTL_CLOFORK flag set on, it is not inherited by the child process. This flag is set with the fcntl callable service.
5. The process and system utilization times for the child are set to zero.
6. Any file locks previously set by the parent are not inherited by the child.
7. The child process has no alarms set (similar to the results of a call to the alarm service with Wait_time specified as zero).
8. The child has no pending signals.
9. The following characteristics of the calling process are changed when the new executable program is given control by the execmvs callable service:
   • The prior process image is replaced with a new process image for the executable program to be run.
   • All open files marked close-on-exec and all open directory streams are closed.
   • All signals that have sigaction settings are reset to their default actions.
10. The input passed to the MVS executable file by the service is consistent with what is passed as input to MVS programs. On input, the MVS program receives a single-entry parameter list pointed to by register 1. The high-order bit of the single parameter entry is set to 1.
    The single parameter entry is the address of a 2-byte length field followed by an argument string. The length field describes the length of the data that follows it. If a null argument and argument length are specified in the call, the length field specifies 0 bytes on input to the executable file.
11. The call can invoke both unauthorized and authorized MVS programs:
   • Unauthorized programs receive control in problem program state, with PSW key 8.
   • Authorized programs receive control in problem program state, with PSW key 8 and APF authorization.
forkexecm

12. The TASKLIB, STEPLIB, or JOBLIB DD data set allocations that are active for the calling task at the time of the call to the execmvs service are propagated to the new process image, if the data sets they represent are found to be cataloged. Uncataloged data sets are not propagated to the new process image. This causes the program that is invoked to run with the same MVS program search order as its invoker.

Example
In the following example, assume that \textit{pgm} was assigned a value earlier in the exec:

\begin{verbatim}
"forkexecm (pgm) 'hello world!'"
\end{verbatim}
Function

`fpathconf` invokes the `fpathconf` callable service to let a program determine the current value of a configurable limit or variable associated with a file or directory. The value is returned in RETVAL.

Parameters

- **fd** The file descriptor (a number) for the file or directory.
- **name** A numeric value that indicates which configurable limit will be returned. You can specify a numeric value (see Appendix A) or one of the following predefined variables used to derive the appropriate numeric value:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>PC_ACL</td>
<td>Test if access control lists (ACLs) are supported.</td>
</tr>
<tr>
<td>PC_ACL_MAX</td>
<td>Maximum number of entries allowed in an ACL.</td>
</tr>
<tr>
<td>PC_LINK_MAX</td>
<td>Maximum value of a file’s link count.</td>
</tr>
<tr>
<td>PC_MAX_CANON</td>
<td>Maximum number of bytes in a terminal canonical input line.</td>
</tr>
<tr>
<td>PC_MAX_INPUT</td>
<td>Minimum number of bytes for which space will be available in a terminal input queue; therefore, the maximum number of bytes a portable application may require to be typed as input before reading them.</td>
</tr>
<tr>
<td>PC_NAME_MAX</td>
<td>Maximum number of bytes in a filename (not a string length; count excludes a terminating null).</td>
</tr>
<tr>
<td>PC_PATH_MAX</td>
<td>Maximum number of bytes in a pathname (not a string length; count excludes a terminating null).</td>
</tr>
<tr>
<td>PCPIPE_BUF</td>
<td>Maximum number of bytes that can be written atomically when writing to a pipe.</td>
</tr>
<tr>
<td>PC_POSIX_CHOWN_RESTRICTED</td>
<td>Change ownership function is restricted to a process with appropriate privileges, and to changing the group ID (GID) of a file only to the effective group ID of the process or to one of its supplementary group IDs.</td>
</tr>
<tr>
<td>PC_POSIX_NO_TRUNC</td>
<td>Pathname components longer than 255 bytes generate an error.</td>
</tr>
<tr>
<td>PC_POSIX_VDISABLE</td>
<td>Terminal attributes maintained by the system can be disabled using this character value.</td>
</tr>
</tbody>
</table>
Usage notes

1. If `name` refers to MAX_CANON, MAX_INPUT, or _POSIX_VDISABLE, the following applies:
   - If `fd` does not refer to a terminal file, the function returns −1 and sets the ERRNO to EINVAL.

2. If `name` refers to NAME_MAX, PATH_MAX, or _POSIX_NO_TRUNC, the following applies:
   - If `fd` does not refer to a directory, the function still returns the requested information using the parent directory of the specified file.

3. If `name` refers to PC_PIPE_BUF, the following applies:
   - If `fd` refers to a pipe or a FIFO, the value returned applies to the referred-to object itself. If `fd` refers to a directory, the value returned applies to any FIFOs that exist or that can be created within the directory. If `fd` refers to any other type of file, the function returns −1 in RETVAL and sets the ERRNO to EINVAL.

4. If `name` refers to PC_LINK_MAX, the following applies:
   - If `fd` refers to a directory, the value returned applies to the directory.

Example

To determine the maximum number of bytes that can be written atomically to the file identified by file descriptor 1:

```
fpathconf 1 (pc_pipe_buf)
```
Function

`f_setfd` invokes the `fcntl` callable service to set file descriptor flags.

Parameters

- `fd` The file descriptor (a number) for the file.
- `close_exec` A numeric value to indicate whether this file descriptor should remain open after an exec:
  - 0 indicates that it should remain open.
  - 1 indicates that it should be closed.

Example

To set the flags for the file identified by file descriptor 0 and indicate that this file descriptor should remain open during an exec:

```
f_setfd 0 0
```
Function
f_setfl invokes the fcntl callable service to set file status flags.

Parameters
fd The file descriptor (a number) for the file.

flags
A value that sets the file status flags. To specify the information, you can specify a numeric value (see Appendix A) or a predefined variable beginning with O_ used to derive the appropriate numeric value. The permitted values are:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>O_APPEND</td>
<td>Set offset to EOF on write.</td>
</tr>
<tr>
<td>O_NONBLOCK</td>
<td>Do not block an open, a read, or a write on the file (do not wait for terminal input).</td>
</tr>
<tr>
<td>O_SYNC</td>
<td>Force synchronous updates.</td>
</tr>
</tbody>
</table>

Example
To set the O_APPEND file status flag for the file identified by file descriptor 1:
"f_setfl 1" o_append
**Function**

*f_setlk* invokes the fcntl callable service to set or release a lock on part of a file.

**Parameters**

<table>
<thead>
<tr>
<th>Parameter</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>fd</strong></td>
<td>The file descriptor (a number) for the file.</td>
</tr>
<tr>
<td><strong>stem</strong></td>
<td>The name of a stem variable that is the flock structure used to set or release the lock. To access the information, you can specify a numeric value (see Appendix A) or the predefined variables beginning with L_ that derive the appropriate numeric value. For example, <em>stem.1</em> and <em>stem.l_type</em> are both the lock-type request.</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>L_LEN</td>
<td>The length of the byte range that is to be set, cleared, or queried.</td>
</tr>
<tr>
<td>L_PID</td>
<td>The process ID of the process holding the blocking lock, if one was found.</td>
</tr>
<tr>
<td>L_START</td>
<td>The starting offset byte of the lock that is to be set, cleared, or queried.</td>
</tr>
<tr>
<td>L_TYPE</td>
<td>The type of lock being set, cleared, or queried. To specify the information, you can specify a numeric value (see Appendix A) or one of the following predefined variables used to derive the appropriate numeric value:</td>
</tr>
</tbody>
</table>

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>F_RDLCK</td>
<td>Shared or read lock. This type of lock specifies that the process can read the locked part of the file, and other processes cannot write on that part of the file in the meantime. A process can change a held write lock, or any part of it, to a read lock, thereby making it available for other processes to read. Multiple processes can have read locks on the same part of a file simultaneously. To establish a read lock, a process must have the file accessed for reading.</td>
</tr>
</tbody>
</table>
f_setlk

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>F_WRLCK</td>
<td>Exclusive or write lock. This type of lock indicates that the process can write on the locked part of the file, without interference from other processes. If one process puts a write lock on part of a file, no other process can establish a read lock or write lock on that same part of the file. A process cannot put a write lock on part of a file if there is already a read lock on an overlapping part of the file, unless that process is the only owner of that overlapping read lock. In such a case, the read lock on the overlapping section is replaced by the write lock being requested. To establish a write lock, a process must have the file accessed for writing.</td>
</tr>
<tr>
<td>F_UNLCK</td>
<td>Unlock. This is used to unlock all locks held on the given range by the requesting process.</td>
</tr>
<tr>
<td>SEEK_CUR</td>
<td>The current offset</td>
</tr>
<tr>
<td>SEEK_END</td>
<td>The end of the file</td>
</tr>
<tr>
<td>SEEK_SET</td>
<td>The specified offset</td>
</tr>
</tbody>
</table>

Usage notes
If the lock cannot be obtained, a RETVAL of −1 is returned along with an appropriate ERRNO and ERRNOJR. You can also use F_SETLK to release locks already held, by setting l_type to F_UNLCK.

Multiple lock requests: A process can have several locks on a file simultaneously, but it can have only one type of lock set on any given byte. Therefore, if a process puts a new lock on part of a file that it had previously locked, the process has only one lock on that part of the file and the lock type is the one given by the most recent locking operation.

Releasing locks: When an f_setlk or f_setlkw request is made to unlock a byte region of a file, all locks held by that process within the specified region are released. In other words, each byte specified on an unlock request is freed from any lock that is held against it by the requesting process.

All of a process’s locks on a file are removed when the process closes a file descriptor for that file. Locks are not inherited by child processes created with the fork service.

Advisory locking: All locks are advisory only. Processes can use locks to inform each other that they want to protect parts of a file, but locks do not prevent I/O on the locked parts. A process that has appropriate permissions on a file can perform whatever I/O it chooses, regardless of which locks are set. Therefore, file locking is only a convention, and it works only when all processes respect the convention.
Example
The following example locks a 40-byte record (rec). Assume that rec was assigned a value earlier in the exec:

```
lock.l_len=40
lock.l_start=rec*40
lock.l_type=f_wrlck
lock.l_whence=seek_set
"f_setlk (fd) lock."
```
f_setlkw

f_setlkw

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>L_LEN</td>
<td>The length of the byte range that is to be set, cleared, or queried.</td>
</tr>
<tr>
<td>L_PID</td>
<td>The process ID of the process holding the blocking lock, if one was found.</td>
</tr>
<tr>
<td>L_START</td>
<td>The starting offset byte of the lock that is to be set, cleared, or queried.</td>
</tr>
<tr>
<td>L_TYPE</td>
<td>The type of lock being set, cleared, or queried. To specify the information, you can specify a numeric value (see Appendix A) or one of the following predefined variables used to derive the appropriate numeric value:</td>
</tr>
<tr>
<td>F_RDLCK</td>
<td>Shared or read lock. This type of lock specifies that the process can read the locked part of the file, and other processes cannot write on that part of the file in the meantime. A process can change a held write lock, or any part of it, to a read lock, thereby making it available for other processes to read. Multiple processes can have read locks on the same part of a file simultaneously. To establish a read lock, a process must have the file accessed for reading.</td>
</tr>
</tbody>
</table>

Function

f_setlkw invokes the fcntl callable service to set or release a lock on part of a file and, if another process has a lock on some or all of the requested range, wait until the specified range is free and the request can be completed.

fd The file descriptor (a number) for the file.

stem

The name of a stem variable that is the flock structure used to set or release the lock. To specify the lock information, you can specify a numeric value (see Appendix A) or the predefined variables beginning with L_ used to derive the appropriate numeric value. For example, stem.1 and stem.l_type are both the lock-type request.
f_setlk

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>F_WRLCK</td>
<td>Exclusive or write lock. This type of lock indicates that the process can write on the locked part of the file, without interference from other processes. If one process puts a write lock on part of a file, no other process can establish a read lock or write lock on that same part of the file. A process cannot put a write lock on part of a file if there is already a read lock on an overlapping part of the file, unless that process is the only owner of that overlapping read lock. In such a case, the read lock on the overlapping section is replaced by the write lock being requested. To establish a write lock, a process must have the file accessed for writing.</td>
</tr>
<tr>
<td>F_UNLCK</td>
<td>Unlock. This is used to unlock all locks held on the given range by the requesting process.</td>
</tr>
<tr>
<td>SEEK_CUR</td>
<td>The current offset</td>
</tr>
<tr>
<td>SEEK_END</td>
<td>The end of the file</td>
</tr>
<tr>
<td>SEEK_SET</td>
<td>The specified offset</td>
</tr>
</tbody>
</table>

Usage notes
If the lock cannot be obtained because another process has a lock on all or part of the requested range, the f_setlk request waits until the specified range becomes free and the request can be completed. You can also use f_setlk to release locks already held, by setting l_type to F_UNLCK.

If a signal interrupts a call to the fcntl service while it is waiting in an f_setlk operation, the function returns with a RETVAL of −1 and an ERRNO of EINTR.

f_setlk operations have the potential for encountering deadlocks. This happens when process A is waiting for process B to unlock a region, and B is waiting for A to unlock a different region. If the system detects that a f_setlk might cause a deadlock, the fcnti service returns with a RETVAL of −1 and an ERRNO of EDEADLK.

See "f_setlk" on page 77 for more information about locks.

Example
The following example locks a 40-byte record (rec). Assume that rec was assigned a value earlier in the exec:

```c
lock.l_len=40
lock.l_start=rec*40
lock.l_type=f_wrlck
lock.l_whence=seek_set
"f_setlkw (fd) lock."
```
Function

The function `f_settag` can be used to directly control basic file tagging for an opened file. The setting of the file tag may be immediate or deferred until the first write, depending upon the input parameters.

**fd**  The file descriptor (a number) for the file. It must be a regular file, FIFO, or character special file.


Example

The following example sets a file tag for an opened file to ASCII ISO88591–1, with the text conversion flag on as well:

```c
int tag = '03338000'x;
"f_settag (fd) tag"
```
Function

The `fstat` function invokes the fstat callable service to get status information about a file that is identified by its file descriptor.

Parameters

- **fd**: The file descriptor (a number) for the file.
- **stem**: The name of a stem variable used to return the status information. Upon return, `stem.0` contains the number of variables that are returned. To access the status values, you can use a numeric value (see Appendix A) or any of the predefined variables that begin with `ST_`. For example, `stem.9` and `stem.st_atime` are both the request for the time of last access.

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>ST_Audit</td>
<td>Auditor audit information</td>
</tr>
<tr>
<td>ST_ACCESSACL</td>
<td>1 if there is an access ACL (access control list)</td>
</tr>
<tr>
<td>ST_ATIME</td>
<td>Time of last access</td>
</tr>
<tr>
<td>ST_AUDITID</td>
<td>RACF File ID for auditing</td>
</tr>
<tr>
<td>ST_BLKSIZE</td>
<td>File block size</td>
</tr>
<tr>
<td>ST_BLOCKS</td>
<td>Blocks allocated</td>
</tr>
<tr>
<td>ST_CCSID</td>
<td>Coded character set ID; first 4 characters are the file tag</td>
</tr>
<tr>
<td>ST_CRTIME</td>
<td>File creation time</td>
</tr>
<tr>
<td>ST_CTIME</td>
<td>Time of last file status change</td>
</tr>
<tr>
<td>ST_DEV</td>
<td>Device ID of the file</td>
</tr>
<tr>
<td>ST_DMODELACL</td>
<td>1 if there is a directory model access control list (ACL)</td>
</tr>
<tr>
<td>ST_EXTLINK</td>
<td>External symbolic link flag, set to 0 or 1</td>
</tr>
<tr>
<td>ST_FID</td>
<td>File identifier</td>
</tr>
<tr>
<td>ST_FILEFMT</td>
<td>Format of the file. To specify the format, you can specify a numeric value</td>
</tr>
<tr>
<td></td>
<td>(see Appendix A) or one of the following predefined variables used to derive</td>
</tr>
<tr>
<td></td>
<td>the appropriate numeric value:</td>
</tr>
<tr>
<td>S_FFBINARY</td>
<td>Binary data</td>
</tr>
<tr>
<td>S_FFCR</td>
<td>Text data delimited by a carriage return character</td>
</tr>
<tr>
<td>S_FFCRLF</td>
<td>Text data delimited by carriage return and line feed characters</td>
</tr>
<tr>
<td>S_FFLF</td>
<td>Text data delimited by a line feed character</td>
</tr>
<tr>
<td>S_FFLFRCR</td>
<td>Text data delimited by a line feed and carriage return characters</td>
</tr>
<tr>
<td>S_FFNA</td>
<td>Text data with the file format not specified</td>
</tr>
<tr>
<td>S_FFNL</td>
<td>Text data delimited by a newline character</td>
</tr>
<tr>
<td>ST_FMODELACL</td>
<td>1 if there is a file model ACL</td>
</tr>
<tr>
<td>ST_GENVALUE</td>
<td>General attribute values</td>
</tr>
<tr>
<td>ST_GID</td>
<td>Group ID of the group of the file</td>
</tr>
<tr>
<td>ST_INO</td>
<td>File serial number</td>
</tr>
</tbody>
</table>
## fstat

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>ST_MAJOR</td>
<td>Major number for a character special file</td>
</tr>
<tr>
<td>ST_MINOR</td>
<td>Minor number for a character special file</td>
</tr>
<tr>
<td>ST_MODE</td>
<td>File mode, permission bits only</td>
</tr>
<tr>
<td>ST_MTIME</td>
<td>Time of last data modification</td>
</tr>
<tr>
<td>ST_NLINK</td>
<td>Number of links</td>
</tr>
<tr>
<td>ST_RTIME</td>
<td>File backup time stamp (reference time)</td>
</tr>
<tr>
<td>ST_SETGID</td>
<td>Set Group ID on execution flag, set to 0 or 1</td>
</tr>
<tr>
<td>ST_SETUID</td>
<td>Set User ID on execution flag, set to 0 or 1</td>
</tr>
<tr>
<td>ST_SIZE</td>
<td>File size for a regular file, in bytes. If file size exceeds 2^31−1 bytes, size is expressed in megabytes, using an M (for example, 3123M).</td>
</tr>
<tr>
<td>ST_SECLABEL</td>
<td>Security Label</td>
</tr>
<tr>
<td>ST_STICKY</td>
<td>Sticky bit flag (keep loaded executable in storage), set to 0 or 1</td>
</tr>
<tr>
<td>ST_TYPE</td>
<td>Numeric value that represents the file type for this file. You can use a numeric value (see Appendix A) or any of the predefined variables that begin with S_ to determine the file type:</td>
</tr>
<tr>
<td>S_ISCHR</td>
<td>Character special file</td>
</tr>
<tr>
<td>S_ISDIR</td>
<td>Directory</td>
</tr>
<tr>
<td>S_ISFIFO</td>
<td>FIFO special file</td>
</tr>
<tr>
<td>S_ISREG</td>
<td>Regular file</td>
</tr>
<tr>
<td>S_ISSYM</td>
<td>Symbolic link</td>
</tr>
<tr>
<td>ST_UAUDIT</td>
<td>Area for user audit information</td>
</tr>
<tr>
<td>ST_UID</td>
<td>User ID of the owner of the file</td>
</tr>
</tbody>
</table>

### Usage notes

All time values returned in stem are in POSIX format. Time is in seconds since 00:00:00 GMT, January 1, 1970. You can use gmtime to convert it to other forms.

### Example

In the following example, assume that *fd* was assigned a value earlier in the exec:

```
fstat (fd) st.
```
Function

fstatvfs invokes the fstatvfs callable service to obtain information about a file system identified by a file descriptor that refers to a file in that file system.

Parameters

fd  A file descriptor referring to a file from the desired file system.

stem
The name of a stem variable used to return the information. On return, stem.0 contains the number of variables returned. You can use the predefined variables beginning with STFS_ or their equivalent numeric values to access the status values they represent. (See Appendix A for the numeric values.) For example, stem.stfs_avail accesses the number of blocks available in the file system.

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>STFS_AVAIL</td>
<td>Space available to unprivileged users in block-size units.</td>
</tr>
<tr>
<td>STFS_BFREE</td>
<td>Total number of free blocks.</td>
</tr>
<tr>
<td>STFS_BLOCKSIZE</td>
<td>Block size.</td>
</tr>
<tr>
<td>STFS_FAVAIL</td>
<td>Number of free file nodes available to unprivileged users.</td>
</tr>
<tr>
<td>STFS_FFREE</td>
<td>Total number of free file nodes.</td>
</tr>
<tr>
<td>STFS_FILES</td>
<td>Total number of file nodes in the file system.</td>
</tr>
<tr>
<td>STFS_FRSIZE</td>
<td>Fundamental file system block size.</td>
</tr>
<tr>
<td>STFS_FSID</td>
<td>File system ID set by the logical file system.</td>
</tr>
<tr>
<td>STFS_INUSE</td>
<td>Allocated space in block-size units.</td>
</tr>
<tr>
<td>STFS_INVARSEC</td>
<td>Number of seconds the file system will remain unchanged.</td>
</tr>
<tr>
<td>STFS_NAMEMAX</td>
<td>Maximum length of file name.</td>
</tr>
<tr>
<td>STFS_NOSEC</td>
<td>Mount data set with no security bit.</td>
</tr>
<tr>
<td>STFS_NOSUID</td>
<td>SETUID and SETGID are not supported.</td>
</tr>
<tr>
<td>STFS_RONLY</td>
<td>File system is read-only.</td>
</tr>
<tr>
<td>STFS_TOTAL</td>
<td>Total space in block-size units.</td>
</tr>
</tbody>
</table>

Example

In the following example, assume that fd was assigned a value earlier in the exec:

"fstatvfs" fd st.
fsync

Function
fsync invokes the fsync callable service to write changes on the direct access storage device that holds the file identified by the file descriptor.

Parameters
fd The file descriptor (a number) for the file.

Usage notes
On return from a successful call, all updates have been saved on the direct access storage that holds the file.

Example
To invoke fsync for file descriptor 1:
"fsync 1"
Function

ftrunc invokes the ftrunc callable service to change the size of the file identified by the file descriptor.

Parameters

fd  The file descriptor (a number) for the file.

file_size  The new size of the file, in bytes.

Usage notes

1. The ftrunc service changes the file size to file_size bytes, beginning at the first byte of the file. If the file was previously larger than file_size, all data from file_size to the original end of the file is removed. If the file was previously shorter than file_size, bytes between the old and new lengths are read as zeros.

2. Full blocks are returned to the file system so that they can be used again, and the file size is changed to the lesser of file_size or the current length of the file.

3. The file offset is not changed.

Example

In the following example, assume that fd was assigned a value earlier in the exec.

To truncate fd to 0 bytes:

"ftrunc" fd 0
Function
getcwd invokes the getcwd callable service to get the pathname of the working directory.

Parameters
variable
  The name of the variable where the pathname of the working directory is to be stored.

Example
To get the pathname of the working directory:
"getcwd cwd"
Function
getegid invokes the getegid callable service to get the effective group ID (GID) of the calling process.

Usage notes
1. The effective GID is returned in RETVAL.
2. If the service fails, the process abends.
Function
geteuid invokes the geteuid callable service to get the effective user ID (UID) of the calling process.

Usage notes
1. The effective UID is returned in RETVAL.
2. If the service fails, the process abends.
Function

`getgid` invokes the getgid callable service to get the real group ID (GID) of the calling process.

Usage notes
1. The GID is returned in RETVAL.
2. If the service fails, the process abends.
getgrent

Function
getgrent invokes the getgrent callable service to retrieve a group database entry.

stem
The name of a stem variable used to return the information. Upon return, stem.0 contains the number of variables returned. To access the status values, you can specify a numeric value (see Appendix A) or the predefined variable beginning with GR_ used to derive the appropriate numeric value. For example, you can specify stem.2 or stem.gr_gid to access the group ID:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>GR_GID</td>
<td>The group ID</td>
</tr>
<tr>
<td>GR_MEM</td>
<td>The stem index of the first member name returned</td>
</tr>
<tr>
<td>GR_MEMBERS</td>
<td>The number of members returned</td>
</tr>
<tr>
<td>GR_NAME</td>
<td>The name of the group</td>
</tr>
</tbody>
</table>

Usage notes
1. The service is intended to be used to search the group database sequentially. The first call to this service from a given task returns the first group entry in the group database. Subsequent calls from the same task return the next group entry found, until no more entries exist, at which time a RETVAL of 0 is returned.
2. The setgrent service can be used to reset this sequential search.

Example
To list all groups in the group data base:

do forever
   "getgrent gr."
   if retval=0 | retval=-1 then
      leave
   say gr.gr_name
end
Function

getgrgid invokes the getgrgid callable service to get information about a group and its members; the group is identified by its group ID (GID).

Parameters

gid
A numeric value.

stem
The name of a stem variable used to return the information. Upon return, stem.0 contains the number of variables returned. To access the status values, you can specify a numeric value (see Appendix A) or the predefined variable beginning with GR_ used to derive the appropriate numeric value. For example, you can specify stem.2 or stem.gr_gid to access the group ID:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>GR_GID</td>
<td>The group ID</td>
</tr>
<tr>
<td>GR_MEM</td>
<td>The stem index of the first member name returned</td>
</tr>
<tr>
<td>GR_MEMBERS</td>
<td>The number of members returned</td>
</tr>
<tr>
<td>GR_NAME</td>
<td>The name of the group</td>
</tr>
</tbody>
</table>

Usage notes

A RETVAL greater than zero indicates success. A RETVAL of −1 or 0 indicates failure. A RETVAL of 0 has an ERRNOJR, but no ERRNO.

Example

In the following example, assume that gid was assigned a value earlier in the exec. To get a list of names connected with a group ID:

```
"getgrgid (gid) gr."
say 'Users connected to group number' gid ':'
do i=gr_mem to gr.0
   say gr.i
end
```
Function

getgrnam invokes the getgrnam callable service to get information about a group and its members. The group is identified by its name.

Parameters

name
A string that specifies the group name as defined to the system.

stem
The name of a stem variable used to return the information. Upon return, stem.0 contains the number of variables returned. To access the status values, you can specify a numeric value (see Appendix A) or the predefined variable beginning with GR_ used to derive the appropriate numeric value. For example, you can specify stem.2 or stem.gr_gid to access the group ID:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>GR_GID</td>
<td>The group ID</td>
</tr>
<tr>
<td>GR_MEM</td>
<td>The stem index of the first member name returned</td>
</tr>
<tr>
<td>GR_MEMBERS</td>
<td>The number of members returned</td>
</tr>
<tr>
<td>GR_NAME</td>
<td>The name of the group</td>
</tr>
</tbody>
</table>

Usage notes

1. A RETVAL greater than zero indicates success. A RETVAL of −1 or 0 indicates failure. A RETVAL of 0 has an ERRNOJR, but no ERRNO.
2. The return values point to data that can change or disappear after the next getgrnam or getgrgid call from that task. Each task manages its own storage separately. Move data to your own dynamic storage if you need it for future reference.
3. The storage is key 0 non-fetch-protected storage that is managed by z/OS UNIX System Services.

Example

To get information about the group named SYS1:

```
"getgrnam SYS1 gr."
say 'Users connected to group SYS1:'
do i= gr_mem to gr.0
   say gr.i
end
```
Function
getgroups invokes the getgroups callable service to get the number of supplementary group IDs (GIDs) for the calling process and a list of those supplementary group IDs.

Parameters
stem
The name of a stem variable used to return the information. Upon return, stem.0 contains the number of variables returned. stem.1 to stem.n (where n is the number of variables returned) each contain a group ID.

Usage notes
A RETVAL greater than zero indicates success. A RETVAL of -1 or 0 indicates failure. A RETVAL of 0 has an ERRNOJR, but no ERRNO.

Example
To invoke getgroups:
"getgroups grps."
Function

getgroupsbyname invokes the getgroupsbyname callable service to get the number of supplementary group IDs (GIDs) for a specified user name and, optionally, get a list of those supplementary group IDs.

Parameters

name
A string that specifies the name of the user as defined to the system.

stem
The name of a stem variable used to return the information. Upon return, stem.0 contains the number of variables returned. stem.1 to stem.n (where n is the number of variables returned) each contain a supplementary group ID for name.

Usage notes

A RETVAL greater than zero indicates success. A RETVAL of −1 or 0 indicates failure. A RETVAL of 0 has an ERRNOJR, but no ERRNO.

Example

To get the number of supplementary group IDs for the user MEGA:

"getgroupsbyname MEGA supgrp."
Function

`getlogin` invokes the `getlogin` callable service to get the user login name associated with the calling process.

Parameters

`variable`

The name of the variable in which the login name is returned.

Usage notes

If the service fails, the process abends.

Example

To invoke `getlogin` and store the login name in the variable `myid`:

```
getlogin myid
```
**getment**

**Function**

*getment* invokes the `w_getmntent` callable service to get information about the mounted file systems and return the information in the format used by the z/OS UNIX callable services. Alternatively, you can use the `getmntent` syscall command to format the mount entries in a stem.

**Parameters**

- **length**
  The size of the specified variable. If the length is 0, RETVAL contains the total number of mounted file systems; otherwise, RETVAL contains the number of entries returned.

- **variable**
  The name of the buffer where the information about the mount entries is to be stored. Clear the buffer on the first call and do not alter it between calls.

**Usage notes**

1. Before a program calls *getment* for the first time, the variable should be dropped, set to blanks, or set to nulls.
2. If more than one call is made to *getment*, use the same variable on each call, because part of the information returned in the variable tells the file system where to continue retrieving its information.
3. *getment* normally returns information about as many file systems as are mounted, or as many as fit in the passed variable. The number of entries contained in the variable is returned. The caller must have a variable large enough to receive information about at least a single mount entry with each call. If a zero-length variable is passed, no information is returned, but the return value contains the total number of mounted file systems. This value could then be used to get enough storage to retrieve information on all these file systems in one additional call.
4. You could also retrieve all mount entries by setting up a loop that continues to call *getment* until a return value of either −1 (in an error) or 0 (no more entries found) is returned.

**Example**

In the following example, assume that `buf` was assigned a value earlier in the exec. This example returns the number of mounted file systems in RETVAL:

```
"getment 0 buf"
```
Function
getmntent invokes the w_getmntent callable service to get information about the mounted file systems, or a specific mounted file system, and return the formatted information in a stem variable.

Parameters
stem
The name of a stem variable used to return the mount table information. The stem has two dimensions: the field name, followed by a period and the number of the mount table entry. For example, you can access the file system name for the first entry in the mount table as stem.mnte_fsname.1.

For the field name, you can specify a numeric value (see Appendix A) or the predefined variable beginning with MNTE_ used to derive the appropriate numeric value. For example, you could use stem.mnte_fsname.1 or stem.6.1 to access the file system name for the first entry:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>MNTE_AGNAME</td>
<td>The name of the zFS aggregate data set.</td>
</tr>
<tr>
<td>MNTE_BYTESREADHW</td>
<td>The number of bytes read (high-word value).</td>
</tr>
<tr>
<td>MNTE_BYTESREADLW</td>
<td>The number of bytes read (low-word value).</td>
</tr>
<tr>
<td>MNTE_BYTESWRITTENHW</td>
<td>The number of bytes written (high-word value).</td>
</tr>
<tr>
<td>MNTE_BYTESWRITTENLW</td>
<td>The number of bytes written (low-word value).</td>
</tr>
<tr>
<td>MNTE_DD</td>
<td>The ddname specified on the mount.</td>
</tr>
<tr>
<td>MNTE_DEV</td>
<td>The device ID of the file system.</td>
</tr>
<tr>
<td>MNTE_DIRIBC</td>
<td>The number of directory I/O blocks.</td>
</tr>
<tr>
<td>MNTE_FILETAG</td>
<td>The file tag.</td>
</tr>
<tr>
<td>MNTE_FROMSYS</td>
<td>The file systems are to be moved from here.</td>
</tr>
<tr>
<td>MNTE_FSNAME</td>
<td>The file system name that was specified on the mount.</td>
</tr>
<tr>
<td>MNTE_FSTYPE</td>
<td>The file system type; for example, HFS.</td>
</tr>
</tbody>
</table>

This statement applies only to HFS and zFS. The file system name is the name of the MVS data set that contains the HFS or zFS file system.
getmntent

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>MNTE_MODE</td>
<td>The file system type mount method. MNTE_MODE contains a decimal number that corresponds to the value as defined in the field mntentfsmode in the BPXYMNTE mapping macro. The REXX predefined variables for this field, as well as other fields on this interface, are a subset of what might be returned by the getmntent service. You can specify a numeric value (see Appendix A) or one of the following predefined variables used to derive the appropriate numeric value. For example, there is no predefined rexx variable corresponding to the bit position defined by mntentfssynchonly in the mntentfsmode.</td>
</tr>
<tr>
<td>MNT_MODE_RDWR</td>
<td>File system mounted read-write.</td>
</tr>
<tr>
<td>MNT_MODE_RDONLY</td>
<td>File system mounted read-only.</td>
</tr>
<tr>
<td>MNT_MODE_AUNMOUNT</td>
<td>The file system can be unmounted if the system’s owner crashes.</td>
</tr>
<tr>
<td>MNT_MODE_CLIENT</td>
<td>The file system is a client.</td>
</tr>
<tr>
<td>MNT_MODE_EXPORT</td>
<td>The file system exported by DFS™.</td>
</tr>
<tr>
<td>MNT_MODE_NOAUTOMOVE</td>
<td>Automove is not allowed.</td>
</tr>
<tr>
<td>MNT_MODE_NOSEC</td>
<td>No security checks are enforced.</td>
</tr>
<tr>
<td>MNT_MODE_NOSETID</td>
<td>SetUID is not permitted for files in this filesystem.</td>
</tr>
<tr>
<td>MNT_MODE_SECACL</td>
<td>ACLs are supported by the security product.</td>
</tr>
<tr>
<td>MNTE_PARDEV</td>
<td>The ST_DEV of the parent file system.</td>
</tr>
<tr>
<td>MNTE_PARM</td>
<td>The parameter specified with mount().</td>
</tr>
<tr>
<td>MNTE_PATH</td>
<td>The mountpoint pathname.</td>
</tr>
<tr>
<td>MNTE_PFSSTATUSNORMAL</td>
<td>Normal status string returned by mount().</td>
</tr>
<tr>
<td>MNTE_PFSSTATUSEXCP</td>
<td>Exception status string returned by the physical file system.</td>
</tr>
<tr>
<td>MNTE_ROSECLABEL</td>
<td>Default security label for objects in a read-only file system.</td>
</tr>
<tr>
<td>MNTE_QJOBNAME</td>
<td>The job name of the quiesce requestor.</td>
</tr>
<tr>
<td>MNTE_QPID</td>
<td>The PID of the quiesce requestor.</td>
</tr>
<tr>
<td>MNTE_QSYSNAME</td>
<td>The name of the quiesce system name.</td>
</tr>
<tr>
<td>MNTE_READCT</td>
<td>The number of reads from filesys.</td>
</tr>
<tr>
<td>MNTE_READIBC</td>
<td>The number of read I/O blocks.</td>
</tr>
<tr>
<td>MNTE_FLAGS</td>
<td>The request flags.</td>
</tr>
<tr>
<td>MNTE_ROOTINO</td>
<td>The_inode of the mountpoint.</td>
</tr>
<tr>
<td>MNTE_ROSECLABEL</td>
<td>The read-only SECLABEL.</td>
</tr>
</tbody>
</table>
## Variable

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>MNTE_STATUS</td>
<td>The status of the file system. To specify the information, you can specify a numeric value (see Appendix A) or one of the following predefined variables used to derive the appropriate numeric value:</td>
</tr>
<tr>
<td>MNT_ASYNCHMOUNT</td>
<td>Asynchronous mount in progress for this file system.</td>
</tr>
<tr>
<td>MNT_FILEACTIVE</td>
<td>File system is active.</td>
</tr>
<tr>
<td>MNT_FILEDEAD</td>
<td>File system is dead.</td>
</tr>
<tr>
<td>MNT_FILEDRAIN</td>
<td>File system is being unmounted with the drain option.</td>
</tr>
<tr>
<td>MNT_FILEFORCE</td>
<td>File system is being unmounted with the force option.</td>
</tr>
<tr>
<td>MNT_FILEIMMED</td>
<td>File system is being unmounted with the immediate option.</td>
</tr>
<tr>
<td>MNT_FILENORM</td>
<td>File system is being unmounted with the normal option.</td>
</tr>
<tr>
<td>MNT_FILERESET</td>
<td>File system is being reset.</td>
</tr>
<tr>
<td>MNT_IMMEDTRIED</td>
<td>File system unmount with the immediate option failed.</td>
</tr>
<tr>
<td>MNT_MOUNTINPROGRESS</td>
<td>Mount in progress for this file system.</td>
</tr>
<tr>
<td>MNT_QUIESCED</td>
<td>File system is quiesced.</td>
</tr>
<tr>
<td>MNTE_STATUS2</td>
<td>The status of the file system.</td>
</tr>
<tr>
<td>MNTE_SUCCESS</td>
<td>Successful moves.</td>
</tr>
<tr>
<td>MNTE_SYSLIST</td>
<td>A list of system names.</td>
</tr>
<tr>
<td>MNTE_SYSNAME</td>
<td>The name of the owning system.</td>
</tr>
<tr>
<td>MNTE_TYPE</td>
<td>The file system type.</td>
</tr>
<tr>
<td>MNTE_WRITEECT</td>
<td>The number of writes done.</td>
</tr>
<tr>
<td>MNTE_WRITEIBC</td>
<td>The number of write I/O blocks.</td>
</tr>
</tbody>
</table>

### devno

An optional parameter, this is the device number for a specific file system for which you want the mount information. Specifying 0 is the equivalent of not specifying a device number.

### Example

Show all mounted file systems:

```addresses
address syscall 'getmntent mounts.'
do i=1 to mounts.0
   say mounts.mnte_fsname.i
say endl
```

---

**getmntent**

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Function

getpgrp invokes the getpgrp callable service to get the process group ID (PGID) of the calling process.

Usage notes
1. The PGID for the calling process is returned in RETVAL.
2. If the service fails, the process abends.
getpid

Function
getpid invokes the getpid callable service to get the process ID (PID) of the calling process.

Usage notes
1. The PID for the calling process is returned in RETVAL.
2. If the service fails, the process abends.
getppid

Function
getppid invokes the getppid callable service to get the parent process ID (PPID) of the calling process.

Usage notes
1. The parent PID for the calling process is returned in RETVAL.
2. If the service fails, the process abends.
Function

getpsent invokes the w_getpsent callable service to provide data describing the status of all the processes to which you are authorized. The data is formatted in a stem. The PS_ variables, or their numeric equivalents, are used to access the fields.

Parameters

stem

Upon return, stem.0 contains the number of processes for which information is returned. stem.1 through stem.n (where n is the number of entries returned) each contain process information for the nth process.

You can use the predefined variables that begin with PS_ or their equivalent numeric values (see Appendix A) to access the information. For example, stem.1.ps_pid is the process ID for the first process returned.

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>PS_CMD</td>
<td>Command</td>
</tr>
<tr>
<td>PS_CONTTY</td>
<td>Controlling tty</td>
</tr>
<tr>
<td>PS_EGID</td>
<td>Effective group ID</td>
</tr>
<tr>
<td>PS_EUID</td>
<td>Effective user ID</td>
</tr>
<tr>
<td>PS_FGPID</td>
<td>Foreground process group ID</td>
</tr>
<tr>
<td>PS_MAXVNODES</td>
<td>Maximum number of vnode tokens allowed</td>
</tr>
<tr>
<td>PS_PATH</td>
<td>Pathname</td>
</tr>
<tr>
<td>PS_PGPID</td>
<td>Process Group ID</td>
</tr>
<tr>
<td>PS_PID</td>
<td>Process ID</td>
</tr>
<tr>
<td>PS_PPID</td>
<td>Parent Process ID</td>
</tr>
<tr>
<td>PS_RGID</td>
<td>Real group ID</td>
</tr>
<tr>
<td>PS_RUID</td>
<td>Real user ID</td>
</tr>
<tr>
<td>PS_SERVERFLAGS</td>
<td>Server flags</td>
</tr>
<tr>
<td>PS_SERVERNAME</td>
<td>Server name supplied on registration</td>
</tr>
<tr>
<td>PS_SERVERTYPE</td>
<td>Server type (File=1; Lock=2)</td>
</tr>
<tr>
<td>PS_SGID</td>
<td>Saved set group ID</td>
</tr>
<tr>
<td>PS_SID</td>
<td>Session ID (leader)</td>
</tr>
<tr>
<td>PS_SIZE</td>
<td>Total size</td>
</tr>
<tr>
<td>PS_STARTTIME</td>
<td>Starting time, in POSIX format (seconds since the Epoch, 00:00:00 on 1 January 1970)</td>
</tr>
<tr>
<td>PS_STAT</td>
<td>Process status</td>
</tr>
</tbody>
</table>
### Variable Description

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>PS_STATE</td>
<td>Process state This value can be expressed as one of the following predefined variables or as an alphabetic value (see <a href="#">Appendix A</a>).</td>
</tr>
<tr>
<td>PS_CHILD</td>
<td>Waiting for a child process</td>
</tr>
<tr>
<td>PS_FORK</td>
<td>fork() a new process</td>
</tr>
<tr>
<td>PS_FREEZE</td>
<td>QUIESCEFREEZE</td>
</tr>
<tr>
<td>PS_MSGRCV</td>
<td>IPC MSGRCV WAIT</td>
</tr>
<tr>
<td>PS_MSGSND</td>
<td>IPC MSGSND WAIT</td>
</tr>
<tr>
<td>PS_PAUSE</td>
<td>MVSPAUSE</td>
</tr>
<tr>
<td>PS_QUIESCE</td>
<td>Quiesce termination wait</td>
</tr>
<tr>
<td>PS_RUN</td>
<td>Running, not in kernel wait</td>
</tr>
<tr>
<td>PS_SEMWT</td>
<td>IPC SEMOP WAIT</td>
</tr>
<tr>
<td>PS_SLEEP</td>
<td>sleep() issued</td>
</tr>
<tr>
<td>PS_WAITC</td>
<td>Communication kernel wait</td>
</tr>
<tr>
<td>PS_WAITF</td>
<td>File system kernel wait</td>
</tr>
<tr>
<td>PS_WAITO</td>
<td>Other kernel wait</td>
</tr>
<tr>
<td>PS_ZOMBIE</td>
<td>Process cancelled</td>
</tr>
<tr>
<td>PS_ZOMBIE2</td>
<td>Process terminated yet still the session or process group leader</td>
</tr>
<tr>
<td>PS_SUID</td>
<td>Saved set user ID</td>
</tr>
<tr>
<td>PS_SYSTIME</td>
<td>System CPU time, a value of the type clock_t, which needs to be divided by sysconf(_SC_CLK_TCK) to convert it to seconds. For z/OS UNIX, this value is expressed in hundredths of a second.</td>
</tr>
<tr>
<td>PS_USERTIME</td>
<td>User CPU time, a value of the type clock_t, which needs to be divided by sysconf(_SC_CLK_TCK) to convert it to seconds. For z/OS UNIX, this value is expressed in hundredths of a second.</td>
</tr>
<tr>
<td>PS_VNODECOUNT</td>
<td>Current number of vnode tokens</td>
</tr>
</tbody>
</table>

### Usage notes

1. Information is returned for only those processes for which RACF allows the user access based on effective user ID, real user ID, or saved set user ID.
2. PS_STARTTIME is in seconds since the Epoch (00:00:00 on 1 January 1970).
3. PS_USERTIME and PS_SYSTIME are task-elapsed times in 1/100ths of seconds.
4. PS_SYSTIME reports the system CPU time consumed for the address space that the process is running in. When only one process is running in the address space, this time represents the accumulated system CPU time for that process. However, when more than one process is running in an address space, the information that is returned is actually the accumulated system CPU time consumed by all of the work running in the address space.

### Example

This exec will produce output similar to the `ps -A` shell command, displaying information on all accessible processes:

```rexx
/* rexx */
address syscall
say right('PID',12) left('TTY',10) ' TIME' 'COMMAND'
```

```
getpsent
```

```
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```
getpsent

'getpsent ps.' /* get process data */
do i=1 to ps.0 /* process each entry returned */
t=(ps.i.ps_usertime + 50) % 100 /* change time to seconds */
'gmtime (t) gm.' /* convert to usable format */
if gm.tm_hour=0 then /* set hours: samp ignores day */
  h=' '
else
  h=right(gm.tm_hour,2,0):'

m=right(gm.tm_min,2,0):'

parse value reverse(ps.i.ps_contty),
  with tty '/' /* get tty filename */
tty=reverse(tty)
say right(ps.i.ps_pid,12), /* display process id */
say right(ps.i.ps_pid,12), /* display process id */
  left(tty,10), /* display controlling tty */
  h || m || right(gm.tm_sec,2,0), /* display process time */
  ps.i.ps_cmd /* display command */
end

return 0
Function
getpwent invokes the getpwent callable service to retrieve a user database entry.

Parameters
stem
The name of a stem variable used to return the information. Upon return, stem.0 contains the number of variables returned. You can use a numeric value (see Appendix A) or the predefined variables beginning with PW_ to access the values:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>PW_DIR</td>
<td>The initial working directory</td>
</tr>
<tr>
<td>PW_GID</td>
<td>The group ID</td>
</tr>
<tr>
<td>PW_NAME</td>
<td>The TSO/E user ID</td>
</tr>
<tr>
<td>PW_SHELL</td>
<td>The name of the initial user program.</td>
</tr>
<tr>
<td>PW_UID</td>
<td>The user ID (UID) as defined to RACF.</td>
</tr>
</tbody>
</table>

Example
To list all users in the user database:

do forever
  "getpwent pw."
  if retval=0 | retval=-1 then
    leave
  say pw pw_name
end
Function
getpwnam invokes the getpwnam callable service to get information about a user, identified by user name.

Parameters

name
The user name as defined to the system.

stem
The name of a stem variable used to return the information. Upon return, stem.0 contains the number of variables returned. To access the information, you can specify a numeric value (see Appendix A) or the predefined variable beginning with PW_ used to derive the appropriate numeric value. For example, to access the name of the user's initial working directory, you can specify stem.4 or stem.pw_dir.

<table>
<thead>
<tr>
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</tr>
</thead>
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<tr>
<td>PW_DIR</td>
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</tr>
<tr>
<td>PW_GID</td>
<td>The group ID</td>
</tr>
<tr>
<td>PW_NAME</td>
<td>The TSO/E user ID</td>
</tr>
<tr>
<td>PW_SHELL</td>
<td>The name of the initial user program.</td>
</tr>
<tr>
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<td>The user ID (UID) as defined to RACF.</td>
</tr>
</tbody>
</table>

Usage notes
1. A RETVAL greater than zero indicates success. A RETVAL of -1 or 0 indicates failure. A RETVAL of 0 has an ERRNOJR, but no ERRNO.
2. If an entry for the specified name is not found in the user database, the RETVAL is 0.

Example
To get information about the user JANET:
"getpwnam JANET pw."
Function
getpwuid invokes the getpwuid callable service to get information about a user, identified by UID.

Parameters
uid
A numeric value that is the user’s UID as defined to the system.

stem
The name of a stem variable used to return the information. Upon return, stem.0 contains the number of variables returned. You can use a numeric value (see Appendix A) or the predefined variables beginning with PW_ to access the values:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>PW_DIR</td>
<td>The initial working directory</td>
</tr>
<tr>
<td>PW_GID</td>
<td>The group ID</td>
</tr>
<tr>
<td>PW_NAME</td>
<td>The TSO/E user ID</td>
</tr>
<tr>
<td>PW_SHELL</td>
<td>The name of the initial user program.</td>
</tr>
<tr>
<td>PW_UID</td>
<td>The user ID (UID) as defined to RACF.</td>
</tr>
</tbody>
</table>

Usage notes
A RETVAL greater than zero indicates success. A RETVAL of −1 or 0 indicates failure. A RETVAL of 0 has an ERRNOJR, but no ERRNO.

Example
To get information about the user with UID 42:

"getpwuid 42 pw."
Function
getrlimit invokes the getrlimit callable service to get the maximum and current resource limits for the calling process.

Parameters
resource
The resource whose limit is being requested. Resources that have limits with values greater than RLIM_INFINITY will return values of RLIM_INFINITY.

You can use the predefined variables beginning with RLIMIT_ or their equivalent numeric values to access the limits. (See Appendix A for the numeric values.)

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>RLIMIT_AS</td>
<td>Maximum address space size for a process</td>
</tr>
<tr>
<td>RLIMIT_CORE</td>
<td>Maximum size (in bytes) of a core dump created by a process</td>
</tr>
<tr>
<td>RLIMIT_CPU</td>
<td>Maximum amount of CPU time (in seconds) used by a process</td>
</tr>
<tr>
<td>RLIMITFSIZE</td>
<td>Maximum files size (in bytes) created by a process</td>
</tr>
<tr>
<td>RLIMIT_NOFILE</td>
<td>Maximum number of open file descriptors for a process</td>
</tr>
</tbody>
</table>

stem
The name of the stem variable used to return the limit. stem.1 is the first word and stem.2 is the second word. The first word contains the current limit; the second word contains the maximum limit. The values for each word are dependent on the resource specified.

Example
To print the maximum and current limit for number of open files allowed:

```
"getrlimit" rlimit_nofile r.
say 'maximum open limit is' r.2
say 'current open limit is' r.1
```
getuid

Function
getuid invokes the getuid callable service to get the real user ID of the calling process.

Usage notes
Upon return, RETVAL contains the UID. If the service fails, the process abends.
Function

`gmtime` converts time expressed in seconds since Epoch into month, day, and year time format.

Parameters

time

A numeric value, the time expressed as “POSIX time,” the number of seconds since the Epoch (00:00:00 on 1 January 1970).

stem

The name of a stem variable used to return the information. Upon return, `stem.0` contains the number of variables returned. To access the status values, you can specify a numeric value (see Appendix A) or the predefined variable beginning with TM__ used to derive the appropriate numeric value. For example, to access the year, you can specify `stem.6` or `stem.tm_year`.

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>TM_HOUR</td>
<td>The hour of the day.</td>
</tr>
<tr>
<td>TM_ISDST</td>
<td>The daylight saving time flag. This flag is always zero (−1) for Greenwich Mean Time (GMT).</td>
</tr>
<tr>
<td>TM_MDAY</td>
<td>The day of the month, 1 to 31.</td>
</tr>
<tr>
<td>TM_MIN</td>
<td>The minutes after the hour, 0 to 59.</td>
</tr>
<tr>
<td>TM_MON</td>
<td>The months since January, 0 to 11.</td>
</tr>
<tr>
<td>TM_SEC</td>
<td>The seconds after the minute, 0 to 59.</td>
</tr>
<tr>
<td>TM_WDAY</td>
<td>The days since Sunday, 0 to 6.</td>
</tr>
<tr>
<td>TM_YDAY</td>
<td>The days since January 1, 0 to 365.</td>
</tr>
<tr>
<td>TM_YEAR</td>
<td>The year.</td>
</tr>
</tbody>
</table>

Example

The `st_ctime` in the following example was set with a stat call:

"gmtime" st.st_ctime "tm."
Function
ioctl invokes the w_ioctl service to issue ioctl commands to a file.

Restriction: Hierarchical file system (HFS) files and pipes do not support ioctl commands.

Parameters
fd  A file descriptor for an open file.
   Restriction: REXX does not support ioctl commands for sockets.

command
   The numeric value for an ioctl command. The commands that can be specified vary by device and are defined by the device driver.

buffer
   The name of a buffer containing the argument to be passed to the device driver.
The argument is limited to 1024 bytes.

length
   An optional numeric value indicating the length of the argument. If length is not specified, the length of the buffer is used.
isatty

Function
isatty invokes the isatty callable service to determine if a file is a terminal.

Parameters
fd The file descriptor (a number) for the file.

Usage notes
On return, RETVAL contains either 0 (not a tty) or 1 (a tty).

Example
To test if file descriptor 0 is a terminal:
"isatty 0"
Function

The `kill` callable service sends a signal to a process or process group. It invokes the `kill` callable service to send a signal to a process or process group.

Parameters

**pid**
A number, the process ID of the process or process group to which the caller wants to send a signal.

**signal**
The signal to be sent. You can set the value by using a numeric value (see Appendix A) or the predefined variable beginning with SIG used to derive the numeric value:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>SIGABND</td>
<td>Abend</td>
</tr>
<tr>
<td>SIGABRT</td>
<td>Abnormal termination</td>
</tr>
<tr>
<td>SIGALRM</td>
<td>Timeout</td>
</tr>
<tr>
<td>SIGBUS</td>
<td>Bus error</td>
</tr>
<tr>
<td>SIGCHLD</td>
<td>Child process terminated or stopped</td>
</tr>
<tr>
<td>SIGCONT</td>
<td>Continue if stopped</td>
</tr>
<tr>
<td>SIGDCE</td>
<td>Exclusive use by DCE</td>
</tr>
<tr>
<td>SIGFPE</td>
<td>Erroneous arithmetic operation, such as division by zero or an operation resulting in an overflow</td>
</tr>
<tr>
<td>SIGHUP</td>
<td>Hangup detected on the controlling terminal</td>
</tr>
<tr>
<td>SIGILL</td>
<td>Termination (cannot be caught or ignored)</td>
</tr>
<tr>
<td>SIGINT</td>
<td>Interactive attention</td>
</tr>
<tr>
<td>SIGIO</td>
<td>Completion of input or output</td>
</tr>
<tr>
<td>SIGIOERR</td>
<td>Error on input/output; used by the C runtime library</td>
</tr>
<tr>
<td>SIGKILL</td>
<td>Termination that cannot be caught or ignored</td>
</tr>
<tr>
<td>SIGPIPE</td>
<td>Write on a pipe with no readers</td>
</tr>
<tr>
<td>SIGPOLL</td>
<td>Pollable event</td>
</tr>
<tr>
<td>SIGPROF</td>
<td>Profiling timer expired</td>
</tr>
<tr>
<td>SIGQUIT</td>
<td>Interactive termination</td>
</tr>
<tr>
<td>SIGSEGV</td>
<td>Detection of an incorrect memory reference</td>
</tr>
<tr>
<td>SIGSTOP</td>
<td>Stop that cannot be caught or ignored</td>
</tr>
<tr>
<td>SIGSYS</td>
<td>Bad system call</td>
</tr>
<tr>
<td>SIGTERM</td>
<td>Termination</td>
</tr>
<tr>
<td>SIGTRAP</td>
<td>Trap used by the ptrace callable service</td>
</tr>
<tr>
<td>SIGTSTP</td>
<td>Interactive stop</td>
</tr>
</tbody>
</table>
### Variable Description

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>SIGTIN</td>
<td>Read from a controlling terminal attempted by a member of a background process group</td>
</tr>
<tr>
<td>SIGTTOU</td>
<td>Write from a controlling terminal attempted by a member of a background process group</td>
</tr>
<tr>
<td>SIGURG</td>
<td>High bandwidth data is available at a socket</td>
</tr>
<tr>
<td>SIGUSR1</td>
<td>Reserved as application-defined signal 1</td>
</tr>
<tr>
<td>SIGUSR2</td>
<td>Reserved as application-defined signal 2</td>
</tr>
<tr>
<td>SIGVTALRM</td>
<td>Virtual timer expired</td>
</tr>
<tr>
<td>SIGXCPU</td>
<td>CPU time limit exceeded</td>
</tr>
<tr>
<td>SIGXFSZ</td>
<td>File size limit exceeded</td>
</tr>
</tbody>
</table>

**option**

The name of a variable that contains exactly four bytes mapped by the BPXYPSSD macro.

**Usage notes**

1. A caller can send a signal if the real or effective user ID of the caller is the same as the real or saved set user ID of the intended recipient. A caller can also send signals if it has appropriate privileges.
2. Regardless of user ID, a caller can always send a **SIGCONT** signal to a process that is a member of the same session as the sender.
3. A caller can also send a signal to itself. If the signal is not blocked, at least one pending unblocked signal is delivered to the sender before the service returns control. Provided that no other unblocked signals are pending, the signal delivered is the signal sent.

See "Using the REXX signal services" on page 9 for information on using signal services.

**Example**

In the following example, assume that *pid* was assigned a value earlier in the exec:

```
"kill" pid sighup
```
Function

`lchown` invokes the `lchown` callable service to change the owner or group for a file, directory, or symbolic link.

Parameters

`pathname`  
The pathname for a file, directory, or symbolic link.

`uid`  
The numeric UID for the new owner or the present UID, or −1 if there is no change to the UID.

`gid`  
The numeric GID for the new group or the present GID, or −1 if there is no change to the GID.

Usage notes

1. If `lchown`'s target is a symbolic link, it modifies the ownership of the actual symbolic link file instead of the ownership of the file pointed to by the symbolic link.
2. The `lchown` service changes the owner UID and owner GID of a file. Only a superuser can change the owner UID of a file.
3. The owner GID of a file can be changed by a superuser, or if a caller meets all of these conditions:
   • The effective UID of the caller matches the file’s owner UID.
   • The `uid` value specified in the change request matches the file’s owner UID.
   • The `gid` value specified in the change request is the effective GID, or one of the supplementary GIDs, of the caller.
4. The set-user-ID-on-execution and set-group-ID-on-execution permissions of the file mode are automatically turned off.
5. If the change request is successful, the change time for the file is updated.
6. Values for both `uid` and `gid` must be specified as they are to be set. If you want to change only one of these values, the other must be set to its present value to remain unchanged.

Example

In the following example, assume that `pathname`, `uid`, and `gid` were assigned a value earlier in the exec:

"lchown (pathname) (uid) (gid)"
Function
link invokes the link callable service to create a hard link to a file (not a directory); that is, it creates a new name for an existing file. The new name does not replace the old one.

Parameters
old_path
A pathname, the current name for the file.
new_path
A pathname, the new name for the file.

Usage notes
1. The link service creates a link named new_path to an existing file named old_path. This provides an alternate pathname for the existing file, so that the file can be accessed by the old name or the new name. The link can be stored in the same directory as the original file, or in a different directory.
2. The link and the file must be in the same file system.
3. If the link is created successfully, the service increments the link count of the file. The link count shows how many links exist for a file. If the link is not created successfully, the link count is not incremented.
4. Links are allowed only to files, not to directories.
5. If the link is created successfully, the change time of the linked-to file is updated and the change and modification times of the directory that holds the link are updated.

Example
To create the hard link /usr/bin/grep to the file /bin/grep:
"link /bin/grep /usr/bin/grep"
Function

`lseek` invokes the `lseek` callable service to change the file offset of a file to a new position. The file offset is the position in a file from which data is next read or to which data is next written.

Parameters

- **fd**: The file descriptor (a number) for the file whose offset you want to change. The file descriptor is returned when the file is opened.
- **position**: A number indicating the number of bytes by which you want to change the offset. If the number is unsigned, the offset is moved forward that number of bytes; if the number is preceded by a – (minus sign), the offset is moved backward that number of bytes.
- **whence**: A numeric value that indicates the point from which the offset is calculated. You can specify a numeric value (see Appendix A) or the predefined variable beginning with SEEK_ used to derive the appropriate numeric value:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>SEEK_CUR</td>
<td>Set the file offset to current offset plus the specified offset.</td>
</tr>
<tr>
<td>SEEK_END</td>
<td>Set the file offset to EOF plus the specified offset.</td>
</tr>
<tr>
<td>SEEK_SET</td>
<td>Set the file offset to the specified offset.</td>
</tr>
</tbody>
</table>

Usage notes

1. **position** gives the length and direction of the offset change. **whence** states where the change is to start. For example, assume that a file is 2000 bytes long, and that the current file offset is 1000:

<table>
<thead>
<tr>
<th>Position Specified</th>
<th>Whence</th>
<th>New File Offset</th>
</tr>
</thead>
<tbody>
<tr>
<td>80</td>
<td>SEEK_CUR</td>
<td>1080</td>
</tr>
<tr>
<td>1200</td>
<td>SEEK_SET</td>
<td>1200</td>
</tr>
<tr>
<td>–80</td>
<td>SEEK_END</td>
<td>1920</td>
</tr>
<tr>
<td>132</td>
<td>SEEK_END</td>
<td>2132</td>
</tr>
</tbody>
</table>

2. The file offset can be moved beyond the end of the file. If data is written at the new file offset, there will be a gap between the old end of the file and the start of the new data. A request to read data from anywhere within that gap completes successfully, and returns bytes with the value of zero in the buffer and the actual number of bytes read.

Seeking alone, however, does not extend the file. Only if data is written at the new offset does the length of the file change.

Example

To change the offset of file descriptor **fd** (assuming that it was assigned a value earlier in the `exec`) to the beginning of the file (offset 0):

```
lseek fd 0 seek_set
```
Function

**lstat** invokes the lstat callable service to obtain status information about a file.

**Parameters**

**pathname**

A pathname for the file.

**stem**

The name of a stem variable used to return the information. Upon return, *stem.0* contains the number of variables returned. To obtain the desired status information, you can use a numeric value (see Appendix A) or the predefined variables beginning with ST_ used to derive the numeric value. See "lstat" on page 83 for a list of those variables.

**Usage notes**

1. If the pathname specified is a symbolic link, the status information returned relates to the symbolic link, rather than to the file to which the symbolic link refers.
2. All time values returned are in POSIX format. Time is in seconds since 00:00:00 GMT, January 1, 1970. You can use *gmtime* to convert it to other forms.

**Example**

In the following example, assume that *pathname* was assigned a value earlier in the exec:

"lstat (pathname) st."
mkdir

Format

```
mkdir pathname mode
```

Purpose

`mkdir` invokes the `mkdir` callable service to create a new, empty directory.

Parameters

`pathname`

A pathname for the directory.

`mode`

A three- or four-digit number, corresponding to the access permission bits. Each digit must be in the range 0–7, and at least three digits must be specified. For more information on permissions, see Appendix B.

Usage

1. The file permission bits specified in `mode` are modified by the file creation mask of the calling process (see "umask" on page 188), and are then used to set the file permission bits of the new directory.
2. The new directory’s owner ID is set to the effective user ID (UID) of the calling process.
3. The `mkdir` service sets the access, change, and modification times for the new directory. It also sets the change and modification times for the directory that contains the new directory.

Examples

In the following example, assume that `pathname` and `mode` were assigned a value earlier in the exec:

```
"mkdir (pathname)" mode
```
mknod

Format

```
  mknod pathname mode
```

Purpose

*mknod* invokes the mknod callable service to create a new FIFO special file.

Parameters

**pathname**
A pathname for the file.

**mode**
A three- or four-digit number, corresponding to the access permission bits. Each digit must be in the range 0–7, and at least three digits must be specified. For more information on permissions, see [Appendix B](#).

Usage

1. The file permission bits specified in *mode* are modified by the process's file creation mask (see "umask" on page 188), and then used to set the file permission bits of the file being created.
2. The file's owner ID is set to the process's effective user ID (UID). The group ID is set to the group ID (GID) of the directory containing the file.
3. The mknod service sets the access, change, and modification times for the new file. It also sets the change and modification times for the directory that contains the new file.

Examples

In the following example, assume that *pathname* was assigned a value earlier in the exec. The mode 777 grants read-write-execute permission to everyone.

```
mknod (pathname) 777
```
mknod

Format

```
mknod pathname mode major minor
```

Purpose

mknod invokes the mknod callable service to create a new character special file. You must be a superuser to use this function.

Parameters

pathname

A pathname for the file.

mode

A three- or four-digit number, corresponding to the access permission bits. Each digit must be in the range 0–7, and at least three digits must be specified. For more information on permissions, see Appendix B.

major

The device major number corresponds to a device driver supporting a class of devices (for example, interactive terminals). For information on specifying the device major number, see Creating special files in z/OS UNIX System Services Planning.

minor

The number that corresponds to a specific device within the class of devices referred to by the device major number. For information on specifying the device minor number, see Creating special files in z/OS UNIX System Services Planning.

Usage

1. The file permission bits specified in mode are modified by the process's file creation mask (see umask on page 188), and then used to set the file permission bits of the file being created.
2. The file's owner ID is set to the process's effective user ID (UID). The group ID is set to the group ID (GID) of the directory containing the file.
3. The mknod service sets the access, change, and modification times for the new file. It also sets the change and modification times for the directory that contains the new file.

Examples

To create /dev/null with read-write-execute permission for everyone:

```
"mknod /dev/null 777 4 0"
```
mount

Format

```
mount pathname name type flags parm
```

Purpose

`mount` invokes the mount callable service to mount a file system, making the files in it available for use. You must be a superuser to use this function.

When used in a sysplex, `mount` may also be used to change some of the mounted file system attributes, including the system that owns that mount.

Parameters

**pathname**

The pathname for the mount point.

**name**

The name of the file system to be mounted. You must specify HFS data set names as fully qualified names in uppercase letters. Do not enclose the data set name in single quotes.

**type**

The type of file system, as defined by the FILESYSTYPE parameter on the BPXPRMxx parmlib member (for example, HFS). Specify this as it is specified on the parmlib member.

**flags**

A value indicating how the file system is to be mounted. To specify the information, you can specify a numeric value (see Appendix A), or one or more of the predefined variables beginning with MTM_. If more than one variable is specified, they are added together like open flags. (RDONLY and RDWR cannot be specified together.) The predefined variables used to derive the appropriate numeric value are:

- **MTM_NOSECURITY**
  Mount with no security.
- **MTM_NOSUID**
  The SETUID and SETGID mode bits on any executable in this file system are to be ignored when the program is run.
- **MTM_RDONLY**
  Mount read-only.
- **MTM_RDWR**
  Mount read-write.
- **MTM_SYNCHONLY**
  Mount must be completed synchronously; that is, `mount()` must not return +1.
**parm**
The name of a variable that contains a parameter string to be passed to the physical file system. The format and content of the string are specified by the physical file system that is to perform the logical mount.

For an HFS file system, *parm* is not used.

**stem**
The name of a stem variable which contains the mount variables. To set the mount variables, you can use a numeric value (see Appendix A) or the predefined variables beginning with MNTE_ used to derive the numeric value. Unused stem variables should be set to the null string.

The following variables are used for mount requests:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>MNTE_FILETAG</td>
<td>4-byte file tag, the contents of which are mapped by the ST_FILETAG structure in the BPXSTAT mapping macro. BPXSTAT is described in BPXSTAT — Map the Response Structure for stat in z/OS UNIX System Services Programming: Assembler Callable Services Reference.</td>
</tr>
<tr>
<td>MNTE_FSTYPE</td>
<td>The name of the HFS data set containing the file system.</td>
</tr>
<tr>
<td>MNTE_MODE</td>
<td>The file system type mount method. You can specify a numeric value (see Appendix A) or one of the following predefined variables used to derive the appropriate numeric value: MNT_MODE_RDWR - File system mounted read-write.  MNT_MODE_RDONLY - File system mounted read-only.  MNT_MODE_AUNMOUNT - The file system can be unmounted if the system’s owner crashes.  MNT_MODE_NOAUTOMOVE - Automove is not allowed.  MNT_MODE_NOSEC - No security checks are enforced.  MNT_MODE_NOSETID - SetUID is not permitted for files in this filesystem.</td>
</tr>
<tr>
<td>MNTE_PARM</td>
<td>The parameter specified with mount().</td>
</tr>
<tr>
<td>MNTE_PATH</td>
<td>The mountpoint pathname.</td>
</tr>
<tr>
<td>MNTE_SYSLIST</td>
<td>A list of system names. For more information about MNTE_SYSLIST, see BPXYMNTE — Map response and element structure of w_getmntent in z/OS UNIX System Services Programming: Assembler Callable Services Reference.</td>
</tr>
<tr>
<td>MNTE_SYSNAME</td>
<td>The name of the file system to be mounted on.</td>
</tr>
</tbody>
</table>

The following variables are used for changing mount attributes:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>MNTE_FSTYPE</td>
<td>The name of the HFS data set containing the file system.</td>
</tr>
<tr>
<td>Variable</td>
<td>Description</td>
</tr>
<tr>
<td>------------------</td>
<td>-------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------------</td>
</tr>
<tr>
<td>MNTE_MODE</td>
<td>The file system type mount method. You can specify a numeric value (see Appendix A) or one of the following predefined variables used to derive the appropriate numeric value:</td>
</tr>
<tr>
<td></td>
<td>MNT_MODE_RDWR                                                                                                                                         File system mounted read-write.</td>
</tr>
<tr>
<td></td>
<td>MNT_MODE_RDONLY                                                                                                                                         File system mounted read-only.</td>
</tr>
<tr>
<td></td>
<td>MNT_MODE_AUNMOUNT                                                                                                                                       The file system can be unmounted if the system’s owner crashes.</td>
</tr>
<tr>
<td></td>
<td>MNT_MODE_NOAUTOMOVE                                                                                                                                     Automove is not allowed.</td>
</tr>
<tr>
<td>MNTE_RFLAGS</td>
<td>Request flags (set to 3 to change the automove attribute, 1 to change mount owner).</td>
</tr>
<tr>
<td>MNTE_SYSNAME</td>
<td>The name of the system to be mounted on.</td>
</tr>
<tr>
<td>MNTE_SYSLIST</td>
<td>A list of system names. For more information about MNTE_SYSLIST, see <a href="https://www.ibm.com/support/knowledgecenter/SS14WS_1.12.0/sboarding/section.html">BPXYMNT — Map response and element structure of w_getmntent in z/OS UNIX System Services Programming: Assembler Callable Services Reference</a>.</td>
</tr>
</tbody>
</table>

**Usage**

1. In order to mount a file system, the caller must be a superuser.
2. The mount service effectively creates a virtual file system. After a file system is mounted, references to the pathname that is mounted refer to the root directory on the mounted file system.
3. A file system can be mounted at only one point.
4. Parameter specifics for the HFS physical file system:
   - The name value must be uppercase and must be the name of the data set.
   - The parm parameter is not used.
5. You can use a wildcard (*) as the last item (or only item) of a system list (MNTE_SYSLIST). A wildcard is allowed only with an INCLUDE list, not with an EXCLUDE list.

**Examples**

1. To mount the HFS data set HFS.BIN.V1R1M0 on the mountpoint /v1r1m0 as a read-only file system:
   
   ```bash
   "mount /v1r1m0 HFS.BIN.V1R1M0 HFS' mtm_rdonly
   ```

2. To mount an HFS read/write with this system as the owner:
   
   ```bash
   m.="  /* set all of the stem variables to the null string */
   m.mnte_mode=mnt_mode_rdwr
   m.mnte_fsname='OMVS.HFS.U.WJS'
   m.mnte_fstype='HFS'
   m.mnte_path='/u/wjs'
   address syscall 'mount m.'
   ```

3. To mount an HFS file system read/write with system SYS2 as the owner and with NOAUTOMOVE:
   
   ```bash
   m.="
   m.mnte_mode=mnt_mode_rdwr+mnt_mode_noautomove
   m.mnte_fsname='OMVS.HFS.U.WJS'
   m.mnte_fstype='HFS'
   m.mnte_path='/u/wjs'
   m.mnte_sysname='SYS2'
   address syscall 'mount m.'
   ```
mnt='''
mnt.mnte_fstype = ''FS1.HFS''
mnt.mnte_path = ''/fs1''
mnt.mnte_mode = mnt_mode_rdwr
mnt.mnte_fstype = ''HFS''

/** build an INCLUDE system list */
syslist = 'SY1 SY2 SY3'
num_systems = d2c(words(syslist),2)
type_syslist = ''0000''x /* or ''0001''x for exclude list */
mnt.mnte_syslist = num_systems || type_syslist

do i = 1 to words(syslist)
mnt.mnte_syslist = mnt.mnte_syslist||left(word(syslist,i),8)
end
address syscall 'mount mnt.'
say 'mount RRR ='' retval errno errnoj'
Function

open invokes the open callable service to access a file and create a file descriptor for it. The file descriptor is returned in RETVAL.

Parameters

pathname
A pathname for the file.

o_flags
One or more numeric values that describe how the file is to be opened. You can specify a numeric value (see Appendix A) or any of the predefined variables that begin with O_ used to derive the appropriate numeric value. For example, the numeric values 128+3 or 131 or the predefined variables o_creat+o_rdwr could be used to specify how the file is to be opened:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>O_APPEND</td>
<td>Set the offset to EOF before each write.</td>
</tr>
<tr>
<td>O_CREAT</td>
<td>Create the file if it does not exist.</td>
</tr>
<tr>
<td>O_EXCL</td>
<td>Fail if the file does exist and O_CREAT is set.</td>
</tr>
<tr>
<td>O_NOCTTY</td>
<td>Do not make this file a controlling terminal for the calling process.</td>
</tr>
<tr>
<td>O_NONBLOCK</td>
<td>Do not block an open, a read, or a write on the file (do not wait for terminal input).</td>
</tr>
<tr>
<td>O_RDONLY</td>
<td>Open for read-only.</td>
</tr>
<tr>
<td>O_RDWR</td>
<td>Open for read and write.</td>
</tr>
<tr>
<td>O_SYNC</td>
<td>Force synchronous updates.</td>
</tr>
<tr>
<td>O_TRUNC</td>
<td>Write, starting at the beginning of the file.</td>
</tr>
<tr>
<td>O_WRONLY</td>
<td>Open for write-only.</td>
</tr>
</tbody>
</table>

mode
A three- or four-digit number, corresponding to the access permission bits. If this optional parameter is not supplied, the mode is defaulted to 000, which is useful for opening an existing file. For an explanation of how mode is handled if you are creating a file, see the usage notes below.

Each digit must be in the range 0–7, and at least three digits must be specified. For more information on permissions, see Appendix B.

Usage notes

When a file is created with the O_CREAT or O_EXCL options, the file permission bits as specified in the mode parameter are modified by the process’s file creation mask (see [umask] on page 188), and then used to set the file permission bits of the file being created.

O_EXCL option: If the O_EXCL bit is set and the create bit is not set, the O_EXCL bit is ignored.
**O_TRUNC option:** Turning on the O_TRUNC bit opens the file as though it had been created earlier but never written into. The mode and owner of the file do not change (although the change time and modification time do); but the file’s contents are discarded. The file offset, which indicates where the next write is to occur, points to the first byte of the file.

**O_NONBLOCK option:** A FIFO special file is a shared file from which the first data written is the first data read. The O_NONBLOCK option is a way of coordinating write and read requests between processes sharing a FIFO special file. It works this way, provided that no other conditions interfere with opening the file successfully:

- If a file is opened read-only and O_NONBLOCK is specified, the open request succeeds. Control returns to the caller immediately.
- If a file is opened write-only and O_NONBLOCK is specified, the open request completes successfully, provided that another process has the file open for reading. If another process does not have the file open for reading, the request ends with RETVAL set to −1.
- If a file is opened read-only and O_NONBLOCK is omitted, the request is blocked (control is not returned to the caller) until another process opens the file for writing.
- If a file is opened write-only and O_NONBLOCK is omitted, the request is blocked (control is not returned to the caller) until another process opens the file for reading.
- If the O_SYNC update option is used, the program is assured that all data updates have been written to permanent storage.

**Example**

To open or create the file `/u/linda/my.exec` with read-write-execute permission for the owner, and to read and write starting at the beginning of the file:

```
"open /u/linda/my.exec" o_rdwr+o_trunc+o_creat 700
```
opendir

Function

opendir invokes the opendir callable service to open a directory stream so that it can be read by rddir. The file descriptor is returned in RETVAL.

Parameters

pathname
A pathname for the directory.

Usage notes

1. You can use opendir and closedir together with the rddir syscall command, but not with the readdir command. The rddir command reads a directory in the readdir callable service format. Alternatively, you can simply use the readdir syscall command to read an entire directory and format it in a stem.

2. The opendir service opens a directory so that the first rddir service (see “rddir” on page 140) starts reading at the first entry in the directory.

3. RETVAL is a file descriptor for a directory only. It can be used only as input to services that expect a directory file descriptor. These services are closedir, rewinddir, and rddir.

Example

To open the directory /u/edman:

"opendir /u/edman"
Function
pathconf invokes the pathconf callable service to determine the current values of a configurable limit or option (variable) that is associated with a file or directory. The limit or option is returned in RETVAL.

Parameters
pathname
A pathname for a file or directory.

name
A numeric value that indicates which limit is returned. You can specify a numeric value (see Appendix A) or the predefined variable beginning with PC_ used to derive the appropriate numeric value.

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>PC_ACL</td>
<td>Test if access control lists (ACLs) are supported.</td>
</tr>
<tr>
<td>PC_ACL_MAX</td>
<td>Maximum number of entries allowed in an ACL.</td>
</tr>
<tr>
<td>PC_LINK_MAX</td>
<td>Maximum value of a file's link count.</td>
</tr>
<tr>
<td>PC_MAX_CANON</td>
<td>Maximum number of bytes in a terminal canonical input line.</td>
</tr>
<tr>
<td>PC_MAX_INPUT</td>
<td>Minimum number of bytes for which space will be available in a terminal input queue; therefore, the maximum number of bytes a portable application may require to be typed as input before reading them.</td>
</tr>
<tr>
<td>PC_NAME_MAX</td>
<td>Maximum number of bytes in a filename (not a string length; count excludes a terminating null).</td>
</tr>
<tr>
<td>PC_PATH_MAX</td>
<td>Maximum number of bytes in a pathname (not a string length; count excludes a terminating null).</td>
</tr>
<tr>
<td>PC_PIPE_BUF</td>
<td>Maximum number of bytes that can be written atomically when writing to a pipe.</td>
</tr>
<tr>
<td>PC_POSIX_CHOWN_RESTRICTED</td>
<td>Change ownership “chown” on page 45 function is restricted to a process with appropriate privileges, and to changing the group ID (GID) of a file only to the effective group ID of the process or to one of its supplementary group IDs.</td>
</tr>
<tr>
<td>PC_POSIX_NO_TRUNC</td>
<td>Pathname components longer than 255 bytes generate an error.</td>
</tr>
<tr>
<td>PC_POSIX_VDISABLE</td>
<td>Terminal attributes maintained by the system can be disabled using this character value.</td>
</tr>
</tbody>
</table>
**Usage notes**

1. If `name` refers to MAX_CANON, MAX_INPUT, or _POSIX_VDISABLE, the following applies:
   - If `pathname` does not refer to a terminal file, the service returns −1 in RETVAL and sets ERRNO to EINVAL.

2. If `name` refers to NAME_MAX, PATH_MAX, or _POSIX_NO_TRUNC, the following applies:
   - If `pathname` does not refer to a directory, the service still returns the requested information using the parent directory of the specified file.

3. If `name` refers to PC_PIPE_BUF, the following applies:
   - If `pathname` refers to a pipe or a FIFO, the value returned applies to the referred-to object itself. If `pathname` refers to a directory, the value returned applies to any FIFOs that exist or can be created within the directory. If `pathname` refers to any other type of file, the pathconf service returns −1 in RETVAL and sets the ERRNO to EINVAL.

4. If `name` refers to PC_LINK_MAX, the following applies:
   - If `pathname` refers to a directory, the value returned applies to the directory.

**Example**

To determine the maximum number of bytes allowed in a pathname in the root directory:

```
"pathconf /" pc_name_max
```
**Function**

*pause* invokes the pause callable service to suspend execution of the calling thread until delivery of a signal that either executes a signal-catching function or ends the thread. See “Using the REXX signal services” on page 9 for more information.

**Usage notes**

1. A thread that calls *pause* does not resume processing until a signal is delivered with an action to either process a signal-handling function or end the thread. Some signals can be blocked by the thread’s signal mask; see “*sigprocmask*” on page 168 for details.
2. If an incoming unblocked signal ends the thread, *pause* never returns to the caller.
3. A return code is set when any failures are encountered that prevent this function from completing successfully.
Function

**pfsctl** invokes the pfsctl service to issue physical file system (PFS) control commands to a PFS. The meaning of the command and argument are specific to and defined by the PFS.

Parameters

**type**
- The type of file system, as defined by the FILESYSTYPE parameter on the BPXPRMxx parmlib member—for example, HFS. Specify this as it is specified on the parmlib member.

**command**
- The name of a PFS control command.

**buffer**
- The name of a buffer containing the argument to be passed to the PFS. The argument is limited to 4095 bytes.

**length**
- An optional numeric value, indicating the length of the argument. If length is not specified, the length of the buffer is used. The maximum length allowed is 4095 bytes.

Usage notes

1. This service is provided for communication between a program running in a user process and a physical file system.
   - It is similar to ioctl, but the command is directed to the physical file system itself rather than to, or for, a particular file or device.

2. As an example of how you could use this function in writing a physical file system, consider the requirement to display status and performance statistics about the physical file system. You can collect this information in the physical file system, but you need a way to display it to the user. For more information about the use of **pfsctl**, see [z/OS DFSMS Using Data Sets](#). With **pfsctl**, your status utility program can easily fetch the information it needs from the physical file system.
Function
pipe invokes the pipe callable service to create a pipe; or an I/O channel that a process can use to communicate with another process, another thread (in this same process or another process), or, in some cases, with itself. Data can be written into one end of the pipe and read from the other end.

Parameters
stem
On return, stem.0 contains the number of variables returned. Two stem variables are returned:

- stem.1 The file descriptor for the end of the pipe that you read from
- stem.2 The file descriptor for the end of the pipe that you write to

Usage notes
When the pipe call creates a pipe, the O_NONBLOCK and FD_CLOEXEC flags are turned off on both ends of the pipe. You can turn these flags on or off by invoking:
- `f_setfl` on page 76 for the flag O_NONBLOCK
- `f_setfd` on page 75 for the flag FD_CLOEXEC

Example
To create a pipe:
"pipe pfd."
Function
`pt3270` invokes the `tcgetattr` and `tcsetattr` callable services to query, set, and reset 3270 passthrough mode.

A REXX program running under a shell started from the OMVS TSO/E command can use this service to send and receive a 3270 data stream or issue TSO/E commands.

Parameters
`fd` The file descriptor for the file.

`option` A number that identifies the service being requested:

<table>
<thead>
<tr>
<th>Number</th>
<th>Service</th>
</tr>
</thead>
<tbody>
<tr>
<td>1</td>
<td>Query 3270 passthrough support for this file. On return, RETVAL will contain a code for the current state of the file descriptor, or -1 if there is an error.</td>
</tr>
<tr>
<td>0 or -1</td>
<td>This file cannot support 3270 passthrough mode.</td>
</tr>
<tr>
<td>1</td>
<td>This file can support 3270 passthrough mode.</td>
</tr>
<tr>
<td>3</td>
<td>This file is currently in 3270 passthrough mode.</td>
</tr>
<tr>
<td>2</td>
<td>Set 3270 passthrough support for this file. If this is attempted on a file that does not support 3270 passthrough mode, on return RETVAL contains -1 and ERRNO contains the value for ENOSYS.</td>
</tr>
<tr>
<td>3</td>
<td>Reset 3270 passthrough support for this file.</td>
</tr>
</tbody>
</table>

Example
The following is an example of a REXX program that can accept a TSO/E command as its argument and issue the command through OMVS using 3270 passthrough mode. This REXXX program would be located in the HFS and run as a command from the shell.

```rexx
/* rexx */
parse arg cmd
if cmd='' then return
address syscall 'pt3270 1 2' /* set passthrough mode on stdout */
if retval=-1 then do
    say 'Cannot set passthrough mode' retval errno errnojr
    return
end
buf='ff51000000010001\x'||, /* OMVS passthrough command */
d2c(length(cmd),4) || cmd
```

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"write 1 buf" */ send command to OMVS */
"read 1 ibuf 1000" /* discard the response */
'pt3270 1 3' /* reset passthrough mode */
return 0
Function

`quiesce` invokes the `quiesce` callable service to quiesce a file system, making the files in it unavailable for use. You must be a superuser to quiesce a file system.

Parameters

`name`

The name of the file system to be quiesced, specified as the name of an HFS data set. You must specify HFS data set names as fully qualified names in uppercase letters. Do not enclose the data set name in single quotes.

Usage notes

1. After a `quiesce` service request, the file system is unavailable for use until a subsequent `unquiesce` service request is received.
2. Users accessing files in a quiesced HFS file system are suspended until an `unquiesce` request for the file system is processed. Other file systems may send an EAGAIN instead of suspending the user.
3. If a file system that is not mounted is quiesced, that file system cannot be mounted until the file system is unquiesced. This ensures that no one can use the file system while it is quiesced.

Example

To quiesce an HFS data set named HFS.USR.SCHOEN:

```
quiesce HFS.USR.SCHOEN
```
Function

**rddir** invokes the readdir callable service to read multiple name entries from a directory and format the information in the readdir callable service format. To format this type of information in a stem, see "readdir" on page 144. The number of entries read is returned in RETVAL.

Parameters

- **fd**
  The file descriptor (a number) for the directory to be read.

- **variable**
  The name of the buffer into which the directory entries are to be read.

- **length**
  The size of the buffer. After the read completes, the length of **variable** is the size of the buffer. The number of entries is returned in RETVAL.

Usage notes

1. You can use this command only with file descriptors opened using the opendir syscall command. The **rddir** syscall command reads a directory in the readdir callable service format. You can use opendir, rewinddir, and closedir together with the **rddir** syscall command, but not with the readdir syscall command. Alternatively, you can use the readdir syscall command to read an entire directory and format it in a stem.

2. The buffer contains a variable number of variable-length directory entries. Only full entries are placed in the buffer, up to the buffer size specified, and the number of entries is returned.

3. Each directory entry returned has the following format:

   - **2-byte Entry_length**
     The total entry length, including itself.

   - **2-byte Name_length**
     Length of the following Member_name subfield.

   - **Member_name**
     A character field of length Name_length. This name is not null-terminated.

   - **File system specific data**
     If name_length + 4 = entry_length, this subfield is not present.

   The entries are packed together, and the length fields are not aligned on any particular boundary.

4. The buffer returned by one call to the readdir service must be used again on the next call to the readdir service to continue reading entries from where you left off. The buffer must not be altered between calls, unless the directory has been rewound.

5. The end of the directory is indicated in either of two ways:
   - A RETVAL of 0 entries is returned.
   - Some physical file systems may return a null name entry as the last entry in the caller's buffer. A null name entry has an Entry_length of 4 and a Name_length of 0.
The caller of the readdir service should check for both conditions.

**Example**
To read the entries from the directory with file descriptor 4 into the buffer named `buf`, which is 300 bytes long:

```
rddir 4 buf 300
```
Function

read invokes the read callable service to read a specified number of bytes from a file into a buffer that you provide. The number of bytes read is returned in RETVAL.

Parameters

fd
The file descriptor (a number) for the file to be read.

variable
The name of the buffer into which the data is to be read.

length
The maximum number of characters to read. After the read completes, the length of variable is the number of bytes read. This value is also returned in RETVAL.

Usage notes

Length:
The value of length is not checked against any system limit.

Access time:
A successful read updates the access time of the file read.

Origin of bytes read:
If the file specified by fd is a regular file, or any other type of file where a seek operation is possible, bytes are read from the file offset associated with the file descriptor. A successful read increments the file offset by the number of bytes read.

For files where no seek operation is possible, there is no file offset associated with the file descriptor. Reading begins at the current position in the file.

Number of bytes read:
When a read request completes, the RETVAL field shows the number of bytes actually read—a number less than or equal to the number specified as length. The following are some reasons why the number of bytes read might be less than the number of bytes requested:

• Fewer than the requested number of bytes remained in the file; the end of file was reached before length bytes were read.

• The service was interrupted by a signal after some but not all of the requested bytes were read. (If no bytes were read, the return value is set to −1 and an error is reported.)

• The file is a pipe, FIFO, or special file and fewer bytes than length specified were available for reading.

There are several reasons why a read request might complete successfully with no bytes read.
(that is, with RETVAL set to 0). For example, zero bytes are read in these cases:

- The call specified a length of zero.
- The starting position for the read was at or beyond the end of the file.
- The file being read is a FIFO file or a pipe, and no process has the pipe open for writing.
- The file being read is a slave pseudoterminal and a zero-length canonical file was written to the master.

Nonblocking: If a process has a pipe open for reading with nonblocking specified, a request to read from the file ends with a return value of −1 a return code of 0, and ERRNO of EAGAIN. But if nonblocking was not specified, the read request is blocked (does not return) until some data is written or the pipe is closed by all other processes that have the pipe open for writing.

Both master and slave pseudoterminals operate this way, too, except that how they act depends on how they were opened. If the master or the slave is opened blocking, the reads are blocked if there is no data. If it is opened nonblocking, EAGAIN is returned if there is no data.

SIGTTOU processing: The read service causes signal SIGTTIN to be sent under the following conditions:

- The process is attempting to read from its controlling terminal, and
- The process is running in a background process group, and
- The SIGTTIN signal is not blocked or ignored, and
- The process group of the process is not orphaned.

If these conditions are met, SIGTTIN is sent. If SIGTTIN has a handler, the handler gets control and the read ends with the return code set to EINTRO. If SIGTTIN is set to default, the process stops in the read and continues when the process is moved to the foreground.

Example
In the following example, assume that fd was assigned a value earlier in the exec. This reads 1000 characters from the file fd into the buffer buf.

"read (fd) buf 1000"
Function

`readdir` invokes the opendir, readdir, and closedir callable services to read multiple name entries from a directory and format the information in a stem.

Parameters

**pathname**
A pathname for a directory.

**stem**
Upon return, `stem.0` contains the number of directory entries returned. `stem.1` through `stem.n` (where `n` is the number of entries returned) each contain a directory entry.

**stem2**
If the optional `stem2` is provided, `stem2.1` through `stem2.n` (where `n` is the number of structures returned) each contain the stat structure for a directory entry.

You can use the predefined variables that begin with ST_ or their equivalent numeric values (see Appendix A) to access the stat values. For example, `stem2.1.st_size` or `stem2.1.8` accesses the file size for the first directory entry. See "`fstat` on page 83" for a list of the ST_ variables.

Usage notes

The `rddir` command reads a directory in the readdir callable service format. You can use `opendir` and `closedir` together with the `rddir` syscall command, but not with the `readdir` syscall command. The `readdir` syscall command reads a directory and formats it in a stem.

Example

To read the root directory and return information about the directory entries and a stat structure for each directory entry:

"readdir / root. rootst."

---

Note: The text provided is a natural representation of the content in the image, converted into plain text format.
Function
readfile invokes the open, read, and close callable services to read from a text file and format it in a stem.

Parameters
pathname A pathname for the file to be read.
stem Upon return, stem.0 contains the number of lines read. stem.1 through stem.n (where n is the number of lines) each contain a line read.

Usage notes
1. The maximum allowable length of a line in the file is 1024 characters. RC=4 is returned when a line is greater than 1024 characters long or is not delimited with a newline character.
2. The newline characters that delimit the lines in a text file are stripped before the lines are saved in the stem.
3. When the RC indicates success, then Retval will be equal to or greater than zero and indicates the corresponding file description number.

Example
In the following example, assume that pathname was assigned a value earlier in the exec:
"readfile {pathname} file."
Function
readlink invokes the readlink callable service to read the contents of a symbolic link. A symbolic link is a file that contains the pathname for another file. The length, in bytes, of the contents of the link is returned in RETVAL. If a variable is specified, the pathname is read into it.

Parameters
pathname A pathname for the symbolic link.
variable The name of the variable to hold the contents of the symbolic link. After the link is read, the length of the variable is the length of the symbolic link.

Example
In the following example, assume that sl and linkbuf were assigned a value earlier in the exec:
"readlink (sl) linkbuf"
Function

**realpath** invokes the realpath callable service to resolve a pathname to a full pathname without any symbolic links.

**Parameters**

- **pathname**: The pathname to resolve.
- **variable**: The name of the variable to contain the resolved pathname that is returned.

**Example**

The following example retrieves the real pathname for the working directory:

```
"realpath . mycwd"
```
Function
rename invokes the rename callable service to change the name of a file or directory.

Parameters
old_pathname
An existing pathname.

new_pathname
A new pathname.

Usage notes
The rename service changes the name of a file or directory from old_pathname to new_pathname. When renaming finishes successfully, the change and modification times for the parent directories of old_pathname and new_pathname are updated.

The calling process needs write permission for the directory containing old_pathname and the directory containing new_pathname. The caller does not need write permission for the files themselves.

Renaming files: If old_pathname and new_pathname are links referring to the same file, rename returns successfully.

If old_pathname is the name of a file, new_pathname must also name a file, not a directory. If new_pathname is an existing file, it is unlinked. Then the file specified as old_pathname is given new_pathname. The pathname new_pathname always stays in existence; at the beginning of the operation, new_pathname refers to its original file, and at the end, it refers to the file that used to be old_pathname.

Renaming directories: If old_pathname is the name of a directory, new_pathname must also name a directory, not a file. If new_pathname is an existing directory, it must be empty, containing no files or subdirectories. If empty, it is removed, as described in rmdir on page 150.

new_pathname cannot be a directory under old_pathname; that is, the old directory cannot be part of the pathname prefix of the new one.

Example
In the following example, assume that old and new were assigned values earlier in the exec:
"rename (old) (new)"
Function
rewinddir invokes the rewinddir callable service to “rewind”, or reset, to the
beginning of an open directory. The next call to rddir reads the first entry in the
directory.

Parameters
fd  The file descriptor (a number) returned from an opendir syscall command. This
    is the file descriptor for the directory to be reset.

Usage notes
You can use this command only with file descriptors opened using the opendir
syscall command. The rddir syscall command reads a directory in the readdir
callable service format. You can use opendir, rewinddir, and closedir together
with the rddir syscall command, but not with the readdir syscall command.
Alternatively, you can use readdir syscall command to read an entire directory and
format it in a stem.

If the contents of the directory you specify have changed since the directory was
opened, a call to the rewinddir service will reset the pointer into the directory to the
beginning so that a subsequent call to the readdir service will read the new
contents.

Example
To rewind the directory associated with file descriptor 4:
"rewinddir 4"
Function

`rmdir` invokes the `rmdir` callable service to remove a directory. The directory must be empty.

Parameters

`pathname` A pathname for the directory.

Usage notes

1. The directory must be empty.
2. If the directory is successfully removed, the change and modification times for the parent directory are updated.
3. If the link count of the directory becomes zero and no process has the directory open, the directory itself is deleted. The space occupied by the directory is freed for new use.
4. If any process has the directory open when the last link is removed, the directory itself is not removed until the last process closes the directory. New files cannot be created under a directory after the last link is removed, even if the directory is still open.

Example

To remove the directory `/u/ehk0`:

```
rmdir /u/ehk0
```
**Function**

`setegid` invokes the `setegid` callable service to set the effective group ID (GID) of the calling process.

**Parameters**

`gid`

The numeric GID that the calling process is to assume.

**Usage notes**

1. If `gid` is equal to the real group ID or the saved set group ID of the process, the effective group ID is set to `gid`.
2. If `gid` is not the same as the real group ID, and the calling process has the appropriate privileges, the effective group ID is set to `gid`.
3. The `setegid` service does not change any supplementary group IDs of the calling process.

**Example**

In the following example, assume that `gid` was assigned a value earlier in the exec:

```
"setegid" gid
```
Function

**seteuid** invokes the seteuid callable service to set the effective user ID (UID) of the calling process.

Parameters

**uid**

The numeric UID that the calling process is to assume.

Usage notes

1. A user can switch to superuser authority (with an effective UID of 0) if the user is permitted to the BPX.SUPERUSER FACILITY class profile within RACF.
2. If **uid** is the same as the process’s real or saved set UID, or the user has the appropriate privilege, the seteuid service sets the effective UID to be the same as **uid**.

Example

In the following example, assume that **uid** was assigned a value earlier in the exec:

```
"seteuid (uid)"
```
Function
setgid invokes the setgid callable service to set the real, effective, and saved set group IDs (GIDs) for the calling process.

Parameters
gid
The numeric GID that the calling process is to assume.

Usage notes
1. If gid is equal to the real group ID or the saved set group ID of the process, the effective group ID is set to gid.
2. If gid is not the same as the real group ID, and the calling process has the appropriate privileges, then the real, saved set, and effective group IDs are set to gid.
3. The setgid service does not change any supplementary group IDs of the calling process.

Example
In the following example, assume that gid was assigned a value earlier in the exec:
*setgid (gid)*
setgrent

setgrent

Function

setgrent invokes the setgrent callable service to rewind, or reset to the beginning, the group database, allowing repeated searches. For more information, see "getgrent" on page 92.
Function

setgroups invokes the setgroups callable service to set the supplemental group list for the process.

Parameters

stem
The name of a stem variable used to set the group list. Upon return, stem.0 contains the number of variables containing a group id. stem.1 to stem.n (where n is the number of variables) each contain one group id number.

Usage notes

A RETVAL of −1 indicates failure.

Example

In the following example, assume that the stem gr. was setup earlier in the exec:
"setgroups gr."
setpgid invokes the setpgid callable service to place a process in a process group. To identify the group, you specify a process group ID. You can assign a process to a different group, or you can start a new group with that process as its leader.

Parameters

- **pid**
  - The numeric process ID (PID) of the process to be placed in a process group. If the ID is specified as 0, the system uses the process ID of the calling process.

- **pgid**
  - The ID of the process group. If the ID is specified as 0, the system uses the process group ID indicated by `pid`.

Usage notes

1. The process group ID to be assigned to the group must be within the calling process's session.
2. The subject process (the process identified by `pid`) must be a child of the process that issues the service, and it must be in the same session; but it cannot be the session leader. It can be the caller.

Example

In the following example, assume that `pid` and `pgid` were assigned values earlier in the exec:

```
"setpgid" pid pgid
```
Function

setpwent invokes the setpwent callable service to effectively rewind the user database to allow repeated searches. For more information, see "getpwent" on page 108.
**Function**

`setregid` invokes the setregid callable service to set the real or effective GIDs for the calling process. If a specified value is set to -1, the corresponding real or effective GID of the calling process is left unchanged.

**Parameters**

- `rgid`  
  The numeric GID value that becomes the real GID for the calling process.

- `egid`  
  The numeric GID value that becomes the effective GID for the calling process.

**Example**

In the following example, assume that `rgid` and `egid` were assigned values earlier in the `exec`:

```
setregid rgid egid
```
Function
setreuid invokes the setreuid callable service to set the real or effective user IDs (UIDs) of the calling process. If a specified value is set to -1, the corresponding real or effective UID of the calling process is left unchanged.

Parameters

ruid
The numeric UID value that becomes the real UID for the calling process.

euid
The numeric UID value that becomes the effective UID for the calling process.

Example
In the following example, assume that ruid and euid were assigned values earlier in the exec:

```
setreuid ruid euid
```
Function

setrlimit invokes the setrlimit callable service to set resource limits for the calling process. A resource limit is a pair of values; one specifies the current limit and the other a maximum limit.

Parameters

resource

The resource whose limit is being set. The maximum resource limit is RLIM_INFINITY.

You can use the predefined variables beginning with RLIMIT_, or their equivalent numeric values, to specify the resource. (See Appendix A for the numeric values.)

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
<th>Allowable Range</th>
</tr>
</thead>
<tbody>
<tr>
<td>RLIMIT_AS</td>
<td>Maximum address space size for a process.</td>
<td>10 485 760—2 147 483 647</td>
</tr>
<tr>
<td>RLIMIT_CORE</td>
<td>Maximum size (in bytes) of a core dump created by a process.</td>
<td>0—2 147 483 647</td>
</tr>
<tr>
<td>RLIMIT_CPU</td>
<td>Maximum amount of CPU time (in seconds) used by a process.</td>
<td>7—2 147 483 647</td>
</tr>
<tr>
<td>RLIMITFSIZE</td>
<td>Maximum files size (in bytes) created by a process.</td>
<td>0—2 147 483 647</td>
</tr>
<tr>
<td>RLIMTNOFILE</td>
<td>Maximum number of open file descriptors for a process.</td>
<td>5—131 072</td>
</tr>
</tbody>
</table>

stem

The name of the stem variable used to set the limit. stem.1 is the first word, which sets the current limit, and stem.2 is the second word, which sets the maximum limit. The values for each word depend on the resource specified. To specify no limit, use RLIM_INFINITY.

Usage notes

1. The current limit may be modified to any value that is less than or equal to the maximum limit. For the RLIMIT_CPU, RLIMIT_Nofile, and RLIMIT_AS resources, if the setrlimit service is called with a current limit that is lower than the current usage, the setrlimit service fails with an EINVAL errno.
2. The maximum limit may be lowered to any value that is greater than or equal to the current limit.
3. The maximum limit can only be raised by a process that has superuser authority.
4. Both the current limit and maximum limit can be changed via a single call to setrlimit.
5. If the setrlimit service is called with a current limit that is greater than the maximum limit, setrlimit returns an EINVAL errno.
6. The resource limit values are propagated across exec and fork. An exception exists for exec. If a daemon process invokes exec and it invoked setuid before invoking exec, the limit values are set based on the limit values specified in parmlib member BPXPRMxx.

7. For a process that is not the only process within an address space, the RLIMIT_CPU and RLIMIT_AS limits are shared with all the processes within the address space. For RLIMIT_CPU, when the current limit is exceeded, action is taken on the first process within the address space. If the action is termination, all the processes within the address space are terminated.

8. In addition to the RLIMIT_CORE limit values, CORE dump defaults are set by SYSMDUMP defaults. See z/OS MVS Initialization and Tuning Guide for information on setting up SYSMDUMP defaults via the IEADMR00 parmlib member.

9. Core dumps are taken in 4160-byte increments. Therefore, RLIMIT_CORE values affect the size of core dumps in 4160-byte increments. For example, if the RLIMIT_CORE current limit value is 4000, core dumps will contain no data. If the RLIMIT_CORE current limit value is 8000, the maximum size of a core dump is 4160 bytes.

10. Limits may have an infinite value of RLIM_INFINITY.

11. If the limit that is specified for RLIMIT_NOFILE is higher than that currently supported by the system, the limit will be reduced to the system maximum when it is used.

12. When setting RLIMIT_NOFILE, the current limit must be set higher than the value of the highest open file descriptor. Attempting to lower the current limit to a value less than or equal to the highest open file descriptor results in an error of EINVAL.

13. When setting RLIMITFSIZE, a limit of 0 prevents the creation of new files and the expansion of existing files.

**Example**

To reduce the maximum number of open files to 100 and the current limit to 50:

```
r.2=100
r.1=50
"setrlimit" rlimit_no file r.
```
Function
setsid invokes the setsid callable service to create a new session with the calling process as its session leader. The caller becomes the group leader of a new process group.

Parameters
sid
The process ID of the calling process, which becomes the session or process group ID of the new process group.

Usage notes
The calling process does not have a controlling terminal.

Example
In the following example, assume that sid was assigned a value earlier in the exec:
"setsid" sid
**Function**

`setuid` invokes the setuid callable service to set the real, effective, and saved set user IDs for the calling process.

**Parameters**

`uid`

The numeric UID the process is to assume.

**Usage notes**

1. A user can switch to superuser authority (with an effective UID of 0) if the user is permitted to the BPX.SUPERUSER FACILITY class profile within RACF.
2. If `uid` is the same as the process’s real UID or the saved set UID, the `setuid` service sets the effective UID to be the same as `uid`.
   
   If `uid` is not the same as the real UID of the process, and the calling process has appropriate privileges, then the real, effective, and saved set UIDs are set to `uid`.

**Example**

In the following example, assume that `uid` was assigned a value earlier in the exec:

```
"setuid" uid
```
Function

`sigaction` invokes the `sigaction` callable service to examine, change, or both examine and change the action associated with a specific signal for all the threads in the process.

**Note:** All threads within a process share the signal handlers (a set of additional signals to be masked) and the flags specified by the `sigaction` callable service.

**Parameters**

**signal**

The signal, as specified by a numeric value or a predefined variable beginning with SIG.

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>SIGABND</td>
<td>Abend</td>
</tr>
<tr>
<td>SIGABRT</td>
<td>Abnormal termination</td>
</tr>
<tr>
<td>SIGALRM</td>
<td>Timeout</td>
</tr>
<tr>
<td>SIGBUS</td>
<td>Bus error</td>
</tr>
<tr>
<td>SIGCHLD</td>
<td>Child process terminated or stopped</td>
</tr>
<tr>
<td>SIGCONT</td>
<td>Continue if stopped</td>
</tr>
<tr>
<td>SIGDCE</td>
<td>Exclusive use by DCE</td>
</tr>
<tr>
<td>SIGFPE</td>
<td>Erroneous arithmetic operation, such as division by zero or an overflow resulting in an overflow</td>
</tr>
<tr>
<td>SIGHUP</td>
<td>Hangup detected on the controlling terminal</td>
</tr>
<tr>
<td>SIGILL</td>
<td>Termination (cannot be caught or ignored)</td>
</tr>
<tr>
<td>SIGINT</td>
<td>Interactive attention</td>
</tr>
<tr>
<td>SIGIO</td>
<td>Completion of input or output</td>
</tr>
<tr>
<td>SIGIOERR</td>
<td>Error on input/output; used by the C runtime library</td>
</tr>
<tr>
<td>SIGKILL</td>
<td>Termination that cannot be caught or ignored</td>
</tr>
<tr>
<td>SIGPIPE</td>
<td>Write on a pipe with no readers</td>
</tr>
<tr>
<td>SIGPOLL</td>
<td>Pollable event</td>
</tr>
<tr>
<td>SIGPROF</td>
<td>Profiling timer expired</td>
</tr>
<tr>
<td>SIGQUIT</td>
<td>Interactive termination</td>
</tr>
<tr>
<td>SIGSEGV</td>
<td>Detection of an incorrect memory reference</td>
</tr>
<tr>
<td>SIGSTOP</td>
<td>Stop that cannot be caught or ignored</td>
</tr>
<tr>
<td>SIGSYS</td>
<td>Bad system call</td>
</tr>
<tr>
<td>SIGTERM</td>
<td>Termination</td>
</tr>
<tr>
<td>SIGTRAP</td>
<td>Trap used by the ptrace callable service</td>
</tr>
</tbody>
</table>
sigaction

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>SIGTSTP</td>
<td>Interactive stop</td>
</tr>
<tr>
<td>SIGTTIN</td>
<td>Read from a controlling terminal attempted by a member of a background process group</td>
</tr>
<tr>
<td>SIGTTOU</td>
<td>Write from a controlling terminal attempted by a member of a background process group</td>
</tr>
<tr>
<td>SIGURG</td>
<td>High bandwidth data is available at a socket</td>
</tr>
<tr>
<td>SIGUSR1</td>
<td>Reserved as application-defined signal 1</td>
</tr>
<tr>
<td>SIGUSR2</td>
<td>Reserved as application-defined signal 2</td>
</tr>
<tr>
<td>SIGVTALRM</td>
<td>Virtual timer expired</td>
</tr>
<tr>
<td>SIGXCPU</td>
<td>CPU time limit exceeded</td>
</tr>
<tr>
<td>SIGXFSZ</td>
<td>File size limit exceeded</td>
</tr>
</tbody>
</table>

**new_handler**

Specifies the new setting for handling the signal. The following predefined variables may be used for `new_handler`:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>SIG_CAT</td>
<td>Set signal handling to catch the signal.</td>
</tr>
<tr>
<td>SIG_DFL</td>
<td>Set signal handling to the default action.</td>
</tr>
<tr>
<td>SIG_IGN</td>
<td>Set signal handling to ignore the signal.</td>
</tr>
<tr>
<td>SIG_QRY</td>
<td>Query the handling for that signal.</td>
</tr>
</tbody>
</table>

**new_flag**

Used to modify the behavior of the specified signal. Specify one of these:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>Use the default behavior.</td>
</tr>
<tr>
<td>SA_NOCLDWAIT</td>
<td>Do not create zombie processes when child processes exit.</td>
</tr>
<tr>
<td>SA_RESETHAND</td>
<td>Reset the signal action to the default, SIG_DFL, when delivered.</td>
</tr>
<tr>
<td>SA_NOCLDSTOP</td>
<td>Do not generate SIGCHLD when the child processes stop.</td>
</tr>
</tbody>
</table>

**old_handler**

A variable name for the buffer where the system returns the old (current) signal handling action.

**old_flag**

The name of the variable that will store the old (current) signal action flags.

**Usage notes**

1. If `new_handler` is set to the action SIG_DFL for a signal that cannot be caught or ignored, the sigaction request is ignored and the return value is set to 0.
2. Setting a signal action to ignore for a signal that is pending causes the pending signal to be discarded.
3. Setting signal action SIG_IGN or catch for signals SIGSTOP or SIGKILL is not allowed.
4. Setting signal action SIG_IGN for SIGCHLD or SIGIO is not allowed.
5. The sigaction caller’s thread must be registered for signals. You can register the thread by calling `syscalls('SIGON')`. If the thread is not registered for signals, the sigaction service fails with an ERRNO of EINVAL and ERRNOJR of JRNotSigSetup.

**Example**
To catch a SIGALRM signal:
```
"sigaction" sigalrm sig_cat 0 "prevhndlr prevflag"
```
sigpending

Function

**sigpending** invokes the sigpending callable service to return the union of the set of signals pending on the thread and the set of signals pending on the process. Pending signals at the process level are moved to the thread that called the sigpending callable service.

**Parameters**

**variable**

The name of the variable that will store a string of 64 characters with values 0 or 1, representing the 64 bits in a signal mask.

**Example**

To invoke sigpending:

```
sigpending sigset
```
sigprocmask

Function

sigprocmask invokes the sigprocmask callable service to examine or change the calling thread’s signal mask.

Parameters

number
To specify the action to be taken on the thread’s signal mask, you can specify a numeric value (see Appendix A) or the predefined variable beginning with SIG_ used to derive the appropriate numeric value. Use one of the following predefined variables:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>SIG_BLOCK</td>
<td>Add the signals in new_mask to those to be blocked for this thread.</td>
</tr>
<tr>
<td>SIG_SETMASK</td>
<td>Replace the thread’s signal mask with new_mask.</td>
</tr>
<tr>
<td>SIG_UNBLOCK</td>
<td>Delete the signals in new_mask from those blocked for this thread.</td>
</tr>
</tbody>
</table>

new_mask
The new signal mask, a string of 64 characters with values 0 or 1. The first character represents signal number 1. A string shorter than 64 characters is padded on the right with zeros. Mask bits set on represent signals that are blocked. For more information on signals, see "Using the REXX signal services" on page 9.

variable
The name of the buffer that will store the old signal mask, a string of 64 characters with values 0 or 1, representing the 64 bits in a signal mask. Mask bits set on represent signals that are blocked. A zero indicates that no signal mask was returned.

Usage notes

1. The sigprocmask service examines, changes, or both examines and changes the signal mask for the calling thread. This mask is called the thread’s signal mask. If there are any pending unblocked signals, either at the process level or at the current thread’s level after changing the signal mask, at least one of the signals is delivered to the thread before the sigprocmask service returns.

2. You cannot block the SIGKILL and the SIGSTOP signals. If you call the sigprocmask service with a request that would block those signals, that part of your request is ignored and no error is indicated.

3. A request to block signals that are not supported is accepted, and a return value of zero is returned.

4. All pending unblocked signals are moved from the process level to the current thread.
Example
In the following example, assume that `newsigset` was assigned a value earlier in the exec:
```
"sigprocmask" sig_setmask newsigset "oldsigset"
```
Function
sigsuspend invokes the sigsuspend callable service to replace a thread’s current signal mask with a new signal mask; then it suspends the caller’s thread until delivery of a signal whose action is either to process a signal-catching service or to end the thread.

Parameters
mask The new signal mask, a string of up to 64 characters with the values 0 or 1. The first character represents signal number 1. A string shorter than 64 characters is padded on the right with zeros. For more information on signals, see “Using the REXX signal services” on page 9.

Usage notes
1. The caller’s thread starts running again when it receives one of the signals not blocked by the mask set by this call, or a system failure occurs that sets the return code to some value other than EINTR.
2. The signal mask represents a set of signals that will be blocked. Blocked signals do not “wake up” the suspended service. The signals SIGSTOP and SIGKILL cannot be blocked or ignored; they are delivered to the program no matter what the signal mask specifies.
3. If the signal action is to end the thread, the sigsuspend service does not return.
4. All pending unblocked signals are moved from the process level to the current thread.

Example
In the following example, assume that sigmask was assigned a value earlier in the exec:
"sigsuspend" sigmask
Function
sleep invokes the sleep callable service to suspend running of the calling thread (process) until either the number of seconds specified by number has elapsed, or a signal is delivered to the calling thread to invoke a signal-catching function or end the thread.

Parameters
number
The number of seconds to suspend the process. For more information on signals, see "Using the REXX signal services" on page 9.

Usage notes
1. The suspension can actually be longer than the requested time, due to the scheduling of other activity by the system.
2. The sleep service suspends the thread running for a specified number of seconds, or until a signal is delivered to the calling thread that invokes a signal-catching function or ends the thread. An unblocked signal received during this time prematurely "wakes up" the thread. The appropriate signal-handling function is invoked to handle the signal. When that signal-handling function returns, the sleep service returns immediately, even if there is "sleep time" remaining.
3. The sleep service returns a zero in RETVAL if it has slept for the number of seconds specified. If the time specified by number has not elapsed when the sleep service is interrupted because of the delivery of a signal, the sleep service returns the unslept amount of time (the requested time minus the time actually slept when the signal was delivered) in seconds. Any time consumed by signal-catching functions is not reflected in the value returned by the sleep service.
4. The following are usage notes for a SIGALRM signal generated by the alarm or kill calls during the execution of the sleep call:
   • If the calling thread has SIGALRM blocked prior to calling the sleep service, the sleep service does not return when SIGALRM is generated, and the SIGALRM signal is left pending when sleep returns.
   • If the calling process has SIGALRM ignored when the SIGALRM signal is generated, then the sleep service does not return and the SIGALRM signal is ignored.
   • If the calling process has SIGALRM set to a signal-catching function, that function interrupts the sleep service and receives control. The sleep service returns any unslept amount of time, as it does for any other type of signal.
5. An EC6 abend is generated when the caller’s PSW key or RB state prevents signals from being delivered.

Example
In the following example, assume that timer was assigned a value earlier in the exec:
"sleep (timer)"
spawn

Function

spawn invokes the spawn callable service to create a new process, called a child process, to run an executable file. It remaps the calling process's file descriptors for the child process.

Parameters

pathname

A pathname for the executable file. Pathnames can begin with or without a slash:
- A pathname that begins with a slash is an absolute pathname, and the search for the file starts at the root directory.
- A pathname that does not begin with a slash is a relative pathname, and the search for the file starts at the working directory.

fd_count

The number of file descriptors that can be inherited by the child process. In the new process, all file descriptors greater than or equal to fd_count are closed.

fd_map

A stem variable. The stem index specifies the child's file descriptor; for example, stem.0 specifies the child's file descriptor 0. This array selects the file descriptors to be inherited. The value assigned to the variable indicates the parent's file descriptor that will be mapped to the child's file descriptor. For example, if stem.0 is 4, the child process inherits the parent's file descriptor 4 as its descriptor 0. Any of the stem variables that contains a negative number or a nonnumeric value is closed in the child.

arg_stem

A stem variable. stem.0 contains the number of arguments you are passing to the program. The first argument should always specify the absolute or relative pathname of the program to be executed. If a relative pathname is used and PATH is specified, PATH is used to resolve the name; otherwise, the name is processed as relative to the current directory. If a PATH environment variable is not passed, the first argument should specify the absolute pathname or a relative pathname for the program.

denv_stem

A stem variable. stem.0 contains the number of environment variables that you want the program to be run with. To propagate the current environment, pass _ _environment. Specify each environment variable as VNAME=value.

Usage notes

1. The new process (called the child process) inherits the following attributes from the process that calls spawn (called the parent process):
   - Session membership.
   - Real user ID.
   - Real group ID.
   - Supplementary group IDs.
2. The new child process has the following differences from the parent process:
   - The child process has a unique process ID (PID) that does not match any active process group ID.
   - The child has a different parent PID (namely, the PID of the process that called `spawn`).
   - If the `fd_count` parameter specified a 0 value, the child has its own copy of the parent’s file descriptors, except for those files that are marked FCTLCLLOEXEC or FCTLCLLOFORK. The files marked FCTLCLLOEXEC or FCTLCLLOFORK are not inherited by the child. If the `filedesc_count` parameter specifies a value greater than 0, the parent’s file descriptors are remapped for the child as specified in the `fd_map` stem with a negative number or non-numeric value.
   - The FCTLCLLOEXEC and FCTLCLLOFORK flags are not inherited from the parent file descriptors to the child’s.
   - The foreground process group of the session remains unchanged.
   - The process and system utilization times for the child are set to zero.
   - Any file locks previously set by the parent are not inherited by the child.
   - The child process has no alarms set (similar to the results of a call to the alarm service with `Wait_time` specified as zero) and has no interval timers set.
   - The child has no pending signals.
   - The child gets a new process image, which is not a copy of the parent process, to run the executable file.
   - Signals set to be caught are reset to their default action.
   - Memory mappings established by the parent via the `shmem` or `mmap` services are not inherited by the child.
   - If the `setuid` bit of the new executable file is set, the effective user ID and saved set-user-ID mode of the process are set to the owner user ID of the new executable file.
   - If the `setgid` bit of the new executable file is set, the effective group ID and saved set-group-ID bit of the process are set to the owner user ID of the new executable file.

   The last parameter that `spawn` passed to the executable file identifies the caller of the file as the exec or spawn service.

3. To control whether the spawned child process runs in a separate address space from the parent address space or in the same address space, you can specify the `_BPX_SHAREAS` environment variable. If `_BPX_SHAREAS` is not specified, is set to NO, or contains an unsupported value, the child process to be created will run in a separate address space from the parent process.

   `_BPX_SHAREAS=YES` indicates that the child process to be created is to run in the same address space as the parent. If the program to be run is a set-user-ID or set-group-ID program that will cause the effective user-ID or group-ID of the
spawn

child process to be different from that of the parent process, the
_BPX_SHAREAS=YES value is ignored and the child process runs in its own
address space.

4. In addition to recognizing the _BPX_SHAREAS environment variable, spawn
recognizes all of the environment variables that are recognized by the fork and
exec callable services.

5. The fd_count parameter can be 0, which means that all file descriptors are
inherited by the child.

6. The fd_count parameter is limited to a maximum value of 1000.

7. When the executable file to be run is a REXX exec, the first argument should
be the pathname of the REXX exec. Subsequent arguments for the exec can
follow this.

Example
In the following example, _bin/l_ is run mapping its STDOUT and STDERR to file
descriptors 4 and 5, which were previously opened in the exec, and STDIN is
closed:

map.0=-1
map.1=4
map.2=5
parm.0=2
parm.1='//bin/l_'
parm.2=''
'spawn /bin/l_ 3 map. parm. _ _environment.'
Function

spawnp invokes the spawn callable service and creates a new process, called a child process, to run a hierarchical file system (HFS) executable file. spawnp functions identically to the spawn function, except that it uses the PATH environment variable to resolve relative filenames.

See "spawn" on page 172 for more information.
**Function**

*stat* invokes the *stat* callable service to obtain status about a specified file. You specify the file by its name. If the pathname specified refers to a symbolic link, the symbolic link name is resolved to a file and the status information for that file is returned. To obtain status information about a symbolic link, rather than the file it refers to, see **"lstat" on page 121.**

To use a file descriptor to obtain this information, see **"fstat" on page 83.**

**Parameters**

*pathname*

A pathname for the file.

*stem*

The name of a stem variable used to return the information. Upon return, *stem.0* contains the number of variables returned. You can use the predefined variables beginning with ST_ or their equivalent numeric values to access the values they represent. (See **Appendix A** for the numeric values.)

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>ST_AUDIT</td>
<td>Auditor audit information</td>
</tr>
<tr>
<td>ST_ACCESSACL</td>
<td>1 if there is an access ACL (access control list)</td>
</tr>
<tr>
<td>ST_ATIME</td>
<td>Time of last access</td>
</tr>
<tr>
<td>ST_AUDITID</td>
<td>RACF File ID for auditing</td>
</tr>
<tr>
<td>ST_BLKSIZE</td>
<td>File block size</td>
</tr>
<tr>
<td>ST_BLOCKS</td>
<td>Blocks allocated</td>
</tr>
<tr>
<td>ST_CCSID</td>
<td>Coded character set ID; first 4 characters are the file tag</td>
</tr>
<tr>
<td>ST_CRTIME</td>
<td>File creation time</td>
</tr>
<tr>
<td>ST_CTIME</td>
<td>Time of last file status change</td>
</tr>
<tr>
<td>ST_DEV</td>
<td>Device ID of the file</td>
</tr>
<tr>
<td>ST_DMODELACL</td>
<td>1 if there is a directory model access control list (ACL)</td>
</tr>
<tr>
<td>ST_EXTLINK</td>
<td>External symbolic link flag, set to 0 or 1</td>
</tr>
<tr>
<td>ST_FID</td>
<td>File identifier</td>
</tr>
<tr>
<td>Variable</td>
<td>Description</td>
</tr>
<tr>
<td>--------------</td>
<td>-------------</td>
</tr>
<tr>
<td>ST_FILEFMT</td>
<td>Format of the file. To specify the format, you can specify a numeric value (see Appendix A) or one of the following predefined variables used to derive the appropriate numeric value:</td>
</tr>
<tr>
<td>S_FFBINARY</td>
<td>Binary data</td>
</tr>
<tr>
<td>S_FFCR</td>
<td>Text data delimited by a carriage return character</td>
</tr>
<tr>
<td>S_FFCRLF</td>
<td>Text data delimited by carriage return and line feed characters</td>
</tr>
<tr>
<td>S_FFLF</td>
<td>Text data delimited by a line feed character</td>
</tr>
<tr>
<td>S_FFLFCR</td>
<td>Text data delimited by a line feed and carriage return characters</td>
</tr>
<tr>
<td>S_FFNA</td>
<td>Text data with the file format not specified</td>
</tr>
<tr>
<td>S_FFNL</td>
<td>Text data delimited by a newline character</td>
</tr>
<tr>
<td>ST_FMODEACL</td>
<td>1 if there is a file model ACL</td>
</tr>
<tr>
<td>ST_GENVALUE</td>
<td>General attribute values</td>
</tr>
<tr>
<td>ST_GID</td>
<td>Group ID of the group of the file</td>
</tr>
<tr>
<td>ST_INO</td>
<td>File serial number</td>
</tr>
<tr>
<td>ST_MAJOR</td>
<td>Major number for a character special file</td>
</tr>
<tr>
<td>ST_MINOR</td>
<td>Minor number for a character special file</td>
</tr>
<tr>
<td>ST_MODE</td>
<td>File mode, permission bits only</td>
</tr>
<tr>
<td>ST_MTIME</td>
<td>Time of last data modification</td>
</tr>
<tr>
<td>ST_NLINK</td>
<td>Number of links</td>
</tr>
<tr>
<td>ST_RTIME</td>
<td>File backup time stamp (reference time)</td>
</tr>
<tr>
<td>ST_SETGID</td>
<td>Set Group ID on execution flag, set to 0 or 1</td>
</tr>
<tr>
<td>ST_SETUID</td>
<td>Set User ID on execution flag, set to 0 or 1</td>
</tr>
<tr>
<td>ST_SIZE</td>
<td>File size for a regular file, in bytes. If file size exceeds $2^{31} - 1$ bytes, size is expressed in megabytes, using an M (for example, 3123M).</td>
</tr>
<tr>
<td>ST_SECLABEL</td>
<td>Security Label</td>
</tr>
<tr>
<td>ST_STICKY</td>
<td>Sticky bit flag (keep loaded executable in storage), set to 0 or 1</td>
</tr>
<tr>
<td>ST_TYPE</td>
<td>Numeric value that represents the file type for this file. You can use a numeric value (see Appendix A) or any of the predefined variables that begin with S_ to determine the file type:</td>
</tr>
<tr>
<td>S_ISCHR</td>
<td>Character special file</td>
</tr>
<tr>
<td>S_ISDIR</td>
<td>Directory</td>
</tr>
<tr>
<td>S_ISFIFO</td>
<td>FIFO special file</td>
</tr>
<tr>
<td>S_ISREG</td>
<td>Regular file</td>
</tr>
<tr>
<td>S_ISSYM</td>
<td>Symbolic link</td>
</tr>
<tr>
<td>ST_UAUDIT</td>
<td>Area for user audit information</td>
</tr>
<tr>
<td>ST_UID</td>
<td>User ID of the owner of the file</td>
</tr>
</tbody>
</table>

The stem variable `stem.st_type` is a number that represents the file type for this file. You can use the predefined variables beginning with `S_` or their equivalent numeric values to determine the file type. For example, if `stem.st_type` is `S_ISDIR`, the file is a directory.
Usage notes
All time fields in `stem` are in POSIX format. You can use `gmtime` to convert it to other forms.

Example
In the following example, assume that `path` was assigned a value earlier in the exec:
"`stat (path) st.`"
Function

**statfs** invokes the statfs callable service to obtain status information about a specified file system.

Parameters

**name**

The name of the file system to be mounted, specified as the name of an HFS data set. You must specify the HFS data set name as a fully qualified name in uppercase letters. Do not enclose the data set name in single quotes.

**stem**

The name of a stem variable used to return the information. On return, **stem.0** contains the number of variables returned. You can use the predefined variables beginning with STFS_ or their equivalent numeric values to access the status values they represent. (See Appendix A for the numeric values.) For example, **stem.stfs_avail** accesses the number of blocks available in the file system.

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>STFS_AVAIL</td>
<td>Space available to unprivileged users in block-size units.</td>
</tr>
<tr>
<td>STFS_BFREE</td>
<td>Total number of free blocks.</td>
</tr>
<tr>
<td>STFS_BLOCKSIZE</td>
<td>Block size.</td>
</tr>
<tr>
<td>STFS_FAVAIL</td>
<td>Number of free file nodes available to unprivileged users.</td>
</tr>
<tr>
<td>STFS_FFREE</td>
<td>Total number of free file nodes.</td>
</tr>
<tr>
<td>STFS_FILES</td>
<td>Total number of file nodes in the file system.</td>
</tr>
<tr>
<td>STFS_FRSIZE</td>
<td>Fundamental file system block size.</td>
</tr>
<tr>
<td>STFS_FSID</td>
<td>File system ID set by the logical file system.</td>
</tr>
<tr>
<td>STFS_INUSE</td>
<td>Allocated space in block-size units.</td>
</tr>
<tr>
<td>STFS_INVARSEC</td>
<td>Number of seconds the file system will remain unchanged.</td>
</tr>
<tr>
<td>STFS_NAMEMAX</td>
<td>Maximum length of file name.</td>
</tr>
<tr>
<td>STFS_NOSEC</td>
<td>Mount data set with no security bit.</td>
</tr>
<tr>
<td>STFS_NOSUID</td>
<td>SETUID and SETGID are not supported.</td>
</tr>
<tr>
<td>STFS_RDONLY</td>
<td>File system is read-only.</td>
</tr>
<tr>
<td>STFS_TOTAL</td>
<td>Total space in block-size units.</td>
</tr>
</tbody>
</table>

Example

In the following example, assume that **fsname** was assigned a value earlier in the exec:

```
"statfs (fsname) st."
```
Function

`statvfs` invokes the `statvfs` callable service to obtain status information about a file system, given the name of a file in the file system.

Parameters

**pathname**

The name of a file in a file system for which status information is to be obtained.

**stem**

The name of a stem variable used to return the information. On return, `stem.0` contains the number of variables returned. You can use the predefined variables beginning with `STFS_` or their equivalent numeric values (see Appendix A) to access the status values they represent. For example, `stem.stfs_avail` accesses the number of blocks available in the file system.

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>STFS_AVAIL</td>
<td>Space available to unprivileged users in block-size units.</td>
</tr>
<tr>
<td>STFS_BFREE</td>
<td>Total number of free blocks.</td>
</tr>
<tr>
<td>STFS_BLOCKSIZE</td>
<td>Block size.</td>
</tr>
<tr>
<td>STFS_FAVAIL</td>
<td>Number of free file nodes available to unprivileged users.</td>
</tr>
<tr>
<td>STFS_FFREE</td>
<td>Total number of free file nodes.</td>
</tr>
<tr>
<td>STFS_FILES</td>
<td>Total number of file nodes in the file system.</td>
</tr>
<tr>
<td>STFS_FRSIZE</td>
<td>Fundamental file system block size.</td>
</tr>
<tr>
<td>STFS_FSID</td>
<td>File system ID set by the logical file system.</td>
</tr>
<tr>
<td>STFS_INUSE</td>
<td>Allocated space in block-size units.</td>
</tr>
<tr>
<td>STFS_INWARSEC</td>
<td>Number of seconds the file system will remain unchanged.</td>
</tr>
<tr>
<td>STFS_NAMEMAX</td>
<td>Maximum length of file name.</td>
</tr>
<tr>
<td>STFS_NOSEC</td>
<td>Mount data set with no security bit.</td>
</tr>
<tr>
<td>STFS_NOSUID</td>
<td>SETUID and SETGID are not supported.</td>
</tr>
<tr>
<td>STFS_RDONLY</td>
<td>File system is read-only.</td>
</tr>
<tr>
<td>STFS_TOTAL</td>
<td>Total space in block-size units.</td>
</tr>
</tbody>
</table>

Example

In the following example, assume that `fsname` was assigned a value earlier in the exec:

```
"statvfs (fsname) st."
```
Function
strerror retrieves diagnostic text for error codes and reason codes.

Parameters
error_code
   Hex value for an error code as returned in ERRNO for other SYSCALL host commands. Specify 0 if text for this code is not being requested.

reason_code
   Hex value for the reason code as returned in ERRNOJR for other SYSCALL host commands. Specify 0 if this code is not being requested.

stem
   The name of a stem variable used to return the information. On return, stem.0 contains the number of variables returned. You can use the predefined variables beginning with SE_ or their equivalent numeric values to access the values that they represent. See Appendix A for the numeric values. For example, stem.se_reason accesses the reason code text. If text is unavailable, a null string is returned.

Variable   Description
SE_ERRNO   Text for the error number.
SE_REASON   Text for the reason code.
SE_ACTION   Text for any action to be taken to correct this error. This variable will be available only when a reason code is requested.
SE_MODID    Name of the module that detected the error. This variable will be available only when a reason code is requested.

Example
To get error text for the last syscall error:
"strerror" errno errnojr "err."
Function

`symlink` invokes the symlink callable service to create a symbolic link to a pathname. This creates a symbolic link file.

Parameters

pathname
A pathname for the file for which you are creating a symbolic link.

linkname
The pathname for the symbolic link.

Usage notes

Like a hard link (described in "link" on page 119), a symbolic link allows a file to have more than one name. The presence of a hard link guarantees the existence of a file, even after the original name has been removed. A symbolic link, however, provides no such assurance; in fact, the file identified by `pathname` need not exist when the symbolic link is created. In addition, a symbolic link can cross file system boundaries.

When a component of a pathname refers to a symbolic link rather than to a directory, the pathname contained in the symbolic link is resolved. If the pathname in the symbolic link begins with `/` (slash), the symbolic link pathname is resolved relative to the process root directory. If the pathname in the symbolic link does not begin with `/`, the symbolic link pathname is resolved relative to the directory that contains the symbolic link.

If the symbolic link is not the last component of the original pathname, remaining components of the original pathname are resolved from there. When a symbolic link is the last component of a pathname, it may or may not be resolved. Resolution depends on the function using the pathname. For example, a rename request does not have a symbolic link resolved when it appears as the final component of either the new or old pathname. However, an open request does have a symbolic link resolved when it appears as the last component. When a slash is the last component of a pathname, and it is preceded by a symbolic link, the symbolic link is always resolved.

Because the mode of a symbolic link cannot be changed, its mode is ignored during the lookup process. Any files and directories to which a symbolic link refers are checked for access permission.

Example

To create a symbolic link named `/bin` for the file `/v.1.1.0/bin`:

```
"symlink /v.1.1.0/bin /bin"
```
Function
sysconf invokes the sysconf callable service to get the value of a configurable system variable.

Parameters
name
A numeric value that specifies the configurable variable to be returned. You can specify a numeric value (see Appendix A) or the predefined variable beginning with SC_ used to derive the appropriate numeric value:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>SC_ARG_MAX</td>
<td>The maximum length of all arguments and environment strings to exec()</td>
</tr>
<tr>
<td>SC_CHILD_MAX</td>
<td>The maximum number of simultaneous processes per real user ID</td>
</tr>
<tr>
<td>SC_CLK_TCK</td>
<td>The number of intervals per second used in defining the type clock_t, which is used to measure process execution times</td>
</tr>
<tr>
<td>SC_JOB_CONTROL</td>
<td>Support for job control</td>
</tr>
<tr>
<td>SC_NGROUPS_MAX</td>
<td>Maximum number of simultaneous supplementary group IDs per process</td>
</tr>
<tr>
<td>SC_OPEN_MAX</td>
<td>Maximum number of simultaneous open files per process</td>
</tr>
<tr>
<td>SC_SAVED_IDS</td>
<td>Support for saved set-user-IDs and set-group-IDs</td>
</tr>
<tr>
<td>SC_THREAD_TASKS_MAX_NP</td>
<td>Constant for querying the maximum number of threaded tasks per calling process</td>
</tr>
<tr>
<td>SC_THREADS_MAX_NP</td>
<td>Constant for querying the maximum number of threads per calling process</td>
</tr>
<tr>
<td>SC_TZNAME_MAX</td>
<td>The number of bytes supported for the name of a time zone</td>
</tr>
<tr>
<td>SC_VERSION</td>
<td>The integer value 199009L</td>
</tr>
<tr>
<td>SC_2_CHAR_TERM</td>
<td>Constant for querying whether the system supports at least one raw mode terminal</td>
</tr>
</tbody>
</table>

Usage notes
SC_MAX_THREADS_NP and SC_MAX_THREAD_TASKS_NP return the limits defined for the caller's process, not the systemwide limits.

Example
To determine the maximum number of files that a single process can have open at one time:

```
"sysconf" sc_open_max
maxopenfiles = retval  /*value of desired configurable system variable is returned in retval*/
```
Function

time returns in RETVAL the time in POSIX format (seconds since the Epoch, 00:00:00 on 1 January 1970). You can use gmtime to convert it to other forms.
Function
times invokes the times callable service to collect information about processor time used by the current process or related processes. The elapsed time since the process was dubbed is returned in RETVAL. This value is of the type clock_t, which needs to be divided by sysconf(_SC_CLK_TK) to convert it to seconds. For z/OS UNIX, this value is expressed in hundredths of a second.

Parameters
stem
The name of a stem variable used to return the information. Upon return, stem.0 contains the number of variables returned. Four variables are returned in the stem. To access the stem variables, use a numeric value or the predefined variables beginning with TMS_ used to derive that numeric value. (For the numeric values, see Appendix A, "REXX predefined variables," on page 289.) For example, you could specify stem.4 or stem.tms_cstime to obtain system CPU values:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>TMS_CSTIME</td>
<td>The sum of system CPU time values and child system CPU time values for all waited-for child processes that have terminated. Zero if the current process has no waited-for children.</td>
</tr>
<tr>
<td>TMS_CUTIME</td>
<td>The sum of user CPU time values and child user CPU time values for all waited-for child processes that have terminated. Zero if the current process has no waited-for children.</td>
</tr>
<tr>
<td>TMS_STIME</td>
<td>The system CPU time of current process in hundredths of a second. This is the task control block (TCB) time accumulated while running in the kernel address space.</td>
</tr>
<tr>
<td>TMS_UTIME</td>
<td>The user CPU time of current process in hundredths of a second. This includes the TCB and service request block (SRB) time of the calling process minus the TCB time accumulated while running in the kernel address space.</td>
</tr>
</tbody>
</table>

Usage notes
Processor times for a child process that has ended are not added to the TMS_CUTIME and TMS_CSTIME of the parent process until the parent issues a wait or waitpid for that child process.

Example
"times tm."
trunc

trunc

Function
trunc invokes the trunc callable service to change the size of the file identified by pathname.

Parameters
pathname
   The pathname of the file.

file_size
   The new size of the file, in bytes.

Usage notes
1.   The file specified must be a regular file to which the calling process has write access.
2.   The file size changes beginning from the first byte of the file. If the file was previously larger than the new size, the data from file_size to the original end of the file is removed. If the file was previously shorter than file_size, bytes between the old and new lengths are read as zeros.
3.   If file_size is greater than the current file size limit for the process, the request fails with EFBIG, and the SIGXFSZ signal is generated for the process.

Example
To set the file size of /tmp/xxx to 1000 bytes:
"trunc /tmp/xxx 1000"
Function

`ttyname` invokes the `ttyname` callable service to obtain the pathname of the terminal associated with the file descriptor.

Parameters

- **fd**  The file descriptor (a number) for the character special file for the terminal.
- **variable**  The name of the variable that stores the pathname for the character special file for the terminal.

Usage notes

This service does not return −1 to indicate a failure. If the file descriptor is incorrect, a null string is returned.

Example

To obtain the pathname for file descriptor 0:

```
"ttyname 0 path"
```
Function
umask invokes the umask callable service to change your process’s file mode creation mask. The file mode creation mask is used by the security package (RACF) to turn off permission bits in the mode parameter specified. Bit positions that are set in the file mode creation mask are cleared in the mode of the created file.

Parameters
mask
A permission bit mask that you specify as a three-digit number. Each digit must be in the range 0 to 7, and all digits must be specified. For more information on permissions, see Appendix B.

Usage notes
1. The umask service changes the process’s file creation mask. This mask controls file permission bits that are set whenever the process creates a file. File permission bits that are turned on in the file creation mask are turned off in the file permission bits of files created by the process. For example, if a call to the open service, BPX1OPN, specifies a mode argument with file permission bits, the process’s file creation mask affects that argument: bits that are on in the mask will be turned off in the mode argument, and therefore in the mode of the created file.
2. Only the file permission bits of new mask are used.

Example
To create a mask that sets read-write-execute permission on for the owner of the file and off for everyone else:

"umask 077"
uname

Function
uname invokes the uname callable service to obtain information about the system you are running on.

Parameters
stem
The name of a stem variable used to return the information. Upon return, stem.0 contains the number of variables returned. To access the information, you can specify a numeric value (see Appendix A) or the predefined variables beginning with U_ that derive the appropriate numeric value. For example, both stem.1 and stem.u_sysname access the name of the operating system.

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>U_MACHINE</td>
<td>The name of the hardware type on which the system is running</td>
</tr>
<tr>
<td>U_NODENAME</td>
<td>The name of this node within the communication network</td>
</tr>
<tr>
<td>U_RELEASE</td>
<td>The current release level of this implementation</td>
</tr>
<tr>
<td>U_SYSNAME</td>
<td>The name of this implementation of the operating system (z/OS)</td>
</tr>
<tr>
<td>U_VERSION</td>
<td>The current version level of this release</td>
</tr>
</tbody>
</table>

Example
"uname uts."
unlink

Function

unlink invokes the unlink callable service to remove a directory entry.

Parameters

pathname
A pathname for the directory entry. The directory entry could be identified by a pathname for a file, the name of a hard link to a file, or the name of a symbolic link.

Usage notes

1. If the name specified refers to a symbolic link, the symbolic link file named by pathname is deleted.
2. If a file is deleted (that is, if the unlink service request is successful and the link count becomes zero), the contents of the file are discarded, and the space it occupied is freed for reuse. However, if another process (or more than one) has the file open when the last link is removed, the file is not removed until the last process closes it.
3. When the unlink service is successful in removing the directory entry and decrementing the link count, whether or not the link count becomes zero, it returns control to the caller with RETVAL set to 0. It updates the change and modification times for the parent directory, and the change time for the file itself (unless the file is deleted).
4. Directories cannot be removed using unlink. To remove a directory, refer to "rmdir" on page 150.

Example

In the following example, assume that file was assigned a value earlier in the exec:
"unlink (file)"
Function

unmount invokes the umount callable service to unmount a file system; that is, it removes a file system from the file hierarchy. You must be a superuser to unmount a file system.

Parameters

name
The name of the file system to be unmounted, specified as the name of an HFS data set. You must specify the HFS data set name as a fully qualified name in uppercase letters. Do not enclose the data set name in single quotes.

flags
The unmount options, expressed as a numeric value. You can specify a numeric value (see Appendix A) or the predefined variable beginning with MTM_ used to derive the appropriate numeric value:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>MTM_DRAIN</td>
<td>An unmount drain request. All uses of the file system are normally ended before the file system is unmounted.</td>
</tr>
<tr>
<td>MTM_FORCE</td>
<td>An unmount force request. The file system is unmounted immediately, forcing any users of the named file system to fail. All data changes made up to the time of the request are saved. If there is a problem saving data, the unmount continues and the data may be lost. So that data will not be lost, you must issue an unmount immediate request before an unmount force request.</td>
</tr>
<tr>
<td>MTM_IMMED</td>
<td>An unmount immediate request. The file system is unmounted immediately, forcing any users of the named file system to fail. All data changes made up to the time of the request are saved. If there is a problem saving data, the unmount request fails.</td>
</tr>
<tr>
<td>MTM_NORMAL</td>
<td>A normal unmount request. If no one is using the named file system, the unmount request is done. Otherwise, the request is rejected.</td>
</tr>
<tr>
<td>MTM_RESET</td>
<td>A reset unmount request. This stops a previous unmount drain request.</td>
</tr>
<tr>
<td>MTM_REMOUNT</td>
<td>Unmounts the file system, changes the mount mode, and remounts the file system. A read/write mount mode changes to read-only. A read-only mount mode changes to read/write.</td>
</tr>
<tr>
<td>MTM_SAMEMODE</td>
<td>Remounts the file system without changing the mount mode. This can be used to regain the use of a file system that has I/O errors.</td>
</tr>
</tbody>
</table>

Usage notes

1. A file system that has file systems mounted on it cannot be unmounted. Any child file systems must be unmounted first.
2. A reset request can stop only an unmount service drain request. There is no effect if it is issued when there is no umount request outstanding.
Example
To request a normal unmount of the file system HFS.USR.CRISP:

"unmount HFS.USR.CRISP" mtm_normal
unquiesce

Function
unquiesce invokes the unquiesce callable service to unquiesce a file system, making the files in it available for use again. You must be a superuser to use this function.

Parameters
name
The name of the file system to be unquiesced. You must specify an HFS data set name as a fully qualified name in uppercase letters. Do not enclose the data set name in single quotes.

flag
A number specifying the type of unquiesce:
0  Normal unquiesce
1  Forced unquiesce. Request is allowed even if the requester process is not the process that made the quiesce request.

Usage notes
An unquiesce service makes a file system available for use again following a previous quiesce request.

Example
To request a normal unquiesce of the file system HFS.USR.ELIZAB:
"unquiesce HFS.USR.ELIZAB 0"
utime

Function
The utime service is used to set the access and modification times of a file.

Parameters

pathname
A pathname for the file.

atime
A numeric value for the new access time for the file, specified as POSIX time (seconds since the Epoch, 00:00:00 1 January 1970).

mtime
A numeric value for the new modification time for the file, specified as POSIX time (seconds since the Epoch, 00:00:00 1 January 1970).

Example
In the following example, assume that file, atm, and mtm were assigned values earlier in the exec:

"utime (file)" atm mtm
**Function**

`wait` invokes the `wait` callable service to obtain the status of any child process that has ended or stopped. You can use the `wait` service to obtain the status of a process that is being debugged with the `ptrace` facilities. The term `child` refers to a child process created by a fork as well as a process attached by `ptrace`.

**Parameters**

**stem**

The name of a stem variable used to return the information. Upon return, `stem.0` contains the number of variables returned. To access the information in the stem variables, you can specify a numeric value (see Appendix A) or the predefined variables beginning with `W_` that derive the appropriate numeric value.

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>W_CONTINUED</td>
<td>Process continued from stop</td>
</tr>
<tr>
<td>W_EXITSTATUS</td>
<td>The exit status of the child process</td>
</tr>
<tr>
<td>W_IFEXITED</td>
<td>The child process ended normally</td>
</tr>
<tr>
<td>W_IFSIGNALED</td>
<td>The child process ended because of a signal that was not caught</td>
</tr>
<tr>
<td>W_IFSTOPPED</td>
<td>Wait if the child process is stopped</td>
</tr>
<tr>
<td>W_STAT3</td>
<td>Byte 3 of the <code>BPXYWAST</code> macro. See <a href="#">BPXYWAST</a> — Map the Wait Status Word in <a href="#">Z/OS UNIX System Services Programming: Assembler Callable Services Reference</a>.</td>
</tr>
<tr>
<td>W_STAT4</td>
<td>Byte 4 of the <code>BPXYWAST</code> macro. See <a href="#">BPXYWAST</a> — Map the Wait Status Word in <a href="#">Z/OS UNIX System Services Programming: Assembler Callable Services Reference</a>.</td>
</tr>
<tr>
<td>W_STOPSIG</td>
<td>The signal number that caused the child process to stop</td>
</tr>
<tr>
<td>W_TERMSIG</td>
<td>The signal number that caused the child process to end</td>
</tr>
</tbody>
</table>

**Usage notes**

1. The `wait` service suspends execution of the calling thread until one of the requested child or debugged processes ends or until it obtains information about the process that ended. If a child or debugged process has already ended but its status has not been reported when `wait` is called, the routine immediately returns with that status information to the caller.

2. If the `WUNTRACED` option is specified, the foregoing also applies for stopped children or stopped debugged processes.

3. The `wait` service always returns status for stopped `debugged` processes, even if `WUNTRACED` is not specified.

4. If status is available for one or more processes, the order in which the status is reported is unspecified.
**wait**

**Note:** A debugged process is one that is being monitored for debugging purposes with the ptrace service.

**Example**
See "Set up a signal to enforce a time limit for a program" on page 212 for an example of signal coding that interprets the stem this returns:
"wait wstat."
Function

`waitpid` invokes the wait callable service to obtain the status of a child process that has ended or stopped. You can use the wait service to obtain the status of a process that is being debugged with the ptrace facilities. The term child refers to a child process created by a fork as well as a process attached by ptrace.

Parameters

**pid**
A numeric value indicating the event the caller is waiting upon:
- A value greater than zero is assumed to be a process ID. The caller waits for the child or debugged process with that specific process ID to end or to stop.
- A value of zero specifies that the caller is waiting for any children or debugged processes with a process group ID equal to the caller’s to end or to stop.
- A value of −1 specifies that the caller is waiting for any of its children or debugged processes to end or to stop.
- If the value is negative and less than −1, its absolute value is assumed to be a process group ID. The caller waits for any children or debugged processes with that process group ID to end or to stop.

**stem**
The name of a stem variable used to return the information. Upon return, `stem.0` contains the number of variables returned. To access the information in the stem variables, you can use numeric values (see Appendix A), or the predefined variables beginning with W_ (see their description in “wait” on page 195).

**options**
A numeric value or its equivalent predefined variable beginning with W_ that indicates the wait options for this invocation of the wait service. These options can be specified separately or together. (For the numeric values, see Appendix A)

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>0</td>
<td>Wait for the child process to end (default processing)</td>
</tr>
<tr>
<td>W_NOHANG</td>
<td>The wait service does not suspend execution of the calling process if status is not immediately available for one of the child processes specified by Process_ID.</td>
</tr>
<tr>
<td>W_UNTRACED</td>
<td>The wait service also returns the status of any child processes specified by Process_ID that are stopped, and whose status has not yet been reported since they stopped. If this option is not specified, the wait service returns only the status of processes that end.</td>
</tr>
</tbody>
</table>
Usage notes

1. Use `waitpid` when you want to wait for a specified child process. The `pid` argument specifies a set of child processes for which status is requested; `waitpid` returns the status of a child process from this set.

2. The `waitpid` service suspends execution of the calling thread until one of the requested child or debugged processes ends or until it obtains information about the process that ended. If a child or debugged process has already ended but its status has not been reported when `waitpid` is called, the routine immediately returns with that status information to the caller.

3. If the WUNTRACED option is specified, the foregoing also applies for stopped children or stopped debugged processes. A debugged process is one that is being monitored for debugging purposes with the ptrace service.

4. The wait service always returns status for stopped debugged processes, even if WUNTRACED is not specified.

5. If status is available for one or more processes, the order in which the status is reported is unspecified.

Example

See "Set up a signal to enforce a time limit for a program" on page 212 for an example of signal coding that interprets the stem this returns:

"waitpid -l wst. 0"
Function

**write** invokes the write callable service to copy data to a buffer and then write it to an open file. The number of bytes written is returned in RETVAL.

Parameters

- **fd**  The file descriptor (a number) for a file.
- **variable**  The name of the variable that is to store the data to be written to the file.
- **length**  The number of bytes to be written to the file identified by **fd**. If you want a length longer than 4096 bytes, specify the **length** parameter. If you do not specify **length**, the length of **variable** is used, up to a maximum length of 4096 bytes. A **variable** longer than 4096 bytes is truncated to 4096, and RC is set to 4.

Usage notes

1. Within the variable, you can use the predefined variables beginning with `ESC_` the same way you use C language escape sequences to avoid code page dependence with some control characters. For example:
   
   ```
   buf='line 1' || esc_n
   ```

   appends a newline character to the string 'line 1'.

<table>
<thead>
<tr>
<th>Code</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>ESC_A</td>
<td>Alert (bell)</td>
</tr>
<tr>
<td>ESC_B</td>
<td>Backspace</td>
</tr>
<tr>
<td>ESC_F</td>
<td>Form feed (new page)</td>
</tr>
<tr>
<td>ESC_N</td>
<td>Newline</td>
</tr>
<tr>
<td>ESC_R</td>
<td>Carriage return</td>
</tr>
<tr>
<td>ESC_T</td>
<td>Horizontal tab</td>
</tr>
<tr>
<td>ESC_V</td>
<td>Vertical tab</td>
</tr>
</tbody>
</table>

2. Return codes:
   - 4 indicates one of these:
     - If **length** was specified, the number of characters specified by **length** is not the same as the length of **variable**. The data is truncated or padded as required. The characters used for padding are arbitrarily selected.
     - If **length** was not specified, 4 indicates that a variable longer than 4096 bytes was truncated to 4096.
   - −24 indicates that storage could not be obtained for the buffer.

3. **File offset**: If **fd** specifies a regular file or any other type of file on which you can seek, the write service begins writing at the file offset associated with that file descriptor. A successful write operation increments the file offset by the
number of bytes written. If the incremented file offset is greater than the previous length of the file, the file is extended: the length of the file is set to the new file offset.

If the file descriptor refers to a file on which you cannot seek, the service begins writing at the current position. No file offset is associated with such a file.

If the file was opened with the “append” option, the write routine sets the file offset to the end of the file before writing output.

4. **Number of bytes written:** Ordinarily, the number of bytes written to the output file is the number you specify in the `length` parameter. (This number can be zero. If you ask to write zero bytes, the service simply returns a return value of zero, without attempting any other action.)

If the `length` you specify is greater than the remaining space on the output device, fewer bytes than you requested are written. When at least 1 byte is written, the write is considered successful. The return value shows the number of bytes written. An attempt to write again to the same file, however, causes an ENOSPC error unless you are using a pseudoterminal. With a pseudoterminal, if there is not enough room in the buffer for the whole write, the number of bytes that can fit are written, and the number of bytes written is returned. However, on the next write (assuming the buffer is still full) there is a block or EAGAIN is returned, depending on whether the file was opened blocking or nonblocking.

Similarly, fewer bytes are written if the service is interrupted by a signal after some but not all of the specified number of bytes are written. The return value shows the number of bytes written. If no bytes were written before the routine was interrupted, the return value is −1 and an EINTR error is reported.

5. The write service causes signal **SIGTTOU** to be sent under all the following conditions:
   - The process is attempting to write to its controlling terminal.
   - TOSTOP is set as a terminal attribute.
   - The process is running in a background process group.
   - The **SIGTTOU** signal is not blocked or ignored.
   - The process is not an orphan.

If all the conditions are met, **SIGTTOU** is sent.

6. Write requests to a pipe (FIFO) are handled in the same as write requests to a regular file, with the following exceptions:
   - There is no file offset associated with a pipe; each write request appends to the end of the pipe.
   - If the size of the write request is less than or equal to the value of the PIPE_BUFF variable (described in the pathconf service), the write is guaranteed to be atomic. The data is not interleaved with data from other write processes on the same pipe. If the size of the write request is greater than the value of PIPE_BUFF, the data can be interleaved, on arbitrary boundaries, with writes by other processes, whether or not the O_NONBLOCK flag is set.

**Example**

In the following example, assume that `fd` and `buf` were assigned values earlier in the exec:

```
"write" fd "buf"
```
Function

`writefile` invokes the open, write, and close callable services to write or append text files, with lines up to 1024 characters long.

Parameters

**pathname**
A pathname for the file.

**mode**
A three- or four-digit number, corresponding to the access permission bits. Each digit must be in the range 0–7, and at least three digits must be specified. For more information on permissions, see [Appendix B](#).

**stem**
The name of a stem variable that contains the information to be written to the file. `stem.0` is set to the number of lines to be written. `stem.1` through `stem.n`, where `n` is the total number of variables written, each contain a line to be written. A newline is written to the file following each line.

Within the stem, you can use the predefined variables beginning with ESC_ the same way you use C language escape sequences to avoid code page dependence with some control characters. See the usage notes for "write" on page 199.

**append**
An optional flag to indicate that the information is to be appended to the file (the file must already exist), as follows:
- 0 Do not append. (This is the default value.)
- 1 Append the information to the file.

Usage notes

File I/O stops when `writefile` sets a return code. `writefile` can set the following return codes:

- **4** A line is longer than 1024 characters.
- **8** A write was attempted, but failed. RETVAL, ERRNO, and ERRNOJR contain the return values from the write callable service.

For further usage notes, see:
- "close" on page 46
- "open" on page 129
- "write" on page 199

Example

In the following example, assume that `fname` and the stem `file` were assigned values earlier in the exec:

"writefile (fname) 600 file."
writefile
Chapter 4. Examples: Using syscall commands

The examples in this chapter are provided to assist you with coding REXX programs that use z/OS UNIX syscall commands.

The first three examples can run in TSO/E, batch, or from the z/OS shells. They begin with the call `syscalls('ON')`.

**Read the root directory into a stem and print it**

This example prints the contents of the root directory to standard output:

```rexx
/* rexx */
call syscalls 'ON'
address syscall
'readdir / root.'
do i=1 to root.0
   say root.i
end
```

The following line saves the results from the previous example in a text file:

```rexx
'writefile /u/schoen/root.dir 777 root.'
```

**Open, write, and close a file**

```rexx
/* rexx */
call syscalls 'ON'
address syscall
path="/u/schoen/my.file'
'open' path,
   O_rdwr+O_creat+O_trunc,
   660
if retval=-1 then
   do
   say 'file not opened, error codes' errno errnojr
   return
   end
fd=retval
rec='hello world'|| esc_n
'write' fd 'rec' length(Rec)
if retval=-1 then
   say 'record not written, error codes' errno errnojr
'close' fd
```

**Open a file, read, and close it**

```rexx
/* rexx */
call syscalls 'ON'
address syscall
path="/u/schoen/my.file'
'open' (path)
   0 rdonly,
   060
if retval=-1 then
   do
   say 'file not opened, error codes' errno errnojr
   return
   end
fd=retval
'read' fd 'bytes 80'
if retval=-1 then
```
say 'bytes not read, error codes' errno errnojr
else
  say bytes
close' fd

Display the working directory and list a specified directory

This REXX program runs in the z/OS shells; it uses both the SH and SYSCALL environments. The program identifies your working directory and lists the contents of a directory that you specify as a parameter or after a prompt.

```rexx
/*
parse arg dir
address syscall 'getcwd cwd'
say 'current directory is' cwd
if dir=' ' then
do
  say 'enter directory name to list'
pull dir
end
'ls -l' dir
return
*/
```

Parse arguments passed to a REXX program: the getopt function

The following simple utility function is used by some of the examples to parse the arguments to a REXX program that is run from a shell. This function parses the stem __argv for options in the format used by most POSIX commands.

```rexx
/*******************************************************************************/
/* Function: getopt */
/* This function parses __argv, stem for options in the format */
/* used by most POSIX commands. This supports simple option */
/* letters and option letters followed by a single parameter. */
/* The stem OPT. is setup with the parsed information. The */
/* options letter in appropriate case is used to access the */
/* variable: op='a'; if opt.op=1 then say 'option a found' */
/* or, if it has a parameter: */
/*   op='a'; if opt.op<>' then say 'option a has value' opt.op */
/* */
/* Parameters: option letters taking no parms */
/*   option letters taking 1 parm */
/* */
/* Returns: index to the first element of __argv. that is not */
/* an option. This is usually the first of a list of files. */
/* A value of 0 means there was an error in the options and */
/* a message was issued. */
/* */
/* Usage: This function must be included in the source for the exec */
/* */
/* Example: the following code segment will call GETOPTS to parse */
/* the arguments for options a, b, c, and d. Options a */
/* and b are simple letter options and c and d each take */
/* one argument. It will then display what options were */
/* specified and their values. If a list of files is */
/* specified after the options, they will be listed. */
/* */
/* parse value 'a b c d' with, */
/* lca lcb lcc lcd . */
/* argx=getopts('ab','cd') */
/* if argx=0 then exit 1 */
/* if opt.a<>0 then */
/* say 'No options were specified' */
/* else */
/* do */
/* if opt.lca<>' then say 'Option a was specified' */```
/**
  if opt.lcb<>' then say 'Option b was specified'
  */
/**
  if opt.lcc<>' then say 'Option c was specified as' opt.lcc
  */
/**
  if opt.lcd<>' then say 'Option d was specified as' opt.lcd
  */
/**
  end
  */
/**
  if __argv.0>=argx then
  */
/**
  say 'Files were specified:'
  */
/**
  else
  */
/**
  say 'Files were not specified'
  */
/**
  do i=argx to __argv.0
  */
/**
  say __argv.i
  */
/**
  end
  */
/**
  */
/**
*******************************************************************************/
getopts: procedure expose opt. __argv.
  parse arg arg0,arg1
  argc=__argv.0
  opt.=' '
  opt.0=0
  optn=0
  do i=2 to argc
    if substr(__argv.i,1,1)<>'-' then leave
    if __argv.i='-' then leave
    opt=substr(__argv.i,2)
    do j=1 to length(opt)
      op=substr(opt,j,1)
      if pos(op,arg0)>0 then
        do
          opt.op=1
          optn=optn+1
        end
        else
          if pos(op,arg1)>0 then
            do
              if substr(opt,j+1)<>' then do
                opt.op=substr(opt,j+1)
              j=length(opt)
            end
            else
              do
                i=i+1
              if i>argc then
                do
                  say 'Option' op 'requires an argument'
                  return 0
                end
              opt.op=__argv.i
            end
            optn=optn+1
          end
          else
            do
              say 'Invalidoption = ' op
            say 'Valid options are:' arg0 arg1
            return 0
          end
        end
      end
    end
  opt.0=optn
  return i

Chapter 4. Examples: Using syscall commands  205
Count newlines, words, and bytes

This is an example of a REXX program that can run as a shell command or filter. It is a REXX implementation of the wc (word count) utility supporting the options −c, −w, and −l.

The program uses open, close, and EXECIO. To read standard input, it accesses the file /dev/fd0.

```rexx
/* rex */
parse value 'l w c' with,
    lcl lcw lcc .
argx=getopts('lwc')
/* parse options */
if argx=0 then return 1
/* return on error */
if opt.0=0 then
/* no opts, set defaults */
    parse value '1 1 1' with,
        opt.lcl opt.lcw opt.lcc .
if __argv.0>argx then
/* multiple files specified */
single=0
else if __argv.0=argx then
/* one file specified */
single=1
else do
/* no files specified, use stdin */
single=2
    __argv.argx='/dev/fd0'
    __argv.0=argx
end
parse value '0 0 0' with,
    twc tcc tlc .
/* clear total counters */
address syscall
do i=argx to __argv.0
/* loop through files */
    fi=__argv.i
    /* get file name */
    parse value '0 0 0' with,
        wc cc lc .
    "open (fi)\" _rdonly 000
    fd= retval
    if fd=-1 then
        /* open failed */
    do
        say 'unable to open' fi
        iterate
    end
end
/* loop reading 1 line at a time */
do forever
    address mvs 'execio diskr' fd 'stem LN.'
    if rc<>0 | ln.0=0 then leave
    if opt.lcw=1 then
        wc=wc+words(ln.1)
    if opt.lcc=1 then
        cc=cc+length(ln.1)+1
    if opt.lcl=1 then
        lc=lc+1
end
'close' fd
/* close file */
twc=twc+wcc
/* accumulate total words */
tlc=tlc+lc
/* accumulate total lines */
tcc=tcc+ccc
/* accumulate total chars */
if opt.lcw<>1 then wc=''
    else wc=right(wc,7)
if opt.lcl<>1 then lc=''
    else lc=right(lc,7)
if opt.lcc<>1 then cc=''
    else cc=right(cc,7)
if single=2 then fi=''
    /* if stdin used clear filename */
say lc wc cc ' 'fi
/* put out counts message */
```
if single=0 then /* if multiple files specified */
do /* format and output totals line */
  if opt.lcw<>1 then twc='else twc=right(twc,7)
  if opt.lcl<>1 then tlc='else tlc=right(tlc,7)
  if opt.lcc<>1 then tcc='else tcc=right(tcc,7)
say tlc twc tcc 'total'
end
return 0

Obtain information about the mounted file system

This REXX program uses getmntent and statfs to list all mount points, the name of the mounted file system, and the available space in the file system. This sample REXX program, as coded, must be run from a shell.

/* rexx */
address syscall
numeric digits 12
'getmntent m.'
do i=1 to m.0
  'statfs' m.mnte_fsname.i 's.'
  j=s.stfs_avail = s.stfs_blocksize
  if length(m.mnte_path.i)>20 then
    say m.mnte_path.i 'strip(m.mnte_fsname.i) 'j'
  else
    say left(m.mnte_path.i,20) strip(m.mnte_fsname.i) 'j'
end

Mount a file system

This REXX program uses mount to mount a file system, an action that requires superuser authority. The name of this program is mountx, and it is in /samples.

The syntax is:

mountx pathname  fsn options

where pathname is the name of the directory where the file system is to be mounted, and fsn is the name of the HFS data set. The options are:

-p parm Parameter data, in the form of a single string.
-r Mount file system as read-only.
-t type File system type (for example, HFS). The default file system type to HFS. If the file system type is HFS, the program changes the file system name to uppercase.

The pathname and the file system name can be specified in any order. stat is used to determine which name is the pathname.

For the getopt function called in this program, see “Parse arguments passed to a REXX program: the getopt function” on page 204.
ix=getopts('r','pt')
/*************************************************************************/
/* If -r specified, mount file system read-only. */
/*************************************************************************/
if opt.lcr<>'' then
  mtm=mtm_rdonly
else
  mtm=mtm_rdwr
/*************************************************************************/
/* If -t 'name' specified, direct mount request to file system 'name'. */
/* Otherwise, direct mount request to file system named HFS. */
/*************************************************************************/
if opt.lct<>'' then
  type=translate(opt.lct)
else
  type='HFS'
/*************************************************************************/
/* Complain if required parameters are missing. */
/*************************************************************************/
if __argv.0<>ix+1 then
  do
    say 'Pathname and file system name are required, and '
    say 'they must follow the options: MOUNT <options> <pathname> <fsn>'
    return 1
  end
/*************************************************************************/
/* Direct function calls to REXX z/OS UNIX services. */
/*************************************************************************/
address syscall
/*************************************************************************/
/* Determine which of the parameters is the mount point name */
/* (it must be a pathname), and which is the file system to be */
/* mounted. */
/*************************************************************************/
'stat (__argv.ix) st.'
if st.st_type<>s_isdir then
  do
    fsn=__argv.ix
    ix=ix+1
    path=__argv.ix
  end
else
  do
    path=__argv.ix
    ix=ix+1
    fsn=__argv.ix
  end
'stat (path) st.'
if st.st_type<>s_isdir then
  do
    say "Can't figure out pathname, neither name is a directory:"
    say path
    say fsn
    return 1
  end
/*************************************************************************/
/* HFS file system requires mounted file systems to be data sets, */
/* so translate the file system name to upper case. */
/*************************************************************************/
if type='HFS' then
  fsn=translate(fsn)
/*************************************************************************/
/* Mount the file system. */
Unmount a file system

This REXX program uses **unmount** to unmount a file system, an action that requires superuser authority. The name of this program is **unmountx**, and it is in `/samples`.

The syntax is:

```
unmountx name
```

or

```
unmountx -t filesystem
```

where:
- *name* is the pathname where the file system is mounted, or the name of an HFS data set.
- *filesystem* is the type name of the physical file system (PFS).

The program assumes that the case of the requested file system name is entered correctly in uppercase. If the unmount fails, it folds the file system name to uppercase and retries the unmount.

```
/* Direct commands to REXX z/OS UNIX services. */
address syscall

/* Verify that exactly one operand or a pfs type was specified. */
if __argv.0=3 & __argv.2='-t' then
   return fstype(__argv.3)
if __argv.0<>2 then
   do
```

say 'Syntax: unmount <name> or unmount -t <filesystype>'
return 2
end
/************************************************************************/
/* Determine if the name is a pathname. If so, determine file system */
/* name via stat(). Otherwise, use the name as entered. */
/************************************************************************/
'stat (argv.2) st.'
if retval =0 & st.st_type=s_isdir then
  do
    getmntent mnt. x2d(st.st_dev)
    fsn=mnt.mnte_fsname.1
  end
else
  fsn=argv.2
/************************************************************************/
/* Unmount the file system, trying both the name as entered */
/* and the name uppercased, since the HFS file system requires */
/* mounted file systems to be data sets. */
/* */
/* Return code Return code Meaning */
/* from system to caller */
/* Not -1 0 Success, file system unmount complete */
/* -1 1 Error, explained by ERRNO and reason code*/
/************************************************************************/
"unmount (fsn)" mtm_normal /* unmount name as specified */
if retval=0 then
  do
    fsn=translate(fsn) /* if fails, upcase name and retry */
    "unmount (fsn)" mtm_normal
  end
if retval<>-1 then
  do
    say 'Unmount complete for' fsn
    return 0
  end
else
  do
    say 'Unmount failed:'
    say ' Error number was' erroxx('x2d(errno)')
    say ' Reason code was ' right(errnojr,8,0)
  end
/************************************************************************/
/* Unmount all file systems for a PFS type */
/************************************************************************/
fstype:
  arg name.
  /* make upper case */
  /* Loop through mount table until unable to do any unmounts. */
  /* This handles cascaded mounts. */
/************************************************************************/
do until didone=0
  didone=0
  'getmntent m.'
  do i=1 to m.0
    if m.mnte_fstype.i=name then
      do
        "unmount (m.mnte_fsname.i)" mtm_normal
        if retval<>-1 then
          do
            didone=1
            say 'unmounted:' m.mnte_fsname.i
            end
        end
      end
    end
  end
/************************************************************************/
Run a shell command and read its output into a stem

This REXX program runs the ls shell command and reads the output from the command into a stem. The program uses pipe, close, and EXECIO. It accesses /dev/fd, where \( n \) is a number that the exec concatenates to /dev/fd.

**Note:** You can use this example to trap output from commands when the output is relatively small (less than 100KB). For command output that could be larger, you should use the spawn service.

```rexx
/* rexx */
address syscall 'pipe p.' /* make a pipe */
'ls>/dev/fd' || p.2 /* run the ls command and redirect output to the write end of the pipe */
address syscall 'close' p.2 /* close output side */
address mvs 'execio * diskr' p.1 '(stem s.' /* read data in pipe */
do i=1 to s.0 /* process the data */
say s.i
say s.1
end
```

Print the group member names

This REXX program uses getgrgid, getgrnam, and write to print the names of the users that are connected to a group. The group can be specified as either a GID or a group name.

```rexx
/* rexx */
arg group .
address syscall
if datatype(group,'W')=1 then /* getgrgid if GID specified */
   'getgrgid (group) gr.'
else /* getgrnam if group name specified */
   'getgrnam (group) gr.'
if retval=0 then /* check for error */
do
   say 'No information available for group' group
   return 1
end
say 'Connected users for group' strip(gr.gr_name)'('gr.gr_gid')'
j=1
do i=gr_mem to gr.0
   buf=' '
```

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Obtain information about a user

This REXX program prints information from the user database for a user. The user can be specified as either a UID or a user name. The program uses getpwuid, getpwnam, getgrgid, getgroupsbyname, and write.

```
/* rexx */
arg user .
address syscall
if datatype(user,'W')=1 then
  'getpwuid (user) pw.' /* use getpwuid if UID specified */
else
  'getpwnam (user) pw.' /* use getpwnam if user name specified */
if retval<=0 then /* check for error */
  do
    say 'No information available for user' user
  return 1
end
say 'Information for user' strip(pw.pw_name)'('pw.pw_uid')'
say ' Home directory:' strip(pw.pw_dir)
say ' Initial program:' strip(pw.pw_shell)
'getgrgid' pw.pw_gid 'gr.'
if retval<=0 then /* check for error */
  do
    say ' Group information not available'
  return 1
end
say ' Primary group: ' strip(gr.gr_name)'('gr.gr_gid')'
'getgroupsbyname' pw.pw_name 's.'
if retval<=0 then /* check for error */
  do
    say ' Supplemental group information not available'
  return 1
end
say ' Supplemental GIDs:'
do i=1 to s.0
  buf=right(s.i,12)
  if i//5=0 | i=s.0 then
    buf=buf || esc_n /* write newline every 5 groups or at end */
  'write 1 buf'
end
return 0
```

Set up a signal to enforce a time limit for a program

This REXX program runs /bin/ls to list files in the /bin directory, and sets up a signal that enforces a time limit of 10 seconds for the program to run:

```
/* REXX */
address syscall
/* initialize file descriptor map (see note 1) */
fd.0=-1
fd.2=-1
'creat /tmp/dirlist 600'
fd.1=retval
/* initialize parameter stem (see note 2) */
parm.1="/bin/ls"
```

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parm.2='-l'
parm.3='/bin'
parm.0=3
/*
spawn
new
process
(see
note
3)
*/
'spawn
/bin/ls
3
fd.
parm. __environment.'
pid=retval
/*
set
up
signals
to
wait
up
to
10
seconds
(see
note
4)
*/
call
syscalls 'SIGON'
'sigaction'
sigalrm
'sig_cat' 0 'oldh oldf'
'sigprocmask'
sig Unblock
'sigaddset(sigsetempty(),sigalrm) 'mask'
'alarm' 10 /* set alarm */
'swaitpid (pid) st. 0' /* wait for process term or alarm */
srv=retval
'alarm 0'
/* make sure alarm is now off */
if srv=-1 then /* if alarm went off */
do
'kill' pid sigkill /* cancel child process */
'swaitpid (pid) st. 0' /* wait for completion */
end
call
syscalls 'SIGOFF' /* turn off signals */
/* determine process status code (see note 5) */
select
  when st.w_ifexited then
    say 'exited with code' st.w_exitstatus
  when st.w_ifsignaled then
    say 'terminated with signal' st.w_termsig
  when st.w_ifstopped then
    say 'stopped by signal' st.w_stopsig
end
return
sigsetempty: return copies(0,64)
sigaddset: return overlay(1,arg(1),arg(2))

Notes:

1. The file descriptor map is set up so that the file descriptor for the new file being
   created is remapped to file descriptor 1 (standard output) for the new process.
   File descriptors 0 and 2 will not be opened in the new process.

2. The first parameter is set to the pathname for the file being spawned. Additional
   parameters are set in the format the program expects. In this case, they specify
   a long directory listing for the /bin directory.

3. The new process is spawned to run /bin/ls. File descriptors greater than or
   equal to 3 are not available to the new process. fd.0 to fd.2 are used in
   remapping file descriptors from the parent. The current environment for the
   parent is passed to the new process.

4. The syscalls function is called to enable signals. If this program were to be
   exec’d (run from the shell), this would not be necessary, and, similarly, the call
   later on to turn off signals would also be unnecessary.

   sigaction is used to set the action for sigalrm to be caught. sigprocmask is
   used to unblock sigalrm. This call also uses sigaddset and sigsetempty to
   create a signal mask with the sigalrm bit.

   The alarm is set by using the alarm service, and the process waits for
   completion of the child or the alarm. "ALARM 0" is used just to make sure the
   alarm is off.

5. A SELECT instruction is used on the status stem returned by the waitpid
   service. This determines if the process was terminated by a signal, if it exited,
   or if the process is stopped. If it exited, the exit status is available; otherwise,
   the signal number is available.
List the ACL entries for a file

This example will show the access ACL entries for a pathname given as a parameter:

```rexx
/* REXX */
parse arg path
address syscall
'aclinit acl' /* init variable ACL to hold acl */
'aclget acl (path) acl_type_access /* get the file access acl */
do i=1 by 1 /* get each acl entry */
   'aclgetentry acl acl.' i
   if rc<0 | retval=-1 then leave /* error, assume no more */
   parse value '- - -' with pr pw px
   if acl.acl_read=1 then pr='R' /* set rwx permissions */
   if acl.acl_write=1 then pw='W'
   if acl.acl_execute=1 then px='X'
   aclid=acl.acl_id /* get uid or gid */
   /* determine acl type */
   if acl.acl_entry_type=acl_entry_user then type='UID='
      else if acl.acl_entry_type=acl_entry_group then type='GID='
      else type='???'
   say pr || pw || px type || aclid /* print line */
end
'aclfree acl' /* must free acl buffer */
```

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Chapter 5. z/OS UNIX REXX functions

The z/OS UNIX REXX functions extend the REXX language on z/OS when it is used in the z/OS UNIX REXX environment. With the exception of bpxwunix() and syscalls(), these functions must be run in a z/OS UNIX environment. The z/OS UNIX REXX functions include functions for standard REXX I/O, and for accessing common file services and environment variables.

All of the functions are fully enabled for large files (>2GB).

All numbers that are used as input on the functions must be integers. The default precision for REXX is 9 digits. If arithmetic is used on large numbers, be sure to change your precision appropriately, using the NUMERIC DIGITS statement.

REXX I/O functions

The REXX input functions are charin() and linein(). The chars() and lines() functions determine if data remains in an input stream.

The REXX output functions are charout() and lineout().

The REXX stream() function controls the processing of I/O streams (file streams and process streams) and obtains their status. A number of commands are used within the stream() function to control stream processing.

Opening a stream implicitly

File streams can be opened implicitly or explicitly. You open a stream implicitly by using a pathname as the stream name for one of the six input or output functions. If the function is charout() or lineout(), the file is opened for output. If the file does not exist, it is created. For newly created files, the permission bits 0666 are applied to your process umask. The position for the first write is set to the end of file, unless the function call explicitly specifies a write location. If one of the four input functions is used, the file is opened for input.

When it is opened implicitly, a stream can be opened for both input and output. The input and output locations are independent of each other. If the stream is opened for both input and output, two file descriptors are used. The stream close command closes the stream and both file descriptors.

Opening a stream explicitly

You open a stream explicitly by using the stream() function.

To explicitly open a file stream, use the open command. The advantage of opening file streams explicitly is that the program can determine that the stream open failed. The program can also have several separate streams for the same file.

The stream open command fails if:
  - The file does not exist.
  - The pathname cannot be accessed for either input or output.

If the file cannot be opened, a message is written to stderr. The I/O function returns as if the stream is empty, at end of file, and the file cannot be extended.
When a stream is explicitly opened, the `stream()` function returns a string that is the name of the stream. This is the only name that can be used to identify that stream. Multiple opens for the same pathname open multiple streams, each with its own name.

**Process streams**

Process streams can only be opened explicitly, with the `popen` command. With process streams, you can run a shell command and provide its input or receive its output. To write the input for the command, use the `popen write` command. To receive the command's output, use the `popen read` command. The `stream()` function returns a string that is the name of the process stream. Use this name on the input functions for `popen read` or the output functions for `popen write`.

A process stream spawns `/bin/sh -c your_command`. The process inherits `stderr` and either `stdin` or `stdout`. If the process is opened for read, `stdout` for the process is a pipe; otherwise, `stdin` is a pipe. The shell completion code is returned as the result for the stream `pclose` command. `pclose` closes your end of the pipe and waits for the process to terminate.
Function

Runs a shell command and optionally:
- Provides its stdin
- Traps its stdout
- Traps its stderr
- Exports a set of environment variables

Parameters

cmd
The shell command that you want to run. The shell is run as `/bin/sh -c` followed by the string you specify as the command.

stdin
An optional argument, stdin is the name of a compound variable (stem) that contains input for the command.

stdin.0 must contain the number of lines that are to be redirected to the command. stdin.1, stdin.2, ... contain the lines.

stdin can also be specified as:
- The string STACK, if the input is on the stack.
- DD:ddname, if the input is to be read from an allocated DD.

If this argument is not specified, your current stdin file is passed to the shell for stdin.

stdout
An optional argument, stdout is the name of a compound variable (stem) that, on return, contains the normal output from the command. stdout.0 is the number of lines output by the command. stdout.1, stdout.2, ... contain the output lines.

stdout can also be specified as:
- The string STACK, if the output is to be returned on the stack
- DD:ddname, if the output is to be written to an allocated DD

If stdout is not specified, your current stdout file is passed to the shell for stdout.

stderr
An optional argument, stderr is the name of a compound variable (stem) that, on return, contains the error output from the command. stderr.0 is the number of lines output by the command. stderr.1, stderr.2, ... contain the output lines.

stderr can also be specified as:
bpwxunix()

- The string STACK, if the output is to be returned on the stack
- DD:ddname, if the output is to be written to an allocated DD

If this argument is not specified, your current stderr file is passed to the shell for stderr.

env

An optional argument, env is the name of a compound variable (stem) that contains environment variables for the command. env.0 must contain the number of environment variables to be passed to the command. env.1, env.2, ... contain the variables in the form variable_name=variable_value. If env is not specified, your current environment is passed to the shell for stdin.

Usage notes

1. The bpwxunix() function runs the shell by passing a single command similar to sh -c cmd. It does not run a login shell.
2. bpwxunix() can be used outside of the z/OS UNIX REXX environment (for example, in TSO/E). In this case, stdin, stdout, stderr, and environment variables are not inherited from the current process environment.
3. If the stdout or stderr stems are specified, they are filled as appropriate:
   - A return value in the range 0-255 is the exit status of the command.
   - A negative return value indicates failure, and is usually a signal number.
   - A number less than −1000 indicates a stop code.
4. The DD names that are used for input and output are processed by the standard REXX input and output services. They have the same restrictions as REXX in terms of the types of allocations they can handle.
5. STDOUT lines cannot exceed 2048 characters.

Example

To trap output from the ls command and display it:
call bpwxunix 'ls -l',out.
do i=1 to out.0
    say out.i
end

To send output from the above example to word count and print byte count:
call bpwxunix 'wc',out.,bc.
    parse var bc.1 . . count
    say 'byte count is' count

To trap output on stack and feed it to word count:
if bpwxunix('ls -l',stack)=0 then
    call bpwxunix 'wc',stack
Function

Returns a string of up to length characters read from the stream specified by name. The location for the next read is the current location increased by the number of characters returned. This function does no editing of the data.

Parameters

name
The name for the stream can be a pathname or a string that was returned from the stream open or popen commands. If name is omitted, the standard input stream is used.

start
For a persistent stream, specifies the byte number in the file where the read begins. start should not be specified for other types of streams.

length
The number of bytes returned by the function. If length bytes are not available in the stream, the function returns the number of bytes that are available and marks an error condition on the stream. For non-persistent streams, this function either blocks until length bytes are available, or returns with fewer bytes, depending on the file type and open flags. If length is 0, no characters are read, a null string is returned, and the read position is set based on the value of start.

Example

To read the next 256 characters:

say charin(file,,256)

To set the read location to the sixth 80-byte record:

call charin file,5*80+1,0
charout()

Function

Returns the number of characters remaining after an attempt to write string to the stream specified by name. The location for the next write is the current location increased by the number of characters written.

Parameters

name
The name for the stream can be a pathname or a string that was returned from the stream open or popen commands. If name is omitted, the standard output stream is used.

string
Data to be written to the stream specified by name.

start
For a persistent stream, specifies the byte number in the file where the write begins. start should not be specified for other types of streams.

Usage notes

1. If string is omitted, no data is written, 0 is returned, and the write location is set to the value start.
2. If start is also omitted, the write position is set to the end of file.

Example

To write the string to stdout:
call charout,'hello world'esc_n

To set the write position to end:
call charout file
Function

Returns the number of characters remaining in the input stream specified by name. For persistent streams, this is the number of characters between the current read location and the end of the stream. If the stream was created by the stream open command, chars(), while the process is active or bytes remain in the stream, returns either 1 or the number of bytes in the stream. After the process has terminated and the stream is empty, chars() returns a value of 0.

Parameters

name

The name for the stream, which can be a pathname or a string that was returned from the stream open or popen commands. If name is omitted, the standard input stream is used.

Example

To get the number of bytes in the stdin stream:

remainder=chars()
Function

Changes the mode for the specified pathname. It returns 0 if the mode for the specified pathname is changed; otherwise, it returns the system call error number.

Parameters

pathname
An absolute or relative pathname for a file.

operation
Specifies whether mode bits are to be set, added, or deleted:
  =  Set the mode bits. If an operation is not specified, = is used.
  +  Add the mode bits to what is currently set for the file.
  −  Remove the mode bits from what is currently set for the file.

mode
A string of octal digits for the new file mode.

Example

To set permissions for owner and group to read-write:
  call chmod file,660

To add read permission for other:
  call chmod file,+4
Function

Converts \textit{timestamp} to POSIX epoch time, and returns the time in seconds past the POSIX epoch (1/1/1970).

Parameters

\textit{timestamp}

A 14-character string in the form \texttt{mmdyyyyHMMS}.

Example

To set POSIX time for 4/21/99 7:15:00:

\texttt{say convd2e('04211999071500')}
Function

Returns the full pathname to the current directory, first changing it to newdirectory if the argument is supplied and you have access to that directory.

Parameters

newdirectory

An absolute or relative pathname for a directory.

Example

To change the current directory to /u/ehk:

call directory '/u/ehk'
environment()  

Function

Queries and alters environment variables. The stem __environment. is not altered through this service. That stem contains the environment variables on entry to the REXX program, and is available for your use. Alterations of the environment are used on subsequent calls to the stream popen command and ADDRESS SH.

Parameters

variable name
The name of the environment variable to operate on. If this is the only argument specified, the value of the variable is returned and the variable is not affected.

new value
A string to replace the value of variable name. The previous value of the variable is returned.

operation
An optional argument that specifies the operation to be performed. Only the first character is significant. The values can be:

- exists Tests the existence of the variable. The function returns 1 if the variable is defined, and 0 if it is not defined.
- delete Deletes the variable from the environment, if it exists. The function returns 0 if the variable is successfully deleted, and 1 if the variable is not defined.

Example

To get the value of the PATH environment variable:

```rexx
path=environment('PATH')
```

To reset PATH to the current directory:

```rexx
call environment 'PATH' , '.'
```

To delete the PATH environment variable:

```rexx
call environment 'PATH' , ',d'
```
Function

Returns the full pathname for the specified file. If the file does not exist, the function returns a null string.

Parameters

filename
A string that names a file

Example

To print the full pathname for the file myfile:

say exists('myfile')
getpass()

Function
Prints prompt on the controlling TTY and reads and returns one line of input with terminal echo suppressed.

Parameters

prompt
A string that prints on the controlling TTY

Example
To prompt for a password and read it:

psw=getpass('enter password')
Function

Returns one line or no lines from the stream specified in name, and sets the location for the next read to the beginning of the next line. The data is assumed to be text. The newline character is the line delimiter and is not returned. A null string is returned if no line is returned; this appears exactly the same as a null line in the file. Use chars() or lines() to determine if you are at end of file. Use stream() to determine if there is an error condition on the stream.

Parameters

ame
The name for the stream can be a pathname or a string that was returned from the stream open or popen commands. If name is omitted, the standard input stream is used.

start
For a persistent stream only, this argument can have the value 1, to begin reading at the beginning of the stream. No other value for start is supported. Do not specify start for other types of streams.

count
Specify 0 or 1 for the number of lines to be returned by the function. If one line is not available in the stream, the function returns a null string and marks an error condition on the stream. For non-persistent streams, this function either blocks until a line is available, or returns as though it is end of file, depending on whether any data remains in the file, on the file open flags, and on the type of file. If count is 0, no lines are read, a null string is returned, and, if start is 1, the read position is set to the beginning of the stream.

Example

To read the next line:

line=linein(file)

To read the first line in the file:

line=linein(file,1)
Function

Returns 1 line or 0 lines that are remaining to be written after an attempt to write
string to the stream specified by name. A newline character is written following
string. If an error occurs on the write, some data may be written to the stream, and
the function returns the value 1.

Parameters

name
The name for the stream can be a pathname or a string that was returned from
the stream open or popen commands. If name is omitted, the standard output
stream is used.

string
Data that is to be written to the stream specified by name. If string is omitted,
no data is written, 0 is returned, and the write position is set based on the value
of start. If both string and start are omitted, the write position is set to the
beginning of the file.

start
For a persistent stream, this argument can have the value 1, to begin writing at
the beginning of the stream. No other value for start is supported. Do not
specify start for other types of streams.

Example

to write the line to stdout:
call lineout , 'hello world'
Function

Returns 1 if data remains in the stream; otherwise it returns 0. Programs should check for a value of 0 or nonzero.

Parameters

name

The name for the stream can be a pathname or a string that was returned from the stream open or popen commands. If name is omitted, the standard input stream is used.

Example

To set more to nonzero, if stdin has data:

```plaintext
more=lines()
```
**Function**

Enables or disables the trapping of output from commands run using ADDRESS TSO, and returns the name of the variable in which trapped output is stored. If trapping is off, the word OFF is returned. Note that `outtrap` does not trap output for ADDRESS SH or any other command environment besides TSO. To trap shell command output, see "bpwunix()" on page 217.

**Parameters**

- **name**
  The name of a stem, a variable prefix used to contain command output, or the string OFF to turn off trapping.

- **max**
  The maximum number of lines to trap. If `max` is not specified, or if it is specified as * or a blank, the number of lines is set to 999 999 999.

- **catopt**
  Specify one of these:
  - **concat**
    Each command output trapping begins following the previous command output.
  - **noconcat**
    Output from each command is trapped starting with the variable concatenated with 1. Unused variables do not have their values altered.

For additional information about using the TSO command environment, see "The TSO command environment" on page 5.

**Example**

To trap TSO command output, use the stem OUT, rewriting the stem on each command:

```rexx
call outtrap 'out.', 'NOCONCAT'
```
procinfo()

Function

Retrieves information about one or more processes.

Parameters

pid
The process ID number of the process for which information is to be returned. If pid is not specified, basic information is retrieved for all processes you have authorization to view.

request
Specify one of these:

process
Retrieve information about the specified process. This is the default.

file
Retrieve file information for the specified process.

thread
Retrieve information for all threads in the specified process.

Results

Information is returned in simple and compound variables. Each variable name has a prefix that defaults to bpxw_. This prefix can be changed using the rexxopt() function. A stem can be used as the prefix, in which case the tails are the simple suffixes set by this function.

If pid is not specified, the following suffixes are set:

<table>
<thead>
<tr>
<th>Suffix</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>PID.n</td>
<td>Each variable contains one process ID number. PID.0 is the number of PIDs returned. PID.1, PID.2, ... are the PID numbers.</td>
</tr>
<tr>
<td>THREADS.n</td>
<td>The number of threads for the process in the corresponding PID.n</td>
</tr>
<tr>
<td>ASID.n</td>
<td>The address space ID for the process in the corresponding PID.n</td>
</tr>
<tr>
<td>JOBNAME.n</td>
<td>The jobname for the process in the corresponding PID.n</td>
</tr>
<tr>
<td>LOGNAME.n</td>
<td>The login name (user ID) for the process in the corresponding PID.n</td>
</tr>
</tbody>
</table>
If process information is requested for a process, the following suffixes are set:

<table>
<thead>
<tr>
<th>Suffix</th>
<th>Description</th>
</tr>
</thead>
</table>
| STATE  | Contains 0 or more of the following strings:  
• MULPROCESS  
• SWAP  
• TRACE  
• STOPPED  
• INCOMPLETE  
• ZOMBIE |
| PID    | Process ID number |
| PPID   | Parent process ID number |
| PGRID  | Process group ID number |
| SID    | Session ID number |
| FPGID  | Foreground process group number |
| EUID   | Effective user ID |
| RUID   | Real user ID |
| SUID   | Saved set user ID |
| EGID   | Effective group ID |
| RGID   | Real group ID |
| SGID   | Saved set group ID |
| SIZE   | Region size |
| SLOWPATH | Number of slow path syscalls |
| USERTIME | Time spent in user code |
| SYSTIME | Time spent in system code |
| STARTTIME | Time when the process started (dub time) |
| OETHREADS | Number of z/OS UNIX threads |
| PTCREATE | Number of threads created using pthread_create |
| THREADS | Number of threads in the process |
| ASID   | Address space ID of the process |
| JOBNAME | Jobname for the process |
| LOGNAME | Login name (user ID) for the process |
| CONTTY | Pathname for the controlling TTY |
| CMDPATH | Pathname for the command that started the process |
| CMDLINE | Command line that started the process |

If file information is requested for a process, the following suffixes are set:

<table>
<thead>
<tr>
<th>Suffix</th>
<th>Description</th>
</tr>
</thead>
</table>
| TYPECD | Type code of the file:  
• rd: root directory  
• cd: current directory  
• fd: file descriptor  
• vd: vnode descriptor |
| TYPE   | Type number of the file |
| OPENF  | Open flags |
### procinfo()

<table>
<thead>
<tr>
<th>Suffix</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>INODE</td>
<td>Inode number</td>
</tr>
<tr>
<td>DEVNO</td>
<td>Device number</td>
</tr>
<tr>
<td>NODES</td>
<td>Count of the total number of file nodes returned</td>
</tr>
</tbody>
</table>

If thread information is requested for a process, the following suffixes are set:

<table>
<thead>
<tr>
<th>Suffix</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>THREAD_ID</td>
<td>Thread ID</td>
</tr>
<tr>
<td>SYSCALL</td>
<td>Current syscall, if in kernel</td>
</tr>
<tr>
<td>TCB</td>
<td>TCB address</td>
</tr>
<tr>
<td>RUNTIME</td>
<td>Time running in milliseconds</td>
</tr>
<tr>
<td>WAITTIME</td>
<td>Time waiting in milliseconds</td>
</tr>
<tr>
<td>SEMNUM</td>
<td>Semaphore number if on a semop (ptrunwait.x='D')</td>
</tr>
<tr>
<td>SEMVALUE</td>
<td>Semaphore value if on a semop (ptrunwait.x='D')</td>
</tr>
<tr>
<td>LATCHPID</td>
<td>Latch that the process waited for</td>
</tr>
<tr>
<td>SIGPENDMASK</td>
<td>Signal pending mask</td>
</tr>
<tr>
<td>LOGNAME</td>
<td>Login name</td>
</tr>
<tr>
<td>LASTSYSCALL.n</td>
<td>Last 5 syscalls</td>
</tr>
<tr>
<td>PTCREATED</td>
<td>Contains J if this was pthread-created</td>
</tr>
<tr>
<td>PTRUNWAIT</td>
<td>Contains one of the following letters, to indicate the current run or wait state of the thread:</td>
</tr>
<tr>
<td>A</td>
<td>msgrecv wait</td>
</tr>
<tr>
<td>B</td>
<td>msgsend wait</td>
</tr>
<tr>
<td>C</td>
<td>communication wait</td>
</tr>
<tr>
<td>D</td>
<td>Semaphore wait</td>
</tr>
<tr>
<td>F</td>
<td>File System Wait</td>
</tr>
<tr>
<td>G</td>
<td>MVS in Pause</td>
</tr>
<tr>
<td>K</td>
<td>Other kernel wait</td>
</tr>
<tr>
<td>P</td>
<td>PTwaiting</td>
</tr>
<tr>
<td>R</td>
<td>Running or non-kernel wait</td>
</tr>
<tr>
<td>S</td>
<td>Sleep</td>
</tr>
<tr>
<td>W</td>
<td>Waiting for child</td>
</tr>
<tr>
<td>X</td>
<td>Fork new process</td>
</tr>
<tr>
<td>Y</td>
<td>MVS wait</td>
</tr>
</tbody>
</table>
### Usage notes

1. This function uses the `__getthent callable service`. For additional information, see the [__getthent (BPX1GTH, BPX4GTH) — Get thread data in z/OS UNIX System Services Programming: Assembler Callable Services Reference](#).

2. The `procinfo()` function returns as its value the prefix used to create variable names. If there is an error returned by the `__getthent` service, the function returns a null string, and sets the variables `ERRNO` and `ERRNOJR` to the hex values of the error codes returned by the `__getthent` service.

### Example

To show the command line for each process:
```
call procinfo
   do i=1 to bpxw_pid.0
      if procinfo(bpxw_pid.i)<>'' then /* ignore processes that ended */
         say bpxw_cmdline
   end
```

To show the current directory for each process:
```
call procinfo
   do i=1 to bpxw_pid.0
      if procinfo(bpxw_pid.i,'file')<>'' then /* ignore processes that ended */
         do j=1 to bpxw_nodes
            if bpxw_typecd.j='cd' then
               do
                  address syscall 'getmntent m.' bpxw_devno.j
                  strm=stream('find m.mnte_path.1 -xdev -inum', bpxw_inode.j, 'c', 'popen read')
                  say right(bpxw_pid.i,12) linein(strm)
                  call stream strm,'c','pclose'
               leave
            end
         end
   end
```

---

**procinfo()**

<table>
<thead>
<tr>
<th>Suffix</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>PTTYPE</td>
<td>Contains one of the following letters to indicate thread type:</td>
</tr>
<tr>
<td>N</td>
<td>Medium-wait thread</td>
</tr>
<tr>
<td>O</td>
<td>Asynchronous thread</td>
</tr>
<tr>
<td>U</td>
<td>Initial process thread</td>
</tr>
<tr>
<td>Z</td>
<td>Process terminated and parent has not completed wait</td>
</tr>
<tr>
<td>PTDETACH</td>
<td>Contains V if the thread is detached</td>
</tr>
<tr>
<td>PTTRACE</td>
<td>Contains E if the thread is in quiesce freeze</td>
</tr>
<tr>
<td>PTTAG</td>
<td>Contains the ptagdata string if it exists</td>
</tr>
<tr>
<td>THREADS</td>
<td>Contains the total number of threads</td>
</tr>
</tbody>
</table>
Function

Sets, resets, or queries z/OS UNIX REXX options.

Parameters

option

Specify one of these:

immed

Associates the immediate command interrupt handler to a signal number.

noimmed

Restores the default action for a signal. arg1 is the signal number.

version

Returns a string showing last compile time for the function package.

varpref

Sets the variable name prefix used by some of the REXX functions. arg1 is the new prefix. If arg1 is not specified, the current value is not changed. If it is specified but null, it is defaulted to bpxw_. This returns a string that is the current setting.

notsoin

Disables any input to the TSO command processor that was started using ADDRESS TSO, including the stack. This returns the string TSO input disabled.

tsoin

Sets back to its default setting the input mode to the TSO command processor that was started using ADDRESS TSO. This returns the string TSO input enabled.

arg1

A signal number.

Example

To make the interrupt signal prompt for a command:

call rexxopt 'immed',sigint
Function

Places the process in a signals-enabled wait, and returns after the wait expires. If a signal interrupts the wait, the function returns the number of seconds remaining for the wait, otherwise it returns 0.

Parameters

seconds

A positive whole number for the number of seconds to sleep

Example

To wait 5 seconds:
call sleep 5
stream

Function

Returns the state of the stream or the result of the command.

Parameters

name
The stream name for the operation. This argument is case-sensitive.

operation
The operation that is to be performed. Only the first character is significant. The operations can be:

S  Returns the current state of the stream. This is the default operation. The following values can be returned:
  ready  The stream is available for normal input or output operations.
  error  The stream encountered an error on an input or output operation. After this call, the error condition is reset, and normal input or output operations can be attempted once again.
  unknown  The specified stream is not open.

D  This operation is almost the same as S, with one exception. When ERROR is returned, it is followed by text describing the error. This text usually contains the error number and reason code if the failure is due to a system call failure.

C  Execute a command for the specified stream. The command is specified as the third argument.

command
The stream command that is to be executed. This argument is valid only when the operation is C. The following commands are supported:

clearfile  Truncates a file to zero bytes for a stream that is opened for write. clearfile should only be used on a persistent stream.
  This example empties the file:
  call stream name,'c','clearfile'

close  Closes a stream. On success, the function returns the string ready. If the stream is not known, it returns the string UNKNOWN.
  This example closes the stream:
  call stream file,'c','close'

infileno  Returns the file descriptor number of the input side of a stream.
stream

The file descriptor can be used on lower-level system calls, using, for example, ADDRESS SYSCALL. This example gets the file descriptor for the read side of the stream:

```plaintext
fd=stream(name,'c','infileno')
```

nosignal

Disables Halt Interruption for stream errors. See nosignal. This example disables the halt signal for the standard input/output stream:

```plaintext
call stream , 'c' , 'nosignal'
```

open <open-type>

Opens a stream. open-type specifies either read or write. If open-type is not specified, read is assumed. The function returns a string that is the name to be used for the stream on subsequent I/O functions. This string is the only name by which this stream will be known. Note that streams do not have to be explicitly opened. A pathname that is used as a stream name on the other I/O functions will cause the stream to be opened, and the name of the stream will be that pathname. Optional arguments may be specified with write, optionally followed by octal permission bits. Write always opens the file with O_CREAT, creating the file if it does not exist. The additional arguments are:

replace

Opens the file with O_TRUNC, setting the file size to 0.

append

Opens the file with O_APPEND, causing all writes to be appended to the file.

This example opens a stream for the file mydata.txt:

```plaintext
file=stream('mydata.txt','c','open write')
```

This example opens a stream for the file mydata.txt, but replaces the file if it exists:

```plaintext
file=stream('mydata.txt','c','open write replace')
```

outfileno

Returns the file descriptor number of the output side of a stream. The file descriptor can be used on lower-level system calls, using, for example, ADDRESS SYSCALL. This example gets the file descriptor for the write side of the stream:

```plaintext
fd=stream(name,'c','outfileno')
```

pclose

Closes a process stream. On success, the function returns the completion code for the process run via the open command.

pid

Returns the process ID number for the shell process opened with open. This example gets the PID for a stream opened with open:

```plaintext
pid=stream(name,'c','pid')
```

popen <open-type>

Opens a pipe to a shell command that is specified by the stream name. open-type must specify read or write. If read is specified, the input functions can be used to read the standard output from the command. If write is specified, the output functions can be used to pipe data to the standard input of the shell command. In either case, the shell command inherits the standard error file for the calling process.
stream

The command that is run is always /bin/sh -c followed by the specified shell command. This means that the completion code is the one returned by the shell. It usually returns the command's completion code. You can obtain the completion code by using the `pclose` command.

The function returns a string that is the name to be used for the stream on subsequent I/O functions. This string is the only name by which this stream will be known.

This example opens a pipe stream to the output from the shell command `ls | wc`:

```plaintext
file=stream('ls | wc','c','popen read')
```

**query <attribute>**

Queries a stream attribute and returns the result:

- **exists**
  - Returns the full pathname of the stream name. This is equivalent to the `exists()` function, but is more portable. This command does not cause a stream to be opened.

- **size**
  - Returns the size of the stream. This is equivalent to the `size stream` command, but is more portable.

This example prints the pathname of the stream:

```plaintext
say stream('myfile','c','query exists')
```

**readpos [location]**

Returns the position in the file where the next read will begin. If `location` is specified, the position is also set to the byte specified by location. `location` is specified as a number optionally preceded by one of the following characters:

- `=` An absolute byte location. This is the default.
- `<` An offset from the end of the stream.
- `+` An offset forward from the current location.
- `-` An offset backward from the current location.

This example gets the read location in the file, and then sets the read location to the sixth 80-byte record:

```plaintext
pos=stream(name,'c','readpos') /* get read location */
call stream name,'c','readpos' 5*80+1 /* set read location*/
```

**signal**

Enables Halt Interruption for stream errors. Note that the NOTREADY REXX signal is not supported. This example enables the halt signal for the standard input/output stream:

```plaintext
call stream , 'c','signal'
```

**size**

Returns the size of the file associated with the stream. This example prints the size of the file:

```plaintext
say stream(name,'c','size')
```

**writepos <location>**

Returns the position in the file where the next write will begin. If `location` is specified, the position is also set to the byte specified by location. `location` is specified as a number optionally preceded by one of the following characters:
An absolute byte location. This is the default.

< An offset from the end of the stream.

+ An offset forward from the current location.

− An offset backward from the current location.

This example sets the position to the end of the file:
call stream name,'c','writepos <0'

This example sets the position to the start of the file:
call stream name,'c','writepos 1'
**submit()**

```c
submit(stem.)
```

**Function**

Submits a job to the primary subsystem (JES), returning the job ID of the submitted job.

**Parameters**

```c
stem.
```

The stem compound variable contains the number of lines in `stem.0`, and each variable from `stem.1`, `stem.2`, ... contains a line for the job that is being submitted.

**Example**

This example reads the file into the stem, sets the number of lines, and submits the job:

```plaintext
do i=1 by 1 while lines(fn)>0
   fn.i=linein(fn)
end
fn.0=i-1
say submit('fn.')
```
Function

Establishes the SYSCALL environment or ends it; or establishes or deletes the signal interface routine (SIR).

Parameters

ON
Establishes the SYSCALL environment. It sets up the REXX predefined variables and blocks all signals.

OFF
Ends the connection between the current task and z/OS UNIX. It is usually not necessary to make a `syscalls (OFF)` call.

SIGON
Establishes the signal interface routine.

SIGOFF
Deletes the signal interface routine, and resets the signal process mask to block all signals that can be blocked.

Usage notes

1. This function can be used outside of the z/OS UNIX REXX environment (for example, in TSO/E). When `syscalls` is used in such an environment, `stdin`, `stdout`, `stderr`, and environment variables are not inherited from the current process environment.
2. For more usage information, see "The SYSCALL environment" on page 2.

Example

See "The SYSCALL environment" on page 2.
Chapter 6. BPXWDYN: a text interface to dynamic allocation and dynamic output

BPXWDYN is a text interface to a subset of the SVC 99 (dynamic allocation) and SVC 109 (dynamic output) services. BPXWDYN supports data set allocation, unallocation, concatenation, the retrieval of certain allocation information, and the addition and deletion of output descriptors. BPXWDYN is designed to be called from REXX, but it may be called from several other programming languages, including Assembler, C, and PL/I.

This interface makes dynamic allocation and dynamic output services easily accessible to programs running outside of a TSO environment; however, it also functions in a TSO environment.

The syntax for allocation is quite similar to that of TSO for the TSO ALLOCATE and FREE commands. It should be possible to provide parameters to BPXWDYN that would be acceptable as a TSO ALLOCATE or FREE command. However, there are keys supported by TSO ALLOCATE that are not currently supported by BPXWDYN. There are also some keys that can be used with BPXWDYN which are not compatible with TSO.

The syntax for accessing dynamic output facilities is similar to that of the TSO OUTDES command, but the name of the output descriptor is identified differently. You associate an output descriptor with a SYSOUT allocation by using the OUTDES key on the SYSOUT allocation request, or by creating a default output descriptor.

Calling conventions

BPXWDYN must be called in 31-bit mode in an environment that permits dynamic allocation and dynamic output requests. To call BPXWDYN from REXX or any other program, three parameter list forms can be used:

- REXX external function parameter list
- Conventional MVS variable-length parameter string
- Null-terminated parameter string

REXX external function parameter list

The external function parameter list allows a REXX program to call the BPXWDYN program as a function or subroutine. The BPXWDYN program must be called with a single string parameter. For example:

```plaintext
if BPXWDYN("alloc dd(sysin) da(my.dataset) shr")<>0 then
call allocfailed
```

Conventional MVS variable-length parameter string

This is the same parameter list as the one generated by ADDRESS LINKMVS with one parameter, and by JCL with EXEC PGM=,PARM=. Any program can easily use this parameter list form. Note that this is a single-item variable-length parameter list. The high bit is on in the parameter address word, and length is a halfword.
PL/I usage might include the following statements:

```pli
DCL PLIRETV BUILTIN;
DCL BPXWDYN EXTERNAL ENTRY OPTIONS(ASM INTER RETCODE);
DCL ALLOC_STR CHAR(100) VAR
  INIT('ALLOC FI(SYSIN) DA(MY.DATASET) SHR');
FETCH BPXWDYN;
CALL BPXWDYN(ALLOC_STR);
```

### Null-terminated parameter string

This parameter list is used most easily from C, passing a string to BPXWDYN defined with a `#pragma` for OS linkage, or extern "OS" for C++. Note that the high bit is on in the parameter address word.

```c
R1 → 1 parm string addr
    
    parameter string 00x
```

C usage might include the following statements:

```c
typedef int EXTF();
#pragma linkage(EXTF,OS)
EXTF *bpxwdyn;
int rc;
bpxwdyn=(EXTF *)fetch("BPXWDYN");
rc=bpxwdyn("alloc fi(sysin) da(my.dataset) shr");
```

### Request types

The request is specified in a string parameter. The request type should be the first keyword in the parameter string. If a valid request type is not the first keyword, or no request type is specified, the request type defaults to ALLOC.

The following request types are supported:

- **ALLOC**  Dynamic allocation of a data set
- **CONCAT** Dynamic concatenation of a ddname list
- **FREE**   Dynamic unallocation of a ddname, or freeing of an output descriptor
- **INFO**   Retrieval of allocation information
OUTDES  Creation of an output descriptor

Keywords

The ordering of keywords on the request is arbitrary. You cannot specify the same keys multiple times for one request.

BPXWDYN does no consistency checking on keys. If some keys are not valid when combined together on a single request, dynamic allocation or dynamic output fails the request.

Some keywords accept arguments:

- Keyword arguments must be specified within parentheses. Spaces are not permitted between the key and the opening parenthesis, or anywhere within the additional argument string through the closing parenthesis, unless the argument string is quoted. Spaces are permitted between key specifications. Where multiple arguments are permitted, the arguments must be separated by commas.
- An argument may be enclosed within single (') or double (") quotes. Two adjacent quotes cannot be used to represent a single quote within a quoted string.
- Arguments that are not quoted are treated as uppercase. Arguments that are quoted are only treated as mixed case if it makes sense for the keyword.
- Arguments that accept REXX variable names will be accepted when BPXWDYN is called from any environment, but are only effective when called from a non-reentrant REXX environment.

Return codes

When BPXWDYN is called as a REXX function or subroutine, the return code can be accessed in RESULT or as the value of the function. When BPXWDYN is called as a program, the return code is available in R15.

BPXWDYN returns the following codes:

0   Success.
20  Invalid parameter list. See "Calling conventions" on page 245 for parameter list formats.
-21 to -9999 Key error. See "Key errors" for more information.
-100nn Message processing error. IEFDB476 returned code nn.
>0, or -1610612737 to -2147483648 Dynamic allocation or dynamic output error codes. See "Error codes for dynamic allocation" on page 248 and "Error codes for dynamic output" on page 248 for more information.

Key errors

The low-order two digits are the number of the key that failed the parse, offset by 20. The first key has the value 21, resulting in a return code of -21.

The high-order two digits are an internal diagnostic code that indicates where in the code the parse failed. If the high digits are 0, the key itself was probably not recognized. Other values usually indicate a problem with the argument for the failing keyword. Likely causes for failure are:

- Blanks in arguments that are not quoted
Error codes for dynamic allocation

The return code contains the dynamic allocation error reason code in the high two bytes, and the information reason code in the low two bytes.

You can use the high four hexadecimal digits to look up the error code in the dynamic allocation error reason codes table found in the "z/OS MVS Programming: Authorized Assembler Services Guide".

Note: For the returned decimal codes (>0, or -1610612737 to -2147483648), convert them to hexadecimal codes before determining the value for the high and low two bytes.

Error codes for dynamic output

Dynamic output errors do not produce messages. The return code contains:

- The dynamic output return code (S99ERROR) in the high two bytes
- The information code (S99INFO) in the low two bytes

You can use the high four hexadecimal digits to look up the error code in the dynamic output return codes table found in the "z/OS MVS Programming: Authorized Assembler Services Guide". You can use the low four digits to look up the information code.

There is no indication of the key that is in error.

Some of the most common reason codes are:

- **300-30C, 312, 380**
  - The arguments may be specified incorrectly.
- **401**
  - The output descriptor already exists.
- **402**
  - The output descriptor does not exist.
- **403**
  - Output descriptors created by JCL cannot be deleted by dynamic output.

Message processing

Dynamic allocation provides message text for failed allocation requests. It is usually easier to use this message text rather than decode allocation return codes.

BPXWDYN can return these messages to a REXX program. By default, the messages are returned in the S99MSG.0 stem. S99MSG.0 contains the number of messages and S99MSG.1 through S99MSG.\(n\) contain the message text.

To change the name of the stem, use the MSG key. A stem is not required. Digits are simply appended to the variable specified by the MSG key.

You can also use the MSG key to request allocation to issue the messages to your job log (write to programmer using WTO).

Under z/OS UNIX, you can have messages written to opened files by providing the file descriptor number as the argument to the MSG key. For example, MSG(2) writes messages to the STDERR file.
Requesting dynamic allocation

To request dynamic allocation, specify `alloc` at the beginning of the parameter string.

BPXWDYN supports single data set allocation. Many of the common allocation keys can be used. For detailed information about those common allocation keys, see [z/OS TSO/E Command Reference](#). Some additional keys are supported to access additional functions.

BPXWDYN supports the following keys:

Table 3. Common keys used for dynamic allocation

<table>
<thead>
<tr>
<th>Key</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>DA(data set name[(member name)])</td>
<td>Specifies the data set name to allocate. The name must be fully qualified and may include a member name. Quotes can be used for TSO compatibility.</td>
</tr>
<tr>
<td>DUMMY</td>
<td>Allocates a dummy data set.</td>
</tr>
<tr>
<td>FI(name)</td>
<td>Specifies the ddname to allocate.</td>
</tr>
<tr>
<td>FILEDATA(TEXT</td>
<td>TELLS the sequential access method services whether the data is to be treated as text or binary.</td>
</tr>
<tr>
<td>OLD</td>
<td>Specifies the data set status. SYSOUT specifies that the data set is to be a system output data set and optionally defines the output class.</td>
</tr>
<tr>
<td>VOL(volser</td>
<td>Specifies the serial numbers for eligible direct access volumes where the data set is to reside.</td>
</tr>
<tr>
<td>DATACLAS(data class)</td>
<td>With SMS, specifies the data class for the data set.</td>
</tr>
<tr>
<td>MGMTCLAS(management class)</td>
<td>With SMS, specifies the management class for the data set.</td>
</tr>
<tr>
<td>STORCLAS(storage class)</td>
<td>With SMS, specifies the storage class for the data set.</td>
</tr>
<tr>
<td>SPACE( primary[, secondary])</td>
<td>Specifies primary and optionally secondary space allocations.</td>
</tr>
<tr>
<td>BLOCK</td>
<td>Specifies the unit of space in blocks.</td>
</tr>
<tr>
<td>TRACKS</td>
<td>Specifies the unit of space in tracks.</td>
</tr>
<tr>
<td>CYL</td>
<td>Specifies the unit of space in cylinders.</td>
</tr>
<tr>
<td>BLKSIZE(block size)</td>
<td>Specifies the block size.</td>
</tr>
<tr>
<td>DIR(directory blocks)</td>
<td>Specifies the number of directory blocks.</td>
</tr>
<tr>
<td>DEST([destination</td>
<td>destination[.user]])</td>
</tr>
<tr>
<td>REUSE</td>
<td>Causes the named DD to be freed before the function is performed.</td>
</tr>
<tr>
<td>HOLD</td>
<td>Specifies that the output data set is to be held until released by the user or operator.</td>
</tr>
<tr>
<td>UNIT(unit name)</td>
<td>Specifies unit name, device type, or unit address.</td>
</tr>
</tbody>
</table>
Table 3. Common keys used for dynamic allocation  (continued)

<table>
<thead>
<tr>
<th>Key</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>MAXVOL(num vols)</td>
<td>Number of volumes for a multi-volume data set.</td>
</tr>
<tr>
<td>KEEP</td>
<td>DELETE</td>
</tr>
<tr>
<td>BUFNO(number)</td>
<td>Sets the number of buffers. This number should be in the range 1–255. Numbers outside that range give unpredictable results.</td>
</tr>
<tr>
<td>LRECL(record length)</td>
<td>Specifies the logical record length.</td>
</tr>
<tr>
<td>RECFM(format[,format...])</td>
<td>Specifies the record format. The valid values are A, B, D, F, M, S, T, U, and V. Several of these can be used in combination.</td>
</tr>
<tr>
<td>DSORG(PS</td>
<td>PO</td>
</tr>
<tr>
<td>COPIES(number of copies)</td>
<td>Specifies the number of copies to print.</td>
</tr>
<tr>
<td>FORMS(forms name)</td>
<td>Specifies the print form.</td>
</tr>
<tr>
<td>LIKE(model data set name)</td>
<td>Copies attributes for the allocation from the model data set.</td>
</tr>
<tr>
<td>OUTDES(output descriptor name)</td>
<td>Specifies the output descriptor name.</td>
</tr>
<tr>
<td>SPIN(UNALLOC)</td>
<td>Spins off a sysout data set at unallocation.</td>
</tr>
<tr>
<td>DSNTYPE(LIBRARY</td>
<td>PDS</td>
</tr>
<tr>
<td>WRITER(external writer name)</td>
<td>Names the external writer.</td>
</tr>
<tr>
<td>PATH(pathname)</td>
<td>Specifies the path name of an HFS file to allocate.</td>
</tr>
<tr>
<td>PATHDISP(KEEP</td>
<td>DELETE[,KEEP</td>
</tr>
<tr>
<td>PATHMODE(path mode list)</td>
<td>Set mode bits for a new allocation. This list is a list of keywords separated with commas. The supported keywords are:</td>
</tr>
<tr>
<td></td>
<td>SIRUSR</td>
</tr>
<tr>
<td></td>
<td>SIWUSR</td>
</tr>
<tr>
<td></td>
<td>SIXUSR</td>
</tr>
<tr>
<td></td>
<td>SIRWXU</td>
</tr>
<tr>
<td></td>
<td>SIRGRP</td>
</tr>
<tr>
<td></td>
<td>SIWGRP</td>
</tr>
<tr>
<td></td>
<td>SIXGRP</td>
</tr>
<tr>
<td></td>
<td>SIRWXG</td>
</tr>
<tr>
<td></td>
<td>SIROTH</td>
</tr>
<tr>
<td></td>
<td>SIWOTH</td>
</tr>
<tr>
<td></td>
<td>SIXOTH</td>
</tr>
<tr>
<td></td>
<td>SIRWXO</td>
</tr>
<tr>
<td></td>
<td>SISUID</td>
</tr>
<tr>
<td></td>
<td>SISGID</td>
</tr>
<tr>
<td></td>
<td>SISVTX</td>
</tr>
</tbody>
</table>
### Table 3. Common keys used for dynamic allocation (continued)

<table>
<thead>
<tr>
<th>Key</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td><strong>PATHOPTS</strong> <em>(path options list)</em></td>
<td>Sets options for path name allocation. The options list is a list of keywords separated with commas. The supported keywords are: ORDWR, OEXCL, OSYNC, OTRUNC, OCREAT, OWRONLY, ORDONLY, OAPPEND, ONOCTTY, ONONBLOCK</td>
</tr>
<tr>
<td><strong>RECORGLS</strong></td>
<td>Creates a VSAM linear data set.</td>
</tr>
<tr>
<td><strong>SEQUENCE</strong> <em>(sequence number)</em></td>
<td>Specifies the relative position (number) of a data set on a tape volume.</td>
</tr>
<tr>
<td><strong>LABEL</strong> <em>(type)</em></td>
<td>Specifies the type of tape label processing to be done, as follows:</td>
</tr>
<tr>
<td>NL</td>
<td>The volume has no label.</td>
</tr>
<tr>
<td>SL</td>
<td>The volume has an IBM standard label.</td>
</tr>
<tr>
<td>NSL</td>
<td>The volume has a non-standard label.</td>
</tr>
<tr>
<td>SUL</td>
<td>The volume has both an IBM standard label and a user label.</td>
</tr>
<tr>
<td>BLP</td>
<td>Bypass label processing for the volume.</td>
</tr>
<tr>
<td>LTM</td>
<td>The system is to check for and bypass a leading tape mark on a DOS unlabeled tape.</td>
</tr>
<tr>
<td>AL</td>
<td>The volume has an American National Standard label.</td>
</tr>
<tr>
<td>AUL</td>
<td>The volume has both an American National Standard label and a user label.</td>
</tr>
<tr>
<td><strong>RETPD</strong> <em>(number of days)</em></td>
<td>Specifies the data set retention period, in days.</td>
</tr>
<tr>
<td><strong>TRTCH</strong> <em>(technique)</em></td>
<td>Specifies the tape recording technique, as follows:</td>
</tr>
<tr>
<td>NONCOMP</td>
<td>Non-compaction mode</td>
</tr>
<tr>
<td>COMP</td>
<td>Compaction mode</td>
</tr>
<tr>
<td>C</td>
<td>Data conversion</td>
</tr>
<tr>
<td>E</td>
<td>Even parity</td>
</tr>
<tr>
<td>ET</td>
<td>Even parity and BCD/EBCDIC translation</td>
</tr>
<tr>
<td>T</td>
<td>BCD/EBCDIC translation</td>
</tr>
<tr>
<td><strong>SUBSYS</strong> <em>(subsystem name[, subsys parm...])</em></td>
<td>Specifies that the data set is a subsystem data set, and also specifies the name of the subsystem and any parameters necessary for the subsystem to process the data set. Case and special characters are preserved for quoted parameters.</td>
</tr>
</tbody>
</table>
The following additional keys are unique to BPXWDYN.

### Table 4. Additional keys used for dynamic allocation

<table>
<thead>
<tr>
<th>Key</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>MSG (WTP</td>
<td>default.S99MSG.</td>
</tr>
<tr>
<td>NORECALL</td>
<td>Fail the allocation request if the data set is migrated.</td>
</tr>
<tr>
<td>PATHPERM(</td>
<td>octal path mode)</td>
</tr>
<tr>
<td>RTDDN(variable)</td>
<td>Return allocated ddname into the REXX variable variable.</td>
</tr>
<tr>
<td>RTDSN(variable)</td>
<td>Return allocated data set name into the REXX variable variable.</td>
</tr>
<tr>
<td>RTVOL(variable)</td>
<td>Return allocated volume name into the REXX variable variable.</td>
</tr>
<tr>
<td>SHORTRC</td>
<td>If the dynamic allocation fails, the dynamic allocation error code (S99ERROR) is returned in R15 and the information code is not returned. This is useful if the application can only process R15 as a halfword, such as with PL/I.</td>
</tr>
<tr>
<td>SYNTAX</td>
<td>When this is used, BPXWDYN parses the request and does not issue the dynamic allocation or dynamic output request.</td>
</tr>
<tr>
<td>RELEASE</td>
<td>Space allocated to an output data set but not used is released when the data set is closed.</td>
</tr>
<tr>
<td>UCOUNT(number)</td>
<td>Specifies the number of devices to be allocated.</td>
</tr>
<tr>
<td>GDGNT</td>
<td>The S99GDGNT flag is set in the S99FLAG1 field. For information about this flag, see <a href="https://www.ibm.com/support/pages/zos-mvs-programming-authorized-assembler-services-guide">z/OS MVS Programming: Authorized Assembler Services Guide</a>.</td>
</tr>
<tr>
<td>MOUNT</td>
<td>Resets the S99NOMNT flag, allowing volumes to be mounted.</td>
</tr>
</tbody>
</table>

### Requesting dynamic concatenation

To request dynamic concatenation, specify **concat** at the beginning of the parameter string.

This function concatenates multiple DDs to a single DD. The DDs are concatenated in the order specified in the DDLIST key.
Table 5. Keys used for dynamic concatenation

<table>
<thead>
<tr>
<th>Key</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>DDLIST([DDname1,DDname2,[DDnamex...]])</td>
<td>Specifies a list of ddnames to concatenate to DDname1. Use this key with the CONCAT key.</td>
</tr>
<tr>
<td>MSG (WTP I default.S99MSG.\ stemname I fdnum)</td>
<td>Directs allocation messages to your job log (WTP), a REXX stem, or a file identified by a file descriptor number. If this key is not specified, messages will be returned in the S99MSG. stem, if possible. If BPXWDYN was not called from a REXX environment, the messages will be lost.</td>
</tr>
<tr>
<td>SHORTRC</td>
<td>If the dynamic allocation fails, the dynamic allocation error code (S99ERROR) is returned in R15 and the information code is not returned. This is useful if the application can only process R15 as a halfword, such as with PL/I.</td>
</tr>
</tbody>
</table>

Requesting dynamic unallocation

To request dynamic unallocation, specify free at the beginning of the parameter string.

BPXWDYN supports single data set allocation. Many of the common unallocation keys can be used. For detailed information about those common keys, see [z/OS TSO/E Command Reference](#). Some additional keys are supported to access additional functions.

Table 6. Common keys used for dynamic unallocation

<table>
<thead>
<tr>
<th>Key</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>DA(data set name[[(member name)]])) I DSN(data set name[[(member name)]]))</td>
<td>Specifies the data set name to free. The name must be fully qualified and may include a member name. Quotes are optional.</td>
</tr>
<tr>
<td>FL(name) I DD(name)</td>
<td>Identifies the ddname to free.</td>
</tr>
<tr>
<td>KEEP</td>
<td>DELETE</td>
</tr>
<tr>
<td>SPIN(UNALLOC)</td>
<td>Spins off a sysout data set at unallocation.</td>
</tr>
<tr>
<td>SYSOUT(class)</td>
<td>Overrides the output class.</td>
</tr>
</tbody>
</table>

The following additional keys are unique to BPXWDYN.

Table 7. Additional keys used for dynamic unallocation

<table>
<thead>
<tr>
<th>Key</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>MSG (WTP I default.S99MSG.\ stemname I fdnum)</td>
<td>Directs allocation messages to your job log (WTP), a REXX stem, or a file identified by a file descriptor number. If this key is not specified, messages will be returned in the S99MSG. stem, if possible. If BPXWDYN was not called from a REXX environment, the messages will be lost. WTP should be specified to obtain messages.</td>
</tr>
</tbody>
</table>
Table 7. Additional keys used for dynamic unallocation (continued)

<table>
<thead>
<tr>
<th>Key</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>SHORTRC</td>
<td>If the dynamic allocation fails, the dynamic allocation error code (S99ERROR) is returned in R15 and the information code is not returned. This is useful if the application can only process R15 as a halfword, such as with PL/I.</td>
</tr>
</tbody>
</table>

Requesting allocation information

To request information about the current dynamic allocation environment, specify `info` at the beginning of the parameter string.

BPXWDYN currently supports the retrieval of the ddnames, data set names, and path names for current allocations; other allocation attributes are not supported.

BPXWDYN supports single data set allocation. You can use specific keys to specify which allocation requests information is to be returned for.

Table 8. Keys used to specify which allocation request information is to be returned for

<table>
<thead>
<tr>
<th>Key</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>DA(data set name[(member name)])</td>
<td>Specifies the data set name of the allocated resource about which you are requesting information. The data set name can contain special characters if the data set name is enclosed in apostrophes. It can also contain system symbols (see the section on using system symbols in z/OS MVS JCL Reference for more information). The maximum length of the data set name is 44 characters, excluding any enclosing apostrophes and compressing any double apostrophes within the data set name. This key is mutually exclusive with the ddname (FI, DD), path name (PATH), and relative entry (INRELNO) keys.</td>
</tr>
<tr>
<td>DSN(data set name[(member name)])</td>
<td>Specifies the ddname of the allocated resource about which you are requesting information. name contains the ddname. This key is mutually exclusive with the data set name (DA, DSN), path name (PATH), and relative entry (INRELNO) keys.</td>
</tr>
<tr>
<td>FI(name)</td>
<td>Specifies the ddname of the allocated resource about which you are requesting information. name contains the ddname. This key is mutually exclusive with the data set name (DA, DSN), path name (PATH), and relative entry (INRELNO) keys.</td>
</tr>
<tr>
<td>DD(name)</td>
<td>Specifies the ddname of the allocated resource about which you are requesting information. name contains the ddname. This key is mutually exclusive with the data set name (DA, DSN), path name (PATH), and relative entry (INRELNO) keys.</td>
</tr>
<tr>
<td>PATH(pathname)</td>
<td>Specifies the path name of the HFS file for which you are requesting information. pathname contains the path name and can contain system symbols (see the section on using system symbols in z/OS MVS JCL Reference for more information). This key is mutually exclusive with the data set name (DA, DSN), ddname (FI, DD), and relative entry (INRELNO) keys.</td>
</tr>
</tbody>
</table>
## Requesting dynamic output

To request dynamic output, specify `OUTDES` at the beginning of the parameter string. This keyword takes an argument that names the output descriptor. The keyword is described in the following table. These are the supported keys:

<table>
<thead>
<tr>
<th>OUTDES(descriptor name)</th>
<th>Names the output descriptor to be added. This must be the first key specified in the parameter string.</th>
</tr>
</thead>
<tbody>
<tr>
<td>Key</td>
<td>Description</td>
</tr>
<tr>
<td>---------------------</td>
<td>-----------------------------------------------------------------------------</td>
</tr>
<tr>
<td>ADDRESS(address[,address...])</td>
<td>Specifies the delivery address. Dynamic output allows up to four arguments to be specified.</td>
</tr>
<tr>
<td>BUILDING(building)</td>
<td>Specifies the building location.</td>
</tr>
<tr>
<td>BURST</td>
<td>Directs output to a stacker.</td>
</tr>
<tr>
<td>CHARS(chars[,chars...])</td>
<td>Names the character arrangement tables.</td>
</tr>
<tr>
<td>CLASS(class)</td>
<td>Assigns the SYSOUT class.</td>
</tr>
<tr>
<td>CONTROL(spacing)</td>
<td>Specifies either line spacing or that the records begin with carriage control characters. The valid values are: Single Single spacing Double Double spacing Triple Triple spacing Program Records begin with carriage control characters</td>
</tr>
<tr>
<td>COPIES(number of copies)</td>
<td>Specifies the number of copies to be printed.</td>
</tr>
<tr>
<td>DEFAULT</td>
<td>Specifies that this is a default output descriptor.</td>
</tr>
<tr>
<td>DEPT(department)</td>
<td>Specifies the department identification.</td>
</tr>
<tr>
<td>DEST(node[,user])</td>
<td>Sends the SYSOUT to the specified destination.</td>
</tr>
<tr>
<td>DPAGELBL</td>
<td>Specifies that a security label be placed on the output.</td>
</tr>
<tr>
<td>FCB(fcb name)</td>
<td>Specifies the FCB image.</td>
</tr>
<tr>
<td>FLASH</td>
<td>Specifies the forms overlay.</td>
</tr>
<tr>
<td>FORMDEF(formdef name)</td>
<td>Names the formdef.</td>
</tr>
<tr>
<td>FORMS(forms name)</td>
<td>Names the forms to print on.</td>
</tr>
<tr>
<td>MODIFY(trc number)</td>
<td>Specifies which character arrangement table is to be used.</td>
</tr>
<tr>
<td>NAME(owner name)</td>
<td>Specifies the owner's name.</td>
</tr>
<tr>
<td>NOTIFY([node.]user)</td>
<td>Sends print completion message to the destination.</td>
</tr>
<tr>
<td>OUTDISP(disposition)</td>
<td>Specifies the data set disposition.</td>
</tr>
<tr>
<td>PAGEDEF(pagedef name)</td>
<td>Names the pagedef.</td>
</tr>
<tr>
<td>PRMODE(print mode)</td>
<td>Identifies the process mode (LINE or PAGE).</td>
</tr>
<tr>
<td>ROOM(room identification)</td>
<td>Specifies the room identification.</td>
</tr>
<tr>
<td>TITLE(separator title)</td>
<td>Specifies the separator page title.</td>
</tr>
<tr>
<td>TRC</td>
<td>Specifies that the data set contains TRC codes.</td>
</tr>
<tr>
<td>UCS(UCS name)</td>
<td>Names the UCS or character arrangement table.</td>
</tr>
<tr>
<td>WRITER(external writer name)</td>
<td>Names an external writer to process the data set.</td>
</tr>
</tbody>
</table>
Table 10. Keys used for dynamic output (continued)

| USERDATA(userdata[,userdata...]) | Specifies installation-specific user data for a dynamic output statement. Dynamic output allows up to 16 1–60-character strings. Case is preserved if the string is quoted. |

Freeing an output descriptor

To request that an output descriptor be freed, specify free at the beginning of the parameter string. Only one keyword is supported:

Table 11. Key used to free an output descriptor

| OUTDES(descriptor name) | Names the output descriptor to be freed. |

Examples: Calling BPXWDYN from a REXX program

Allocate

This example allocates SYS1.MACLIB to SYSLIB and directs messages to z/OS UNIX standard error (stderr):

```rexx
call bpxwdyn "alloc fi(syslib) da(sys1.maclib) shr msg(2)"
```

Info

This example requests that the name of the data set allocated to ddname SYSLIB be returned in the REXX variable dsnvar.

```rexx
call bpxwdyn "info fi(syslib) inrtdsn(dsnvar)"
```

Free

This example frees SYSLIB and traps messages in stem S99 MSG:

```rexx
call bpxwdyn "free fi(syslib)"
```

Concatenate

This example concatenates SYS1.SBPXEXEC to SYSPROC:

```rexx
if bpxwdyn("alloc fi(tmp) da(sys1.sbpexec) shr msg(2)")=0 then
call bpxwdyn "concat ddlist(sysproc,tmp) msg(2)"
```

Create dynamic output descriptor

This example creates descriptor P20 with distribution information:

```rexx
call bpxwdyn "outdes(p20) dest(kgn.p20n10)",
"address('WJ Schoen', 'M/S 619')",
"name(wschoen) bin(0004) dept(64ba)"
```

Free descriptor

This example frees descriptor P20:

```rexx
call bpxwdyn "free outdes(p20)"
```

Example: calling BPXWDYN from C

This example allocates SYS1.MACLIB to SYSLIB and directs messages to z/OS UNIX standard error (stderr):
typedef int EXT();
#pragma linkage(EXTF,OS_UPSTACK)

int call_alloc()
{
    EXTF *bpxwdyn=(EXTF *)fetch("BPXWDYN");
    return bpxwdyn("alloc fi(syslib) da(sys1.maclib) shr msg(2)")
}
Chapter 7. Virtual file system (VFS) server syscall commands

A number of syscall commands are intended for file system server applications, such as a Network File System server. Although it is unlikely that you would implement a Network File System server using REXX, you can access the server callable services using z/OS UNIX REXX syscall commands. Because it is possible for a server to create a file in the file hierarchy with a name that cannot be accessed through conventional C functions—for example, a filename that has a slash (/) in it—these syscall commands may be useful if you need to obtain local access to such a file.

For detailed information about these services, see z/OS UNIX System Services File System Interface Reference.

Security

The file system server services are available only to a registered server. Only a superuser can use the callable service v_reg to register the process as a server.

The server services can bypass system security for file access.

For examples of REXX coding using these commands, see Chapter 8.

Tokens

Many tokens flow across the server interface. The types of tokens are:

**VFS** Represents a mounted file system.

**vnode** Represents a file or directory that is currently in use. This identifier is valid only until the token is released, using the v_rel callable service.

**FID** Uniquely identifies a file or directory in a particular mounted file system. The file or directory may or may not be currently in use. This identifier is valid across mounting and unmounting of the file system, as well as across z/OS UNIX re-IPLs.

You must specify tokens as variable names, not as strings. See "Specifying a syscall command" on page 19 for information on specifying variable names. When a token is returned to the exec, the value of the token is stored in the variable. When a token variable is used as a parameter on a syscall command, the token value is extracted from the variable. The format for a token is 8 bytes of binary data.
v_create

Function

v_create invokes the v_create callable service to create a new file in the directory represented by vntoken. The file can be a regular, FIFO, or character special file.

Parameters

vntoken
A variable name that contains the vnode token for the directory in which the file filename is to be created.

filename
The name of the file. It must not contain null characters.

type
A number used to specify the type of file to be created: a regular, FIFO, or character special file. You can specify one of the predefined variables beginning with S_to set the value—for example, S_ISREG. For a list of the variables, see "fstat" on page 83.

mode
A three-digit number, corresponding to the access permission bits. Each digit must be in the range 0–7, and all three digits must be specified. For more information on permissions, see Appendix B.

stem
The name of a stem variable used to return the status information. Upon return, stem.0 contains the number of variables that are returned. To obtain the desired status information, you can use a numeric value or the predefined variables beginning with ST_ used to derive the numeric value. See "stat" on page 176 for a list of the predefined variables or Appendix A for the numeric values.

vntoken2
A variable name that will contain the vntoken of the created file on return.

major
For a character special file (S_ISCHR), the device major number. For a complete description, see "mknod" on page 124.

minor
For a character special file (S_ISCHR), the device minor number. For a complete description, see "mknod" on page 124.

Usage notes

1. If the file named in filename already exists, the v_create service returns a failing return code, and no vnode token is returned.
2. The caller is responsible for freeing vnode tokens returned by the service by calling to the v_rel service when they are no longer needed.
Example

In the following example, assume that *filenm* and *vnod* were assigned values earlier in the exec:

```
v_create vnod (filenm) s_isreg 777 "st. filetok"
```
v_fstatfs

Function

v_fstatfs invokes the v_fstatfs callable service to return file system status for the file system containing the file or directory represented by the specified vntoken.

Parameters

vntoken
A variable name that contains the vnode token for a file or directory in the file system whose status is to be checked.

stem
The name of a stem variable used to return the status information. Upon return, stem.0 contains the number of variables that are returned. To obtain the desired status information, you can use a numeric value or the predefined variables beginning with STFS_ used to derive the numeric value. For example, stem.stfs_avail accesses the number of blocks available in the file system. See "statfs" on page 179 for a list of the predefined variables, or Appendix A for the numeric values.

Example

In the following example, assume that vnod was assigned a value earlier in the exec:

"v_fstatfs vnod st."
v_get

Function

v_get invokes the v_get callable service to return a vnode token for the file or
directory represented by the input FID within the mounted file system represented
by the input VFS token.

Parameters

vfstoken
  A variable name that contains the VFS token for the file system where the file
  identified by fid resides.

fid
  A variable that contains a file ID. File IDs are returned in the file attribute
  structure in the stem index ST_FID.

vntoken
  A variable name that stores the vnode token for the requested file.

Usage notes

1. The FID (file identifier) uniquely identifies a file in a particular mounted file
   system, and its validity persists across mounting and unmounting of the file
   system, as well as z/OS UNIX re-IPLs. This distinguishes the FID from the
   vnode token, which relates to a file in active use, and whose validity persists
   only until the token is released via the v_rel callable service.
   A server application uses v_get to convert an FID to a vnode token when it is
   preparing to use a file, since the vnode token identifies the file to the other
   services.
2. The FID for a file is returned in a stem variable by such services as v_rpn and
   v_lookup.
3. The caller is responsible for freeing vnode tokens returned by v_get by calling
   to the v_rel service when they are no longer needed.

Example

In the following example, assume that vfs and st.st_fid were assigned values earlier
in the exec:
"v_get vfs st.st_fid vnod"
Function

_v_getattr_ invokes the _v_getattr_ callable service to get the attributes of the file represented by _vntoken_.

Parameters

_vntoken_

A variable name that contains the vnode token of the file for which the attributes are returned.

_stem_

The name of a stem variable used to return the file attribute information. Upon return, _stem.0_ contains the number of attribute variables returned. To obtain the desired information, you can use a numeric value or the predefined variables beginning with ST_ used to derive the numeric value. See "stat" on page 176 for a list of the variables, or Appendix A for the numeric values.

Example

In the following example, assume that _vnod_ was assigned a value earlier in the exec:

"_v_getattr vnod attr."
Function

_v_link_ invokes the _v_link_ callable service to create a link to the file specified by
_vntoken_ in the directory specified by _vntoken2_. The link is a new name, _filename_,
identifying an existing file.

Parameters

_vntoken_
   A variable name that contains the vnode token for the file being linked to.

_filename_
   The new name for the existing file.

_vntoken2_
   A variable name that contains the vnode token for the directory to which
   _filename_ is to be added.

Usage notes

1. _v_link_ creates a link named _filename_ to an existing file specified by _vntoken_.
   This provides an alternative pathname for the existing file, so that you can
   access it by the old name or the new name. You can store the link under the
   same directory as the original file, or under a different directory on the same file
   system.
2. If the link is created successfully, the service routine increments the link count of
   the file. The link count shows how many links to a file exist. (If the link is not
   created successfully, the link count is not incremented.)
3. Links are not allowed to directories.
4. If the link is created successfully, the change time of the linked-to file is
   updated, as are the change and modification times of the directory that contains
   _filename_—that is, the directory that holds the link.

Example

In the following example, assume that _filetok, name_, and _dirtok_ were assigned
values earlier in the exec:

"v_link filetok (name) dirtok"
Function

v_lockctl invokes the v_lockctl callable service to control advisory byte-range locks on a file.

Note: All locks are advisory only. Client and local processes can use locks to inform each other that they want to protect parts of a file, but locks do not prevent I/O on the locked parts. A process that has appropriate permissions on a file can perform whatever I/O it chooses, regardless of which locks are set. Therefore, file locking is only a convention, and it works only when all processes respect the convention.

Parameters

v_command

The name of a predefined command variable that is used to control the lock. You can specify a numeric value (see Appendix A) or a predefined VL_command variable that derives the appropriate numeric value. The command variables are

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>VL_REGLOCKER</td>
<td>Register the lock server (locker).</td>
</tr>
<tr>
<td>VL_UNREGLOCKER</td>
<td>Unregister the locker.</td>
</tr>
<tr>
<td>VL_LOCK</td>
<td>Set a lock in a specified byte range.</td>
</tr>
<tr>
<td>VL_LOCKWAIT</td>
<td>Set a lock in a specified byte range or wait to set the lock until the byte range is free.</td>
</tr>
<tr>
<td>VL_UNLOCK</td>
<td>Unlock all locks in a specified byte range.</td>
</tr>
<tr>
<td>VL_QUERY</td>
<td>Query for lock information about a file.</td>
</tr>
<tr>
<td>VL_PURGE</td>
<td>Release all locks on all files, held by a locker or a group of lockers.</td>
</tr>
</tbody>
</table>

stem

The name of a stem variable that is the structure used to obtain information about the lock. To access the information, you can specify a numeric value (see Appendix A) or the predefined variables beginning with VL_ or L_ that derive the appropriate numeric value—for example, stem.vl_serverpid. The variables beginning with VL_ are:

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>VL_SERVERPID</td>
<td>Server's PID</td>
</tr>
<tr>
<td>VL_CLIENTPID</td>
<td>Server's client's process ID (PID)</td>
</tr>
<tr>
<td>VL_LOCKERTOK</td>
<td>Token for locker</td>
</tr>
<tr>
<td>VL_CLIENTTID</td>
<td>Client's thread's TID. The TID is the individual lock owner within a locker.</td>
</tr>
<tr>
<td>VL_OBJCLASS</td>
<td>The class for an object (a single locked file)—for example, HFS for an HFS file, MVS for an MVS data set, LFSESA for a LAN file server.</td>
</tr>
</tbody>
</table>
### v_lockctl

<table>
<thead>
<tr>
<th>Variable</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>VL_OBJID</td>
<td>The unique ID for an object (locked file) within its class. For an HFS file, the VL_OBJID contains the device number and FID of the file.</td>
</tr>
<tr>
<td>VL_OBJTOK</td>
<td>A token to identify the object (locked file) on a subsequent lock request.</td>
</tr>
<tr>
<td>VL_DOSMODE</td>
<td>DOS file-sharing field</td>
</tr>
<tr>
<td>VL_DOSACCESS</td>
<td>DOS file-sharing field</td>
</tr>
</tbody>
</table>

For a description of the variables beginning with L_, see ["f_getlk" on page 68].

### Usage notes

1. The v_lockctl service locks out other cooperating lockers from part of a file, so that the locker can read or write to that part of the file without interference from others.

2. Each locker must be registered before issuing any lock requests. On a REGLOCKER command, the caller must provide stem variables with these suffixes:
   - VL_SERVERPID
   - VL_CLIENTPID
   The VL_LOCKERTOK variable is returned to the caller; it is a token to identify the locker on subsequent lock requests.

3. On a QUERY, LOCK, LOCKWAIT, or UNLOCK command, the caller provides stem variables with these suffixes:
   - VL_LOCKERTOK
   - VL_CLIENTTID
   - VL_OBJCLASS
   - VL_OBJID
   - VL_OBJTOK (This is optional, but it will improve performance for multiple lock requests)

   To describe the byte range for the command, the caller must also provide stem variables with the following L_ suffixes:
   - QUERY L_TYPE, L_START, L_LEN, L_WHENCE
   - LOCK L_TYPE, L_START, L_LEN, L_WHENCE
   - LOCKWAIT L_TYPE, L_START, L_LEN, L_WHENCE
   - UNLOCK L_START, L_LEN, L_WHENCE

   The L_ variables are described in ["f_getlk" on page 68].

   VL_OBJTOK is returned to the caller; it is a token to identify the object on a subsequent lock request. On a QUERY, lock information describing a lock that would prevent the proposed lock from being set is returned to the caller.

4. Stem variables with the suffixes L_TYPE, L_START, and L_LEN are needed whether the request is for setting a lock, releasing a lock, or querying a particular byte range for a lock. L_WHENCE is always treated as SEEK_SET, the start of the file.

   The L_TYPE variable is used to specify the type of lock to be set or queried. (L_TYPE is not used to unlock.) You can use a numeric value (see Appendix A) or one of the following predefined variables used to derive the appropriate value:

<table>
<thead>
<tr>
<th>Type</th>
<th>Description</th>
</tr>
</thead>
<tbody>
<tr>
<td>F_RDLCK</td>
<td>A read lock, also known as a shared lock. This</td>
</tr>
</tbody>
</table>
**v_lockctl**

A type of lock specifies that the locker can read the locked part of the file, and other lockers cannot write on that part of the file in the meantime. A locker can change a held write lock, or any part of it, to a read lock, thereby making it available for other lockers to read. Multiple lockers can have read locks on the same part of a file simultaneously.

**F_WRLCK**

A write lock, also known as an exclusive lock. This type of lock indicates that the locker can write on the locked part of the file, without interference from other lockers. If one locker puts a write lock on part of a file, no other locker can establish a read lock or write lock on that same part of the file. A locker cannot put a write lock on part of a file if there is already a read lock on an overlapping part of the file, unless that locker is the only owner of that overlapping read lock. In such a case, the read lock on the overlapping section is replaced by the write lock being requested.

**F_UNLCK**

Returned on a query, when there are no locks that would prevent the proposed lock operation from completing successfully.

The L_WHENCE variable specifies how the byte range offset is to be found within the file; L_WHENCE is always treated as SEEK_SET, which stands for the start of the file.

The L_START variable is used to identify the part of the file that is to be locked, unlocked, or queried. The part of the file affected by the lock begins at this offset from the start of the file. For example, if L_START has the value 10, a lock request attempts to set a lock beginning 10 bytes past the start of the file.

**Note:** Although you cannot request a byte range that begins or extends beyond the beginning of the file, you can request a byte range that begins or extends beyond the end of the file.

The L_LEN variable is used to give the size of the locked part of the file, in bytes. The value specified for L_LEN cannot be negative. If a negative value is specified for L_LEN, a RETVAL of −1 and an EINVAL ERRNO are returned. If L_LEN is zero, the locked part of the file begins at L_START and extends to the end of the file.

The L_PID variable identifies the VL_CLIENTPID of the locker that holds the lock found on a query request, if one was found.

5. You can set locks by specifying a VL_LOCK as the command parameter. If the lock cannot be obtained, a RETVAL of −1 is returned along with an appropriate ERRNO and ERRNOJR.

You can also set locks by specifying a VL_LOCKWAIT as the command parameter. If the lock cannot be obtained because another process has a lock on all or part of the requested range, the LOCKWAIT request waits until the specified range becomes free and the request can be completed.

If a signal interrupts a call to the v_lockctl service while it is waiting in a LockWait operation, the function returns with a RETVAL of −1, and the ERRNO EINTR.
LockWait operations have the potential for encountering deadlocks. This happens when locker A is waiting for locker B to unlock a byte range, and B is waiting for A to unlock a different byte range. If the system detects that a LockWait might cause a deadlock, the v_lockctl service returns with a RETVAL of -1 and the ERRNO EDEADLK.

6. A process can determine locking information about a file using VL_QUERY as the command parameter. The stem should describe a lock operation that the caller would like to perform. When the v_lockctl service returns, the structure is modified to describe the first lock found that would prevent the proposed lock operation from finishing successfully.

If a lock is found that would prevent the proposed lock from being set, the query request returns a stem whose L_WHENCE value is always SEEK_SET, whose L_START value gives the offset of the locked portion from the beginning of the file, whose L_LEN value is set to the length of the locked portion of the file, and whose L_PID value is set to the ClientProcessID of the locker that is holding the lock. If there are no locks that would prevent the proposed lock operation from finishing successfully, the returned structure is modified to have an L_TYPE of F_UNLCK, but otherwise it remains unchanged.

7. A locker can have several locks on a file simultaneously but can have only one type of lock set on any given byte. Therefore, if a locker sets a new lock on part of a file that it had previously locked, the locker has only one lock on that part of the file, and the lock type is the one given by the most recent locking operation.

8. When a VL_UNLOCK command is issued to unlock a byte range of a file, all locks held by that locker within the specified byte range are released. In other words, each byte specified on an unlock request is freed from any lock that is held against it by the requesting locker.

9. Each locker should be unregistered when done issuing lock requests. On a VL_UNLOCKER command, the caller provides the stem variable VL_LOCKERTOK to identify the locker to unregister.

10. The VL_PURGE command releases all locks on all files, held by a locker or a group of lockers. The following stem variables are provided by the caller:
    - VL_SERVERPID
    - VL_CLIENTPID
    - VL_CLIENTTID

Example

This example illustrates several calls to v_lockctl to register a locker, lock a range, unlock a range, and unregister a locker:

```plaintext
/* rexx */
address syscall
'v_reg 2 RxLocker'    /* register server as a lock server */
/*****************************/
/* register locker */
/*****************************/
lk.vl_serverpid=0    /* use my pid as server pid */
lk.vl_clientpid=1    /* set client process id */
'l_lockctl' vl_reglocker 'lk.' /* register client as a locker */
c1Tok=lk.vl_lockertok /* save client locker token */
/*****************************/
/* lock a range */
/*****************************/
lk.vl_lockertok=c1Tok    /* set client locker token */
lk.vl_clienttid='thread1' /* invent a thread id */
lk.vl_objclass=1        /* invent an object class */
```
v_lockctl

lk.vl_objid='objectname' /* invent an object name */
lk.vl_objtok='' /* no object token */
lk.l_len=40 /* set length of range to lock */
lk.l_start=80 /* set start of range to lock */
lk.l_whence=seek_set /* start of range is absolute */
lk.l_type=f_wrlck /* set write lock */
'v_lockctl' vl_lock 'lk.' /* try to do the lock */
obj1=lk.vl_objtok /* keep returned object token */

/***********************************************************/
/* unlock a range */
/***********************************************************/
lk.l_lockertok=cltok /* set client locker token */
lk.l_clienttid='thread1' /* invent a thread id */
lk.l_objclass=1 /* invent an object class */
lk.l_objid='objectname' /* invent an object name */
lk.l_objtok=obj1 /* set object token */
lk.l_len=40 /* set length of range to lock */
lk.l_start=80 /* set start of range to lock */
lk.l_whence=seek_set /* start of range is absolute */
lk.l_type=f_unlck /* set unlock */
'v_lockctl' vl_unlock 'lk.' /* unlock the range */

/***********************************************************/
/* unregister locker */
/***********************************************************/
lk.l_lockertok=cltok /* set client locker token */
'v_lockctl' vl_unreglocker 'lk.' /* unregister client as a locker */
return

v_lockctl
Function

v_lookup invokes the v_lookup callable service to search a directory for a file. v_lookup accepts a vnode token representing a directory and a name identifying a file. If the file is found in the directory, a vnode token for the file and the attributes of the file are returned.

Parameters

vntoken
A variable name that contains the vnode token for the directory in which filename is looked up.

filename
The name of the file.

stem
The same file attribute information is returned in stem as if a v_getattr had been used on the file looked up. Upon return, stem.0 contains the number of attribute variables returned. To access the attribute values, you can use a numeric value or the predefined variables beginning with ST_ used to derive the numeric value. See "stat" on page 176 for a list of the predefined variables, or Appendix A for the numeric values.

vntoken2
The variable name for the buffer that, when returned, will contain the vnode token for the looked-up file.

Usage notes

The caller is responsible for freeing vnode tokens returned by the v_lookup service by calling to the v_rel service when they are no longer needed.

Example

In the following example, assume that dirtok and file were assigned values earlier in the exec:
"v_lookup dirtok (file) st. outtok"
Function

v_mkdir invokes the v_mkdir callable service to create a new empty directory in the directory represented by vntoken, with the permission specified in mode.

Parameters

  vntoken
  A variable name that contains the vnode token for the directory in which filename is to be created.

  directoryname
  The name of the directory.

  mode
  A three-digit number, corresponding to the access permission bits for the directory. Each digit must be in the range 0–7, and all three digits must be specified. For more information on permissions, see Appendix B.

  stem
  The same file attribute information is returned in stem as if a v_getattr had been used on the file specified. Upon return, stem.0 contains the number of attribute variables returned. To access the attribute information, you can use a numeric value or the predefined variables beginning with ST_. used to derive the numeric value. See [stat on page 176](#) for a list of the predefined variables, or Appendix A for the numeric values.

  vntoken2
  The variable name for the buffer that will contain the vnode token for the newly created directory.

Usage notes

1. If the directory that is named in directoryname already exists, the v_mkdir service returns a failing return code, and no vntoken2 is returned.
2. The caller is responsible for freeing vnode tokens returned by the v_mkdir service by calling to the v_rel service when they are no longer needed.

Example

In the following example, assume that dirtok, file, and perm were assigned values earlier in the exec:

"v_mkdir dirtok (file)" perm "st. newtok"
Function

v_read invokes the v_rdwr callable service to accept a vnode token representing a file and to read data from the file. The file attributes are returned when the read completes. The number of bytes read is returned in RETVAL.

Parameters

vntoken

A variable name that contains the vnode token for the file to be read.

variable

The name of the buffer into which data will be read.

length

The maximum number of characters to read. After the read completes, the length of variable is the number of bytes read. This value is also returned in RETVAL.

position

The file offset where the read is to begin, specified in bytes.

stem

The name of a stem variable used to return the file attribute information. Upon return, stem.0 contains the number of attribute variables returned. The same information is returned in stem as if a v_getattr had been used on the file. To obtain the attribute information, you can use a numeric value or the predefined variables beginning with ST_ used to derive the numeric value. See "stat" on page 176 for a list of the variables, or Appendix A for the numeric values.

access_check

Specify a 0 for no access check, or 1 to indicate the system is to check the user for read access to the file. The user is defined by the effective UID and GID. Setting the effective GID does not affect supplemental groups. Also, there is no support for altering the MVS user identity of the task using the ACEE (Accessor Environment Element).

Example

In the following example, assume that filetok, bytes, and pos were assigned values earlier in the exec:

"v_read filetok buffer" bytes pos stem. 0
Function

`v_readdir` invokes the `v_readdir` callable service to accept a vnode token representing a directory and return directory entries from this directory.

Parameters

**vntoken**
A variable name that contains the vnode token for the directory to be read.

**stem**
The name of a stem variable used to return the directory entries. Upon return, `stem.0` contains the number of directory entries returned. `stem.1` through `stem.n` (where `n` is the number of entries returned) each contain a directory entry.

**Note:** Only small directories can be read in a single call. To ensure that you read to the end of the directory, make calls to `v_readdir` until no entries are returned.

**start**
The number of the first directory entry to be returned. The numbers 0 and 1 both indicate that the read should start at the beginning of the directory.

**access_check**
Specify a 0 for no access_check, or 1 to specify that the system is to check the user for read access to the file. The user is defined by the effective UID and GID. Setting the effective GID does not affect supplemental groups. Also, there is no support for altering the MVS user identity of the task using the ACEE.

Example

In the following example, assume that `dirtok` was assigned a value earlier in the exec:

```
v_readdir dirtok dir 0
```
Function

v_readlink invokes the v_readlink callable service to read the symbolic link file represented by the vnode token and return its contents in variable.

Parameters

vntoken

A variable name that contains the vnode token for the symbolic link to be read.

The attribute stem returned on call to another function (for example, v_getattr) identifies whether the symbolic link is a link to an external name in the stem index ST_EXTLINK. An external name is the name of an object outside the HFS.

variable

The name of the buffer that, on return, contains the contents of the symbolic link.

Example

In the following example, assume that symtok was assigned a value earlier in the exec:

"v_readlink symtok link"
v_reg

v_reg

Function

v_reg invokes the v_reg callable service to register a process as a server. A process must be registered using this service before it can use any other vnode interface services.

Parameters

type
A numeric value that defines the type of server. You can specify:
  1 to indicate a file server
  2 to indicate a lock server

name
The name of the server, a character string up to 32 bytes in length. There are no restrictions on the name; for example, it does not have to be unique in the system.

Usage notes

1. Only a superuser can register a process as a server.
2. The DOMVS command uses the values supplied in type and name fields to display information about the currently active servers.

Example

To register a server:

`v_reg 1 "My REXX server"`


v_rel

Function

v_rel invokes the v_rel callable service to accept a vnode token representing a file descriptor for a file or directory, and to release that token.

Parameters

vntoken
A variable name that contains the vnode token for the file descriptor to be released.

Usage notes

1. The vnode token is no longer valid and cannot be used for subsequent requests after v_rel has successfully processed it.
2. This service must be used to release all vnode tokens obtained from other operations.

Example

In the following example, assume that vntok was assigned a value earlier in the exec:
"v_rel vntok"
v_remove

v_remove

v_remove invokes the v_remove callable service to remove a directory entry.

Parameters

vntoken
A variable name that contains the vnode token for the directory from which
filename is to be removed.

filename
The name for the directory entry. The directory entry could be identified by a
name for a file, the name of a hard link to a file, or the name of a symbolic link.

Usage notes

1. If the name specified refers to a symbolic link, the symbolic link file named by
filename is deleted.
2. If the v_remove service is successful and the link count becomes zero, the file
is deleted. The contents of the file are discarded, and the space it occupied is
freen for reuse. However, if another process (or more than one) has the file
open or has a valid vnode token when the last link is removed, the file contents
are not removed until the last process closes the file or releases the vnode
token.
3. When the v_remove service is successful in removing a directory entry and
decrementing the link count, whether or not the link count becomes zero, it
returns control to the caller with RETVAL 0. It updates the change and
modification times for the parent directory, and the change time for the file itself
(unless the file is deleted).
4. You cannot remove a directory using v_remove. To remove a directory, refer to
"v_rmdir" on page 280.

Example

In the following example, assume that dirtok and file were assigned values earlier in
the exec:
"v_remove dirtok (file)"
v_rename

Function

_vrename_vntoken_oldname_vntoken2_newname_

_vrename_ invokes the _vrename_ callable service to rename a file or directory to a new name.

Parameters

_vntoken_
A variable name that contains the vnode token for the directory that contains the filename _oldname_.

_vntoken2_
A variable name that contains the vnode token for the directory that is to contain the filename _newname_.

(oldname)
The existing name for the file or directory.

Usage notes

1. The _v_rename_ service changes the name of a file or directory from _oldname_ to _newname_. When renaming completes successfully, the change and modification times for the parent directories of _oldname_ and _newname_ are updated.

2. The calling process needs write permission for the directory containing _oldname_ and the directory containing _newname_. If _oldname_ and _newname_ are the names of directories, the caller does not need write permission for the directories themselves.

3. Renaming files: If _oldname_ and _newname_ are links referring to the same file, _v_rename_ returns successfully and does not perform any other action.

   _oldname_ is the name of a file, _newname_ must also name a file, not a directory. If _newname_ is an existing file, it is unlinked. Then the file specified as _oldname_ is given _newname_. The pathname _newname_ always stays in existence; at the beginning of the operation, _newname_ refers to its original file, and at the end, it refers to the file that used to be _oldname_.

4. Renaming directories: If _oldname_ is the name of a directory, _newname_ must also name a directory, not a file. If _newname_ is an existing directory, it must be empty, containing no files or subdirectories. If it is empty, it is removed.

   _newname_ cannot be a directory under _oldname_, that is, the old directory cannot be part of the pathname prefix of the new one.

Example

In the following example, assume that _olddir_, _oldfile_, _newdir_, and _newfile_ were assigned values earlier in the exec:

"v_rename olddir (oldfile) newdir (newfile)"
v_rmdir

Function

v_rmdir invokes the v_rmdir callable service to remove an empty directory.

Parameters

vntoken
   A variable name that contains the vnode token for the directory from which
dirname is to be removed.

dirname
   The name of the empty directory to be removed.

Usage notes

1. The directory specified by dirname must be empty.
2. If the directory is successfully removed, the change and modification times for
   the parent directory are updated.
3. If any process has the directory open when it is removed, the directory itself is
   not removed until the last process closes the directory. New files cannot be
   created under a directory that is removed, even if the directory is still open.

Example

In the following example, assume that dirtok and dirname were assigned values
earlier in the exec:
"v_rmdir dirtok (dirname)"
v_rpn invokes the v_rpn callable service to accept a pathname of a file or directory and return a vnode token that represents this file or directory and the VFS token that represents the mounted file system that contains the file or directory.

Parameters

- **pathname**
  The absolute pathname to be resolved, specified as a string.

- **vfstoken**
  The name of a variable in which the VFS token for the resolved file is stored.

- **vntoken**
  The name of a variable in which the vnode token for the resolved file is stored.

- **stem**
  The name of a stem variable used to return the mount entry for the file system. To access mount table information, you can use a numeric value or the predefined variables beginning with MNTE_ used to derive the numeric value. See "getmntent" on page 99 for a description of the MNTE_ variables; see Appendix A for the numeric values.

- **stem2**
  Upon return, stem2.0 contains the number of attribute variables returned. The same information is returned in stem2 as if a v_getattr had been used on the file that was just resolved. You can use the predefined variables beginning with ST_ to access those respective values. For example, stem2.st_size accesses the file size. See "stat" on page 176 for a description of the ST_ variables.

Usage notes

1. The mount point pathname is not available in the MNTE_ structure returned by the variable stem2.mnte_path.
2. The caller is responsible for freeing vnode tokens returned by the v_rpn service, by calling to the v_rel service when they are no longer needed.

Example

In the following example, assume that path was assigned a value earlier in the exec:

"v_rpn (path) vfstok filetok mnt. attr."
Function

v_setattr invokes the v_setattr callable service to set the attributes associated with
the file represented by the vnode token. You can change the mode, owner, access
time, modification time, change time, reference time, audit flags, general attribute
flags, and file size.

Parameters

vntoken
A variable name that contains the vnode token for the file for which the
attributes are to be set.

attribute_list
A list of attributes to be set and their values. The attributes are expressed either
as numeric values or as the predefined variables beginning with ST_, followed
by arguments for that attribute. For the predefined variables beginning with ST_,
see "chattr" on page 38, for the numeric values, see Appendix A.

Usage notes

For usage notes, see "chattr" on page 38.

Example

In the following example, assume that vntok was assigned a value earlier in the
exec. This example truncates a file to 0 length and sets the mode to 600:
"v_setattr vntok" st_size 0 st_mode 600
Function

_v_symlink invokes the _v_symlink callable service to create a symbolic link to a pathname or external name. The contents of the symbolic link file is _pathname_.

Parameters

_vntoken
A variable name for the directory that contains the vnode token in which _filename_ is being created.

_filename_
The name for the symbolic link.

_pathname_
The absolute or relative pathname of the file you are linking to (the contents of the symbolic link).

_extlink_
Specify 1 if this is a symbolic link to an external name rather than to a pathname in the file hierarchy. An external name is the name of an object outside of the file hierarchy.

Usage notes

1. Like a hard link (described in [v_link](#) on page 265), a symbolic link allows a file to have more than one name. The presence of a hard link guarantees the existence of a file, even after the original name has been removed. A symbolic link, however, provides no such assurance; in fact, the file identified by _pathname_ need not exist when the symbolic link is created. In addition, a symbolic link can cross file system boundaries, and it can refer to objects outside of a hierarchical file system.

2. When a component of a pathname refers to a symbolic link (but not an external symbolic link) rather than to a directory, the pathname contained in the symbolic link is resolved. For _v_rpn_ or other _z/OS UNIX_ callable services, a symbolic link in a pathname parameter is resolved as follows:
   - If the pathname in the symbolic link begins with / (slash), the symbolic link pathname is resolved relative to the process root directory.
   - If the pathname in the symbolic link does not begin with /, the symbolic link pathname is resolved relative to the directory that contains the symbolic link.
   - If the symbolic link is not the last component of the original pathname, remaining components of the original pathname are resolved from there.
   - When a symbolic link is the last component of a pathname, it may or may not be resolved. Resolution depends on the function using the pathname. For example, a rename request does not have a symbolic link resolved when it appears as the final component of either the new or old pathname. However, an open request does have a symbolic link resolved when it appears as the last component.
v_symmlink

- When a slash is the last component of a pathname, and it is preceded by a symbolic link, the symbolic link is always resolved.
- Because it cannot be changed, the mode of a symbolic link is ignored during the lookup process. Any files and directories to which a symbolic link refers are checked for access permission.

3. The external name contained in an external symbolic link is not resolved. The filename cannot be used as a directory component of a pathname.

Example

In the following example, assume that dirtok, file, and linkname were assigned values earlier in the exec:

"v_symlink dirtok (file) (linkname)"
Function

v_rdwr invokes the v_rdwr callable service to accept a vnode token representing a file and to write data to the file. The number of bytes written and the file attributes are returned in RETVAL when the write completes.

Parameters

vntoken
A variable name that contains the vnode token for the file to be written.

variable
The name of the buffer from which data is to be written.

length
The number of characters to write.

position
The file offset where the write is to start from, specified in bytes.

stem
The name of a stem variable used to return the file attributes. Upon return, stem.0 contains the number of attribute variables returned. The same information is returned in stem as if a v_getattr was used on the file. To access the file attributes, you can use a numeric value or the predefined variables beginning with ST_ used to derive the numeric value. See "stat" on page 176 for more information about the ST_ predefined variables; see Appendix A for the numeric values.

access_check
Specify 0 for no access check, or 1 for the system to check the user for read access to the file. The user is defined by the effective UID and GID. Setting the effective GID does not affect supplemental groups. Also, there is no support for altering the MVS user identity of the task using the ACEE.

Example

In the following example, assume that filetok, buf, and pos were assigned values earlier in the exec:

"v_write filetok buf" length(buf) pos
Chapter 8. Examples: Using virtual file system syscall commands

These are examples of REXX programs that use the virtual file system syscall commands.

List the files in a directory

Given a directory pathname, this example lists the files in the directory.

```rexx
/* rexx */
parse arg dir                           /* take directory path as argument */
if dir=' ' then do
    say 'directory argument required'
    return
end

call syscalls 'ON'
address syscall
've_reg 1 dirlist'                        /* register as a file server */
if retval=-1 then do
    say 'error registering as a server - error codes:' errno errnojr
    return
end

've_rpn (dir) vfs vn mnt. st.'            /* resolve the directory pathname */
if retval=-1 then do
    say 'error resolving path' dir ' - error codes:' errno errnojr
    return
end

i=1                                /* next dir entry to read is 1 */
do forever                          /* loop reading directory */
    'v_readdir vn d.' i               /* read starting at next entry */
    if retval=-1 then do
        say 'error reading directory - error codes:' errno errnojr
        leave
    end
    if d.0=0 then leave                /* if nothing returned then done */
    do j=1 to d.0                      /* process each entry returned */
        say d.j
    end
    i=1+d.0                            /* set index to next entry */
end
've_rel vn'                          /* release the directory vnode */
return
```

Remove a file or empty directory

Given a directory pathname and the filename to delete, this example removes the file or an empty directory.

```rexx
/* rexx */
if __argv.0<>3 then                   /* check for right number of args */
do
    say 'directory and filename required'
    return
end

/* __argv.1 is program name */
dir= __argv.2                        /* 1st arg is directory pathname */
file= __argv.3                        /* 2nd arg is file name */
```
call syscalls 'ON'
address syscall
"v_reg 1 'remove file'" /* register as a file server */
if retval=-1 then
  do
    say 'error registering as a server - error codes:' errno errnojr
    return
  end
'v_rpn (dir) vfs vn mnt. st.' /* resolve the directory pathname */
if retval=-1 then
  do
    say 'error resolving path' dir ' - error codes:' errno errnojr
    return
  end
'v_lookup vn (file) fst. fvn' /* look up the file */
if retval=-1 then
  do
    say "error locating file" file
    say " in directory" dir
    say " error codes:" errno errnojr
  end
else
  do
    if fst.st_type=s_isdir then /* if the file is a directory */
      'v_rmdir vn (file)' /* then delete it with v_rmdir */
    else
      'v_remove vn (file)' /* else delete it with v_remove */
    if retval=-1 then
      say 'error deleting file - error codes:' errno errnojr
      'v_rel fvn' /* release the file vnode */
    end
  'v_rel vn' /* release the directory vnode */
return
Appendix A. REXX predefined variables

Predefined variables make symbolic references easier and more consistent. Instead of using a numeric value, you can use the predefined variable that will derive that numeric value. The following list shows the data type and numeric value for each predefined variable. Each variable is discussed in the section about the syscall command with which it can be used. Most variable names correspond to the POSIX-defined names in the C runtime library include files (also known as header files).

Except for the error numbers and signal numbers, all variables contain an underscore. This is also true of most of the POSIX names in the C include files.

The data types are:
- **Bin** Binary: 2-byte hexadecimal format
- **Char** Character
- **Dec** Decimal
- **Hex** Hexadecimal: 4-byte hexadecimal format
- **Tok** Token
- **Oct** Octal

The predefined variables listed alphabetically are:

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<thead>
<tr>
<th>Variable</th>
<th>Data Type</th>
<th>Numeric Value</th>
</tr>
</thead>
<tbody>
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Table 12. List of predefined variables (continued)

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Table 12. List of predefined variables (continued)

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Table 12. List of predefined variables (continued)

<table>
<thead>
<tr>
<th>Variable</th>
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<th>Numeric Value</th>
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<td>W_UNTRACED</td>
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<tr>
<td>X_OK</td>
<td>Dec</td>
<td>1</td>
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</table>
Appendix B. Setting permissions for files and directories

Typically, octal permissions are specified with three or four numbers, in these positions:

1234

Each position indicates a different type of access:
- Position 1 is the bit that sets permission for set-user-ID on access, set-group-ID on access, or the **sticky bit**. Specifying this bit is optional.
- Position 2 is the bit that sets permissions for the owner of the file. Specifying this bit is required.
- Position 3 is the bit that sets permissions for the group that the owner belongs to. Specifying this bit is required.
- Position 4 is the bit that sets permissions for others. Specifying this bit is required.

Position 1

Specifying the bit in position 1 is optional. You can specify one of these values:

0   Off
1   Sticky bit on
2   Set-group-ID-on execution
3   Set-group-ID-on execution and set the sticky bit on
4   Set-user-ID on execution
5   Set-user-ID on execution and set the sticky bit on.
6   Set-user-ID and set-group-ID on execution
7   Set-user-ID and set-group-ID on execution and set the sticky bit on

Positions 2, 3, and 4

Specifying these bits is required. For each type of access—owner, group, and other—there is a corresponding octal number:

0   No access (---)
1   Execute-only access (-x-)
2   Write-only access (-w-)
3   Write and execute access (-wx)
4   Read-only access (r--)
5   Read and execute access (r-x)
6   Read and write access (rw-)
7   Read, write, and execute access (rwx)

To specify permissions for a file or directory, you use at least a 3-digit octal number, omitting the digit in the first position. When you specify just three digits, the first digit describes owner permissions, the second digit describes group permissions, and the third digit describes permissions for all others. When the first digit is not
set, some typical 3-digit permissions are specified in octal as shown in Figure 1.

<table>
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<tr>
<th>Octal Number</th>
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<td>666</td>
<td>owner (rw-)</td>
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<tr>
<td></td>
<td>group (rw-)</td>
</tr>
<tr>
<td></td>
<td>other (rw-)</td>
</tr>
<tr>
<td>700</td>
<td>owner (rwx)</td>
</tr>
<tr>
<td></td>
<td>group (---)</td>
</tr>
<tr>
<td></td>
<td>other (---)</td>
</tr>
<tr>
<td>755</td>
<td>owner (rwx)</td>
</tr>
<tr>
<td></td>
<td>group (r-x)</td>
</tr>
<tr>
<td></td>
<td>other (r-x)</td>
</tr>
<tr>
<td>777</td>
<td>owner (rwx)</td>
</tr>
<tr>
<td></td>
<td>group (rwx)</td>
</tr>
<tr>
<td></td>
<td>other (rwx)</td>
</tr>
</tbody>
</table>

Figure 1. Three-digit permissions specified in octal

**Example: using BITOR and BITAND to set mode bits**

To set a file's mode bits, use the REXX functions BITOR() and BITAND() with the octal numbers.

For example, if you have obtained a file's permission bits and want to use `chmod` to turn on the write bits, you could code:

```rexx
'chmod (file) BITOR(st.st_mode, 222)
```

To turn the same bits off, you could code:

```rexx
'chmod (file) BITAND(st.st_mode, 555)
```
Appendix C. Accessibility

Accessibility features help a user who has a physical disability, such as restricted mobility or limited vision, to use software products successfully. The major accessibility features in z/OS enable users to:

- Use assistive technologies such as screen readers and screen magnifier software
- Operate specific or equivalent features using only the keyboard
- Customize display attributes such as color, contrast, and font size

Using assistive technologies

Assistive technology products, such as screen readers, function with the user interfaces found in z/OS. Consult the assistive technology documentation for specific information when using such products to access z/OS interfaces.

Keyboard navigation of the user interface

Users can access z/OS user interfaces using TSO/E or ISPF. Refer to z/OS TSO/E Primer, z/OS TSO/E User’s Guide, and z/OS ISPF User’s Guide Vol I for information about accessing TSO/E and ISPF interfaces. These guides describe how to use TSO/E and ISPF, including the use of keyboard shortcuts or function keys (PF keys). Each guide includes the default settings for the PF keys and explains how to modify their functions.

z/OS information

z/OS information is accessible using screen readers with the BookServer/Library Server versions of z/OS books in the Internet library at: http://www.ibm.com/systems/z/os/zos/bkserv/
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